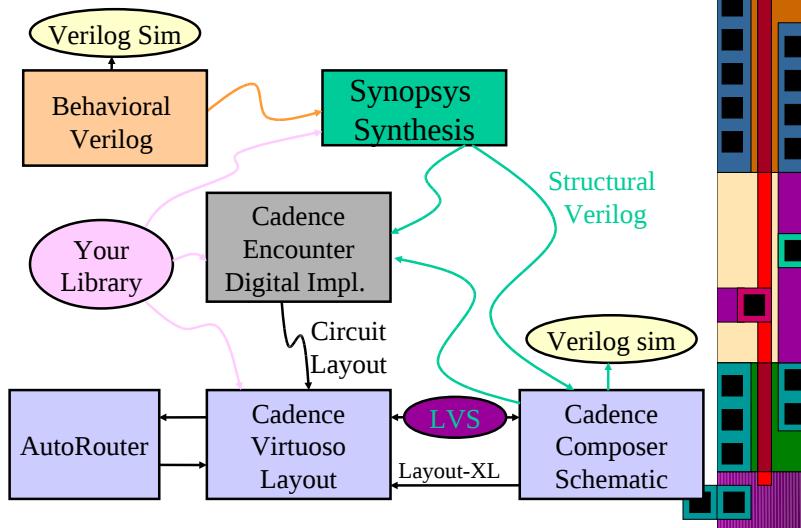


Digital VLSI ASIC Tool Suite

Verilog is the Key Tool



- ▶ Behavioral Verilog is synthesized into Structural Verilog
- ▶ Structural Verilog represents net-lists
 - ▶ From Behavioral
 - ▶ From Schematics
 - ▶ High-level (Synthesizer will flatten these)
- ▶ Verilog is used for testing all designs
 - ▶ Behavioral & Structural & Schematic & High-level
 - ▶ NC_Verilog, vcs (Synopsys Verilog simulator), modelsim (Mentor Verilog simulator)

Verilog has a Split Personality

- ▶ Hardware Description Language (HDL)
 - ▶ Reliably & Readably
 - ▶ Create hardware
 - ▶ Document hardware
- ▶ Testbench creation language
 - ▶ Create external test environment
 - ▶ Time & Voltage
 - ▶ Files & messages
- ▶ Are these two tasks
 - ▶ Related?
 - ▶ Compatible?

Verilog as HDL

- ▶ Want high level modeling
 - ▶ unification at all levels
 - ▶ from fast functional simulation, accurate device simulation
 - ▶ support simulation based validation (verification?)
- ▶ How could we do this?
 - ▶ behavioral model mapped to transistors
 - ▶ pragmas: throughput, latency, cycle time, power...
- ▶ Reality
 - ▶ we rely on designers to do most of these xforms
 - ▶ therefore:
 - ▶ different algorithms => try before you buy...
 - ▶ use only a subset of the language.
 - ▶ RTL and schematic design used to support Verilog
 - ▶ System-C and other HLD models for co-simulation, etc.

Synthesis

This lecture is only about synthesis...

Quick Review

```

Module name (args...);
begin
  parameter ...; // define parameters
  input ...; // define inputs
  output ...; // define outputs
  wire ...; // internal wires
  reg ...; // internal regs, possibly output
  // the parts of the module body are
  // executed concurrently
  <continuous assignments>
  <always blocks>
endmodule
  
```

Quick Review

- ▶ Continuous assignments to `wire` vars
 - ▶ `assign variable = exp;`
 - ▶ Results in combinational logic
- ▶ Procedural assignment to `reg` vars
 - ▶ Always inside procedural blocks (`always` blocks in particular for synthesis)
 - ▶ blocking
 - ▶ `variable = exp;`
 - ▶ non-blocking
 - ▶ `variable <= exp;`
- ▶ Can result in combinational or sequential logic

Verilog Description Styles

- ▶ Verilog supports a variety of description styles
 - ▶ Structural
 - ▶ explicit structure of the circuit
 - ▶ e.g., each logic gate instantiated and connected to others
 - ▶ Behavioral
 - ▶ program describes input/output behavior of circuit
 - ▶ many structural implementations could have same behavior
 - ▶ e.g., different implementation of one Boolean function

Synthesis: Data Types

- ▶ Possible Values (`wire` and `reg`):
 - ▶ `0`: logic 0, false
 - ▶ `1`: logic 1, true
 - ▶ `Z`: High impedance
- ▶ Digital Hardware
 - ▶ The domain of Verilog
 - ▶ Either `logic` (gates)
 - ▶ Or `storage` (registers & latches)
- ▶ Verilog has two relevant data types
 - ▶ `wire`
 - ▶ `reg`

Synthesis: Data Types

- ▶ Register declarations
 - ▶ `reg a;` // a scalar register
 - ▶ `reg [3:0] b;` // a 4-bit vector register
 - ▶ `output g;` // an output can be a reg
 - ▶ `reg g;`
- ▶ Wire declarations
 - ▶ `wire d;` // a scalar wire
 - ▶ `wire [3:0] e;` // a 4-bit vector wire
 - ▶ `output f;` // an output can be a wire

Parameters

- ▶ Used to define constants
 - ▶ `parameter size = 16, foo = 8;`
 - ▶ `wire [size-1:0] bus;` // defines a 15:0 bus

Synthesis: Assign Statement

- ▶ The `assign` statement creates combinational logic
 - ▶ `assign LHS = expression;`
 - ▶ `LHS` can only be wire type
 - ▶ `expression` can contain either wire or reg type mixed with operators
 - ▶ `wire a,c; reg b; output out;`
`assign a = b & c;`
 - ▶ `assign out = ~(a & b);` // output as wire
 - ▶ `wire [15:0] sum, a, b;`
`wire cin, cout;`
`assign {cout,sum} = a + b + cin;`

Synthesis: Basic Operators

- ▶ Bit-Wise Logical
 - ▶ \sim (not), $\&$ (and), $|$ (or), \wedge (xor), $\sim\wedge$ or $\sim\wedge$ (xnor)
 - ▶ Simple Arithmetic Operators
 - ▶ Binary: $+$, $-$
 - ▶ Unary: $-$
 - ▶ Negative numbers stored as 2's complement
 - ▶ Relational Operators
 - ▶ $<$, $>$, \leq , \geq , $=$, \neq
 - ▶ Logical Operators
 - ▶ $!$ (not), $\&\&$ (and), $\|$ (or)
- ```
assign a = (b > 4'b0110) && (c <= 4'd5);
assign a = (b > 4'b0110) && !(c > 4'd5);
```

## Synthesis: Operand Length

- ▶ When operands are of unequal bit length, the shorter operator is zero-filled in the most significant bit position

```
wire [3:0] sum, a, b; wire cin, cout, d, e, f, g;

assign sum = f & a;
assign sum = f | a;
assign sum = {d, e, f, g} & a;
assign sum = {4{f}} | b;
assign sum = {4{f == g}} & (a + b);
assign sum[0] = g & a[2];
assign sum[2:0] = {3{g}} & a[3:1];
```

## Synthesis: More Operators

- ▶ Concatenation
    - ▶  $\{a,b\}$   $\{4\{a==b\}\}$  { a,b,4'b1001,{4{a==b}} }
  - ▶ Shift (logical shift)
    - ▶  $<<$  left shift
    - ▶  $>>$  right shift
- ```
assign a = b >> 2; // shift right 2, division by 4  
assign a = b << 1; // shift left 1, multiply by 2
```
- ▶ Arithmetic
 - assign a = b * c; // multiply b times c
 - assign a = b * 'd2; // multiply b times constant (2)
 - assign a = b / 'b10; // divide by 2 (constant only)
 - assign a = b % 'h3; // b modulo 3 (constant only)

Synthesis: Operand Length

- ▶ Operator length is set to the longest member (both RHS & LHS are considered). Be careful.

```
wire [3:0] sum, a, b; wire cin, cout, d, e, f, g;  
wire[4:0]sum1;  
  
assign {cout,sum} = a + b + cin;  
assign {cout,sum} = a + b + {4'b0,cin};  
  
assign sum1 = a + b;  
assign sum = (a + b) >> 1; // what is wrong?
```

Synthesis: Extra Operators

- ▶ Funky Conditional
 - ▶ $cond_exp ? true_expr : false_expr$
 - ▶ `wire [3:0] a,b,c; wire d;`
 - ▶ `assign a = (b == c) ? (c + 'd1): '05;`
 - ▶ Reduction Logical
 - ▶ Named for impact on recreational time
 - ▶ Unary operators that perform bit-wise operations on a single operand, reduce it to one bit
 - ▶ $\&$, $\sim\&$, $|$, $\sim|$, \wedge , $\sim\wedge$, $\wedge\wedge$
- ```
assign d = &a || ~^a ^ ~^b ^ ~^c;
```

## Synthesis: Assign Statement

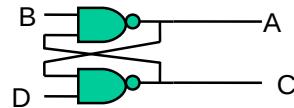
- ▶ The assign statement is sufficient to create all combinational logic
- ▶ What about this:

```
assign a = ~(b & c);
assign c = ~(d & a);
```

## Synthesis: Assign Statement

- ▶ The assign statement is sufficient to create all combinational logic
- ▶ What about this:

```
assign a = ~(b & c);
assign c = ~(d & a);
```



## Simple Behavioral Module

```
// Behavioral model of NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;
 assign out = ~(in1 & in2);
endmodule
```

## Simple Behavioral Module

```
// Behavioral model of NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;
 // Uses Verilog builtin nand function
 // syntax is func id (args);
 nand i0(out, in1, in2);
endmodule
```

## Simple Structural Module

```
// Structural Module for NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;
 wire w1;
 // call existing modules by name
 // module-name ID (signal-list);
 AND2X1 u1(w1, in1, in2);
 INVX1 u2(out, w1);
endmodule
```

## Simple Structural Module

```
// Structural Module for NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;
 wire w1;
 // call existing modules by name
 // module-name ID (signal-list);
 // can connect ports by name...
 AND2X1 u1(.Q(w1), .A(in1), .B(in2));
 INVX1 u2(.A(w1), .Q(out));
endmodule
```

## Primitive Gates

- ▶ Multiple input gates
  - ▶ `<gatename> [delay] [id] (out, in1, in2, in3...);`
  - ▶ and, or, nand, nor, xor, xnor
- ▶ Multiple output gates
  - ▶ `<gatename> [delay] [id] (out1, out2, ... outn, in);`
  - ▶ buf, not
- ▶ Tristate gates
  - ▶ `<gatename> [delay] [id] (out, in, ctrl);`
  - ▶ bufif1, bufif0,notif1, notif0

## Primitive Gates

- Delay: three types for gates
  - #(delaytime) same delay for all transitions
  - #(rise,fall) different delay for rise and fall
  - #(rise, fall, turnoff) for tristate gates
- Each delay number can be:
  - single number i.e. #(2) or #(2,3)
  - min/typ/max triple i.e. #(2:3:4) or #(2:3:4, 3:2:5)

## Primitive Gates

- and (out, a, b);
- nand i0 (out a b c d e f g);
- xor #(2,3) (out a b c);
- buf (Y A);
- buf #(2:3:4, 3:4;5) \_i1 (y, a);
- bufif1 (out, in, ctl);
- notif0 #(1, 2, 3) (Y, A, S);

## Primitive Gates

- OR – you can skip the delays on each gate, and use a specify block for the whole module
  - Specifies from module input to module outputs
  - Outputs must be driven by a primitive gate
  - The syntax defines the delay for each path from input to output

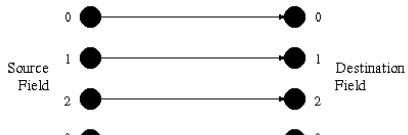
## Simple Behavioral Module

```
// Behavioral model of NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;
 nand _i0(out, in1, in2);

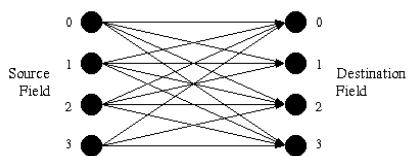
 // include specify block for timing
 specify
 (in1 => out) = (1.0, 1.0);
 (in2 => out) = (1.0, 1.0);
 endspecify
endmodule
```

## Specify Block Types

- Parallel Connection => (one to one)



- Full Connection \*-> (one to many)



## Parallel Specify

```
module A (q, a, b, c, d)
 input a, b, c, d;
 output q;
 wire e, f;

 // specify block containing delay statements
 specify
 (a => q) = 6; // delay from a to q
 (b => q) = 7; // delay from b to q
 (c => q) = 7; // delay from c to q
 (d => q) = 6; // delay from d to q
 endspecify

 // module definition
 or o1(e, a, b);
 or o2(f, c, d);
 exor ex1(q, e, f);
endmodule
```

## Full Specify

```
module A (q, a, b, c, d)
 input a, b, c, d;
 output q;
 wire e, f;

 // specify block containing full connections
 specify
 (a, d => q) = 6; // delay from a and d to q
 (b, c => q) = 7; // delay from b and c to q
 endspecify

 // module definition
 or o1(e, a, b);
 or o2(f, c, d);
 exor ex1(q, e, f);
endmodule
```

## Full Specify

```
// a[63:0] is a 64 bit input register and
// q[7:0] is an 8 bit output register
// this would require 64 x 8 = 512 parallel
// connections, but only 1 full

specify
 (a *-> q) = 8; // equivalent to 512 parallel
 connections
endspecify
```

## Specify needs a gate!

```
module DCX1 (CLR, D, CLK, Q);
 input CLR, D, CLK;
 output Q;
 reg Q_i;

 always @ (posedge CLK or negedge CLR)
 if (CLR == 0) Q_i = 1'b0;
 else Q_i = D;

 buf _i0 (Q, Q_i);

 specify
 (CLK => Q) = (1.0, 1.0);
 (CLR => Q) = (1.0, 1.0);
 $setuphold (posedge CLK, D, 0.1, 0.0);
 $recovery (negedge CLR, posedge CLK, 0.0);
 endspecify
endmodule
```

## Simple Behavioral Module

```
// Behavioral model of NAND gate
module NAND (out, in1, in2);
 output out;
 input in1, in2;

 nand _i0(out, in1, in2);

 // include specify block for timing
 specify
 (in1 => out) = (1.0, 1.0);
 (in2 => out) = (1.0, 1.0);
 endspecify
endmodule
```

## Procedural Assignment

- ▶ Assigns values to **register** types
- ▶ They involve data storage
  - ▶ The register holds the value until the next procedural assignment to that variable
- ▶ They occur only within procedural blocks
  - ▶ **initial** and **always**
  - ▶ **initial** is NOT supported for synthesis!
- ▶ They are triggered when the flow of execution reaches them

## Always Blocks

- ▶ When is an always block executed?
  - ▶ **always**
    - ▶ Starts at time 0
  - ▶ **always @ (a or b or c)**
    - ▶ Whenever there is a change on a, b, or c
    - ▶ Used to describe combinational logic
  - ▶ **always @ (posedge foo)**
    - ▶ Whenever foo goes from low to high
    - ▶ Used to describe sequential logic
  - ▶ **always @ (negedge bar)**
    - ▶ Whenever bar goes from high to low

## Synthesis: Always Statement

- The always statement creates...
  - `always @sensitivity LHS = expression;`
    - @sensitivity controls *when*
    - LHS can only be reg type
    - expression can contain either wire or reg type mixed with operators
  - Logic
    - `reg c, b; wire a;`
    - `always @(a, b) c = ~(a & b);`
    - `always @* c = ~(a & b);`
  - Storage
    - `reg Q; wire clk;`
    - `always @(posedge clk) Q <= D;`

## Procedural Control Statements

- Conditional Statement
  - `if ( <expression> ) <statement>`
  - `if ( <expression> ) <statement>`  
`else <statement>`
    - “else” is always associated with the closest previous if that lacks an else.
  - You can use begin-end blocks to make it more clear
    - `if (index >0)`  
`if (rega > regb)`  
`result = rega;`  
`else result = regb;`

## Multi-Way Decisions

- Standard if-else-if syntax

```
If (<expression>)
 <statement>
else if (<expression>)
 <statement>
else if (<expression>)
 <statement>
else <statement>
```

## Procedural NAND gate

```
// Procedural model of NAND gate
module NAND (out, in1, in2);
 output out;
 reg out;
 input in1, in2;
 // always executes when in1 or in2
 // change value
 always @ (in1 or in2)
 begin
 out = ~(in1 & in2);
 end
endmodule
```

## Procedural NAND gate

```
// Procedural model of NAND gate
module NAND (out, in1, in2);
 output out;
 reg out;
 input in1, in2;
 // always executes when in1 or in2
 // change value
 always @ (in1 or in2)
 begin
 out <= ~(in1 & in2);
 end
endmodule
```

*Is out combinational?*

## Synthesis: NAND gate

```
input in1, in2;
reg n1, n2; // is this a flip-flop?
wire n3, n4;

always @ (in1 or in2) n1 = ~(in1 & in2);
always @* n2 = ~(in1 & in2);
assign n3 = ~(in1 & in2);
nand u1(n4, in1, in2);
```

- Notice always block for combinational logic
  - Full sensitivity list, but `@*` works (2001 syntax)
  - Can then use the always goodies
  - Is this a good coding style?

## Procedural Assignments

- Assigns values to `reg` types
  - Only useable inside a procedural block Usually synthesizes to a register
    - But, under the right conditions, can also result in combinational circuits
- Blocking** procedural assignment
  - LHS = timing-control exp `a = #10 1;`
  - Must be executed before any assignments that follow (timing control is optional)
  - Assignments proceed in order even if no timing is given
- Non-Blocking** procedural assignment
  - LHS <= timing-control exp `b <= 2;`
  - Evaluated simultaneously when block starts
  - Assignment occurs at the end of the (optional) time-control

## Procedural Synthesis

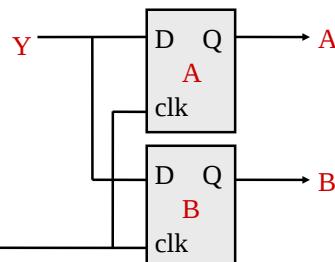
- Synthesis ignores all that timing stuff
- So, what does it mean to have blocking vs. non-blocking assignment for synthesis? How does this relate to reg?

`begin`      `begin`  
`A=B;`      `A<=B;`  
`B=A;`      `B<=A;`  
`end`      `end`

`begin`      `begin`  
`A=Y;`      `A<=Y;`  
`B=A;`      `B<=A;`  
`end`      `end`

## Synthesized Circuits

- `begin`  
`A = Y;`  
`B = A;`  
`end`
- `begin`  
`A <= Y;`  
`B <= A;`  
`end`
- `begin`  
`B = A;`  
`A = Y;`  
`end`



## Synthesized Circuits

```
always @ (posedge clk)
begin
 A = Y;
 B = A;
end
```

```
always @ (posedge clk)
begin
 B = A;
 A = Y;
end
```

```
always @ (posedge clk)
begin
 A <= Y;
 B <= A;
end
```

```
always @ (posedge clk)
begin
 B <= A;
 A <= Y;
end
```

## Assignments and Synthesis

- Note that different circuit structures result from different types of procedural assignments
  - Therefore you can't mix assignment types in the same `always` block
  - Non-blocking is often a better model for hardware
    - Real hardware is often concurrent...
  - Blocking is often better for setting subsets of signals
    - Set them to defaults at beginning, then reset only the ones that change

- Using continuous assignment
  - Concurrent execution of assignments

```
Module comp (a, b, Cgt, Clt, Cne);
parameter n = 4;
input [n-1:0] a, b;
output Cgt, Clt, Cne;
 assign Cgt = (a > b);
 assign Clt = (a < b);
 assign Cne = (a != b);
endmodule
```

## Comparator Example

- ▶ Using procedural assignment
  - ▶ Non-blocking assignment implies concurrent

```
Module comp (a, b, Cgt, Clt, Cne);
parameter n = 4;
input [n-1:0] a, b;
output Cgt, Clt, Cne;
reg Cgt, Clt, Cne;
always @(a or b)
begin
 Cgt <= (a > b);
 Clt <= (a < b);
 Cne <= (a != b);
end
endmodule
```

## Modeling a Flip Flop

- ▶ Use an **always** block to wait for clock edge

```
Module dff (clk, d, q);
 input clk, d;
 output q;
 reg q;
 always @ (posedge clk)
 d = q;
endmodule
```

## Synthesis: Always Statement

- ▶ This is a simple D Flip-Flop

```
reg Q;
always @ (posedge clk) Q <= D;
```

- ▶ **@(posedge clk)** is the sensitivity list
- ▶ The **Q <= D**; is the block part
- ▶ The block part is always “entered” whenever the sensitivity list becomes true (positive edge of clk)
- ▶ The LHS of the **<=** must be of data type **reg**
- ▶ The RHS of the **<=** may use **reg** or **wire**

## Synthesis: Always Statement

- ▶ This is an asynchronous clear D Flip-Flop

```
reg Q;
always @ (posedge clk, posedge rst)
 if (rst) Q <= 'b0; else Q <= D;
```

- ▶ Notice **,** instead of **or**
- ▶ Verilog 2001...
- ▶ Positive reset

## Synthesis: Always Statement

```
reg Q;
always @ (posedge clk, posedge rst, posedge set)
 if (rst) Q <= 'b0;
 else if (set) Q <= 'b1;
 else Q <= D;
```

- ▶ What is this?
- ▶ What is synthesized?

> beh2str foo.v foo\_str.v UofU\_Digital.db

## Synthesis: Always Statement

```
reg Q;
always @ (posedge clk, posedge rst, posedge set)
 if (rst) Q <= 'b0;
 else if (set) Q <= 'b1;
 else Q <= D;
```

- ▶ What is this?
- ▶ What is synthesized?

~/home/atanner/1U\_UAW/synopsys/nt.v

| Register Name | Type      | Width | Bus | MB | AR | AS | SR | SS | ST |
|---------------|-----------|-------|-----|----|----|----|----|----|----|
| rout_reg      | Flip-flop | 1     | 1   | N  | N  | Y  | Y  | I  | I  |

Presto compilation completed successfully.

Current design is now '/home/atanner/IC\_CAD/synopsys/nt.db;nt'

## Synthesis: Always Statement

```
reg Q;
always @(posedge clk, posedge rst, posedge set)
 if (rst) Q <= 'b0;
 else if (set) Q <= 'b1;
 else Q <= D;
```

► What is this?

► What is synthesized?

```
module nt (clk, rst, set, a, b, c, d, rout);
 input clk, rst, set, a, b;
 output c, d, rout;
 wire n2, n4, n5;

 NAND2 U6 (.A(c), .B(b), .Y(d));
 //***FFGEN** rout_reg (.next_state(a), .clocked_on(clk), .force_00(n2),
 //force_01(rst), .force_10(n4), .force_11(n2), .Q(rout));
 TIELO U9 (.Y(n2));
 NOR2 U10 (.A(rst), .B(n5), .Y(n4));
 INV U11 (.A(set), .Y(n5));
 NAND2 U12 (.A(d), .B(a), .Y(c));
endmodule
```

## Synthesis: Always Statement

```
reg Q;
always @(posedge clk)
 if (rst) Q <= 'b0;
 else if (set) Q <= 'b1;
 else Q <= D;
```

► What is this?

## Synthesis: Always Statement

```
reg Q;
always @(posedge clk)
 if (rst) Q <= 'b0;
 else if (set) Q <= 'b1;
 else Q <= D;
```

► What is this?

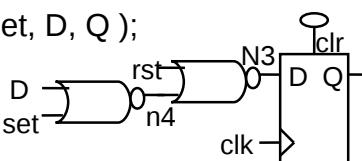
Inferred memory devices in process  
in routine set line 5 in file  
'/home/elb/IC\_CAD/syn-f06/set.v'.

| Register Name | Type      | Width | Bus | MB | AR | AS | SR | SS | ST |
|---------------|-----------|-------|-----|----|----|----|----|----|----|
| Q_reg         | Flip-flop | 1     | N   | N  | N  | N  | N  | N  | N  |

```
module foo (clk, rst, set, D, Q);
 input clk, rst, set, D;
 output Q;
 wire N3, n2, n4;
```

```
dff Q_reg (.D(N3), .G(clk), .CLR(n2), .Q(Q));
tiehi U6 (.Y(n2));
nor2 U7 (.A(rst), .B(n4), .Y(N3));
nor2 U8 (.A(D), .B(set), .Y(n4));
endmodule
```

```
module foo (clk, rst, set, D, Q);
 input clk, rst, set, D;
 output Q;
 wire N3, n2, n4;
```



```
dff Q_reg (.D(N3), .G(clk), .CLR(n2), .Q(Q));
tiehi U6 (.Y(n2));
nor2 U7 (.A(rst), .B(n4), .Y(N3));
nor2 U8 (.A(D), .B(set), .Y(n4));
endmodule
```

| A | B | Out |
|---|---|-----|
| 0 | 0 | 1   |
| 0 | 1 | 0   |
| 1 | 0 | 0   |
| 1 | 1 | 0   |

## Synthesis: Always Statement

```
reg P,Q;
reg [3:0] R;
always @(posedge clk)
begin
 Q <= D;
 P <= Q;
 R <= R + 'h1;
end
```

► What is this?

► Will it synthesize? Simulate?

## Synthesis: Always Statement

```
module testme (D, P, Q, R, clk);
output [3:0] R;
input D, clk;
output P, Q;
wire N0, N1, N2, N3, n1, n7, n8, n9;
dff Q_reg (.D(D), .G(clk), .CLR(n1), .Q(Q));
dff P_reg (.D(Q), .G(clk), .CLR(n1), .Q(P));
dff R_reg_0_ (.D(N0), .G(clk), .CLR(n1), .Q(R[0]));
dff R_reg_1_ (.D(N1), .G(clk), .CLR(n1), .Q(R[1]));
dff R_reg_2_ (.D(N2), .G(clk), .CLR(n1), .Q(R[2]));
dff R_reg_3_ (.D(N3), .G(clk), .CLR(n1), .Q(R[3]));
tiehi U9 (.Y(n1));
xor2 U10 (.A(R[3]), .B(n7), .Y(N3));
nor2 U11 (.A(n8), .B(n9), .Y(n7));
xor2 U12 (.A(n8), .B(n9), .Y(N2));
invX1 U13 (.A(R[2]), .Y(n9));
nand2 U14 (.A(R[1]), .B(R[0]), .Y(n8));
xor2 U15 (.A(R[1]), .B(R[0]), .Y(N1));
invX1 U16 (.A(R[0]), .Y(N0));
endmodule
```

## Synthesis: Always Statement

- This is a simple D Flip-Flop  
reg Q;  
always @(posedge clk) Q <= D;
- So is this  
reg Q;  
always @(posedge clk) Q = D;
- = is for blocking assignments
- <= is for nonblocking assignments

## Constants

- **parameter** used to define constants
  - parameter size = 16, foo = 8;
  - wire [size-1:0] bus; \\\ defines a 15:0 bus
  - externally modifiable
  - scope is local to module
- **localparam** not externally modifiable
  - localparam width = size \* foo;
- **`define** macro definition
  - `define value 7'd53
  - assign a = (sel == `value) & b;
  - scope is from here on out

## Example: Counter

```
module counter (clk, clr, load, in, count);
parameter width=8;
input clk, clr, load;
input [width-1 : 0] in;
output [width-1 : 0] count;
reg [width-1 : 0] tmp;

always @ (posedge clk or negedge clr)
begin
 if (!clr)
 tmp = 0;
 else if (load)
 tmp = in;
 else
 tmp = tmp + 1;
end
assign count = tmp;
endmodule
```

## Synthesis: Modules

```
module the_top (clk, rst, a, b, sel, result);
 input clk, rst;
 input [3:0] a,b; input [2:0] sel;
 output reg [3:0] result;
 wire[3:0] sum, dif, alu;

 adder u0(a,b,sum);
 subber u1(.subtrahend(a), .subtractor(b), .difference(dif));

 assign alu = {4{ (sel == 'b000)}} & sum
 | {4{ (sel == 'b001)}} & dif;

 always @ (posedge clk or posedge rst)
 if(rst) result <= 'h0;
 else result <= alu;

endmodule
```

## Synthesis: Modules

// Verilog 1995 syntax

```
module adder (e,f,g);
 parameter SIZE=2;
 input [SIZE-1:0] e, f;
 output [SIZE-1:0] g;
 assign g = e + f;
endmodule
```

// Verilog 2001 syntax

```
module subber #(parameter SIZE = 3)
 (input [SIZE-1:0] c,d, output [SIZE-1:0] difference);
 assign difference = c - d;
endmodule
```

## Synthesis: Modules

```
module the_top (clk, rst, a, b, sel, result);
parameter SIZE = 4;
input clk, rst;
input [SIZE-1:0] a,b;
input [2:0] sel;
output reg [SIZE-1:0] result;
wire[SIZE-1:0] sum, dif, alu;

adder #(SIZE(SIZE)) u0(a,b,sum);
subber #(4) u1(.c(a), .d(b), .difference(dif));

assign alu = {SIZE{sel == 'b000}} & sum
 | {SIZE{sel == 'b001}} & dif;

always @(posedge clk or posedge rst)
 if (rst) result <= 'h0;
 else result <= alu;
endmodule
```

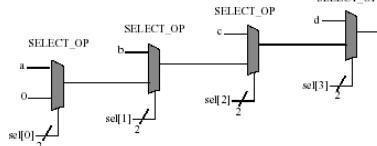
## Multi-Way Decisions

► Standard if-else-if syntax

```
If (<expression>)
 <statement>
else if (<expression>)
 <statement>
else if (<expression>)
 <statement>
else <statement>
```

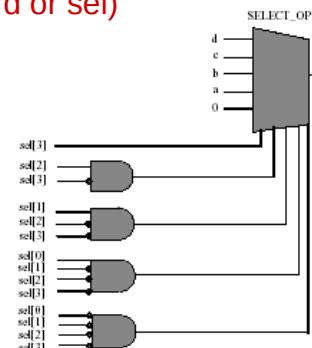
## Priority vs. Parallel Choice (if)

```
module priority (a, b, c, d, sel, z);
input a,b,c,d;
input [3:0] sel;
output z;
reg z;
always @ (a or b or c or d or sel)
begin
 z = 0;
 if (sel[0]) z = a;
 if (sel[1]) z = b;
 if (sel[2]) z = c;
 if (sel[3]) z = d;
end
endmodule
```



## Priority vs. Parallel Choice

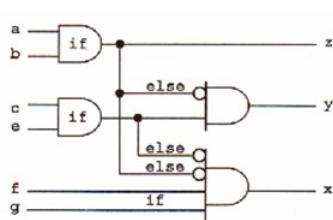
```
module parallel (a, b, c, d, sel, z);
input a,b,c,d;
input [3:0] sel;
output z;
reg z;
always @ (a or b or c or d or sel)
begin
 z = 0;
 if (sel[3]) z = d;
 else if (sel[2]) z = c;
 else if (sel[1]) z = b;
 else if (sel[0]) z = a;
end
endmodule
```



## Priority Encoders

```
/*
// Priority encoders
// Allen Tanner
*/

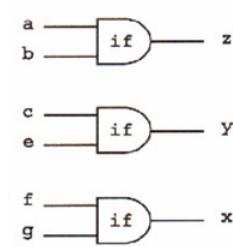
module prior_enc(x,y,z, a,b,c,d,e,f);
 output reg x,y,z;
 input a,b,c,d,e,f;
 always@ (a,b,c,d,e,f,g)
 begin
 {x,y,z} = 3'b0;
 if ((a==1) && (b==1)) z = 1;
 else if ((c==1) && (d==1)) y = 1;
 else if ((e==1) && (f==1)) x = 1;
 end
endmodule // prior_enc
```



## Priority Encoders

```
/*
// Priority encoders
// Allen Tanner
*/

module prior_enc(x,y,z, a,b,c,d,e,f);
 output reg x,y,z;
 input a,b,c,d,e,f;
 always@ (a,b,c,d,e,f,g)
 begin
 {x,y,z} = 3'b0;
 if ((a==1) && (b==1)) z = 1;
 if ((c==1) && (d==1)) y = 1;
 if ((e==1) && (f==1)) x = 1;
 end
endmodule // prior_enc
```



## Case Statements

- ▶ Multi-way decision on a single expression

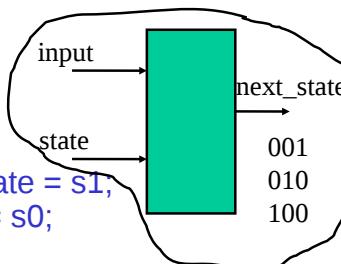
```
case (<expression>)
 <expression>: <statement>
 <expression>, <expression>: <statement>
 <expression>: <statement>
 default: <statement>
endcase
```

## Case Example

```
reg [1:0] sel;
reg [15:0] in0, in1, in2, in3, out;
case (sel)
 2'b00: out = in0;
 2'b01: out = in1;
 2'b10: out = in2;
 2'b11: out = in3;
endcase
```

## Another Case Example

```
// simple counter next-state logic
// one-hot state encoding...
parameter [2:0] s0=3'h1, s1=3'h2, s2=3'h4;
reg[2:0] state, next_state;
always @(input or state)
begin
 case (state)
 s0: if (input) next_state = s1;
 else next_state = s0;
 s1: next_state = s2;
 s2: next_state = s0;
 endcase
end
```



## Weird Case Example

- ▶ Verilog allows you to put a value in the case slot, and test which variable currently has that value...

```
reg [2:0] curr_state, next_state;
parameter s1=3'b001, s2=3'b010, s3=3'b100
case (1)
 curr_state[0] : next_state = s2;
 curr_state[1] : next_state = s3;
 curr_state[2] : next_state = s1;
endcase
```

## Latch Inference

- ▶ Incompletely specified `if` and `case` statements cause the synthesizer to infer latches

```
always @(cond)
begin
 if (cond) data_out <= data_in;
end
```

- ▶ This infers a latch because it doesn't specify what to do when cond = 0
  - ▶ Fix by adding an `else` if you want combinational logic
- ▶ In a case, fix by including `default`:

## Full vs. Parallel

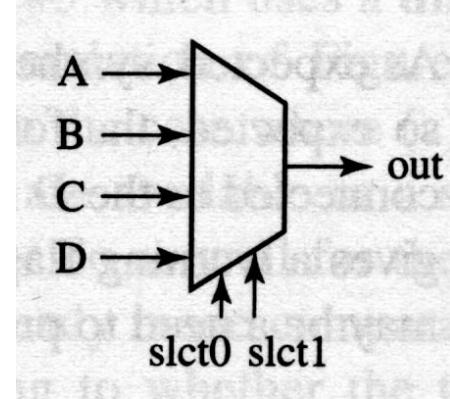
- ▶ `Case` statements check each case in sequence
- ▶ A `case` statement is **full** if all possible outcomes are accounted for
- ▶ A `case` statement is **parallel** if the stated alternatives are mutually exclusive
- ▶ These distinctions make a difference in how `cases` are translated to circuits...
  - ▶ Similar to the `if` statements previously described

## Case full-par example

```
// full and parallel = combinational logic
module full-par (slct, a, b, c, d, out);
 input [1:0] slct;
 input a, b, c, d;
 output out;
 reg out; // optimized away in this example
 always @(slct or a or b or c or d)
 case (slct)
 2'b11 : out <= a;
 2'b10 : out <= b;
 2'b01 : out <= c;
 default : out <= d; // really 2'b10
 endcase
 endmodule
```

## Synthesis Result

- ▶ Note that full-par results in combinational logic

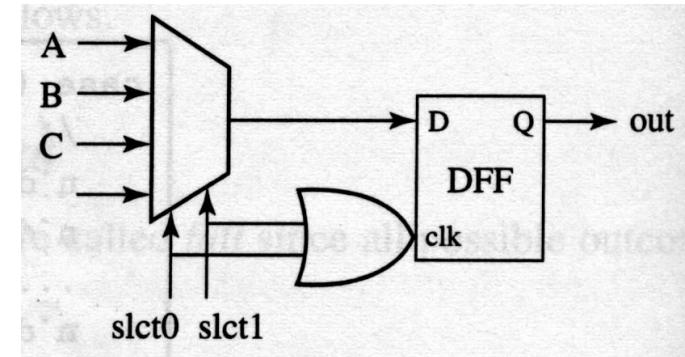


## Case notfull-par example

```
// a latch is synthesized because case is not full
module notfull-par (slct, a, b, c, d, out);
 input [1:0] slct;
 input a, b, c, d;
 output out;
 reg out; // NOT optimized away in this example
 always @(slct or a or b or c)
 case (slct)
 2'b11 : out <= a;
 2'b10 : out <= b;
 2'b01 : out <= c;
 endcase
 endmodule
```

## Synthesized Circuit

- ▶ Because it's not full, a latch is inferred...

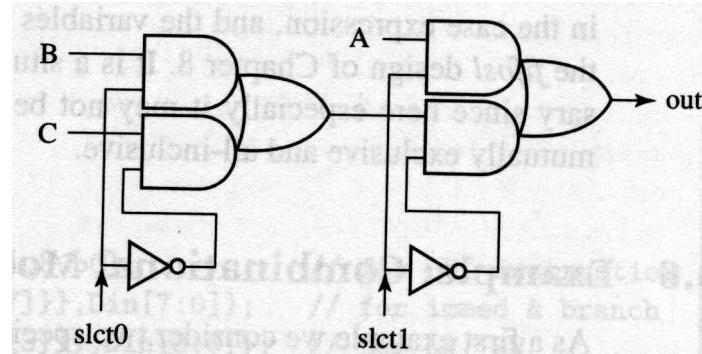


## Case full-notpar example

```
// because case is not parallel - priority encoding
// but it is still full, so no latch...
// this uses a casez which treats ? as don't-care
module full-notpar (slct, a, b, c, out);
 ...
 always @(slct or a or b or c)
 casez (slct)
 2'b1? : out <= a;
 2'b?1 : out <= b;
 default : out <= c;
 endcase
 endmodule
```

## Synthesized Circuit

- ▶ It's **full**, so it's combinational, but it's **not parallel** so it's a priority circuit instead of a "check all in parallel" circuit



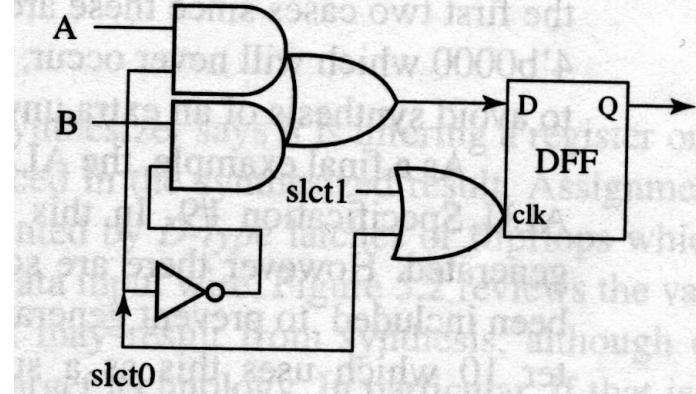
## Case notfull-notpar example

```
// because case is not parallel - priority encoding
// because case is not full - latch is inferred
// uses a casez which treats ? as don't-care
module full-notpar (slct, a, b, c, out);

...
always @(slct or a or b or c)
 casez (slct)
 2'b1? : out <= a;
 2'b?1 : out <= b;
 endcase
endmodule
```

## Synthesized Circuit

- Not full and not parallel, infer a latch



## Get off my Case

### Verification

- CASE matches all (works like ==)
- CASEX uses "z", "x", "?" as don't care
- CASEZ uses "z", "?" as don't care
- Beware: Matches **first** valid case

### Synthesis

- CASE works like ==
- CASEX uses "?" as don't care
- CASEZ uses "?" as don't care

## Get off my Case

```
// Case tests
// Allen Tanner
reg [15:0] opcode;
initial begin
 $display("We've only just begun");
 #1 $display ("Driving add");
 opcode = 16'b1000_1000_1010_1111;
 #1 $display ("Driving subtract");
 opcode = 16'b0100_0100_1010_1111;
 #1 $display ("Driving multiply");
 opcode = 16'b0010_0010_1010_1111;
 #1 $finish;
end

always @ (opcode)
casex (opcode)
 16'b1xxz?_xxxx_xxxx_xxxx : $display("CASEX. Opcode: add");
 16'bxx1x_xxxx_xxxx_xxxx : $display("CASEX. Opcode: subtract");
 16'bxx1x_xxxx_xxxx_xxxx : $display("CASEX. Opcode: multiply");
endcase

always @ (opcode)
casex (opcode)
 16'b1?1?_zzzz_zzzz_zzzz : $display("CASEZ. Opcode: add");
 16'b1?1?_zzzz_zzzz_zzzz : $display("CASEZ. Opcode: subtract");
 16'b1?1?_zzzz_zzzz_zzzz : $display("CASEZ. Opcode: multiply");
endcase
```

We've only just begun  
 Driving add  
**CASEZ.** Opcode: add  
**CASEX.** Opcode: add  
 Driving subtract  
**CASEX.** Opcode: subtract  
**CASEZ.** Opcode: subtract  
 Driving multiply  
**CASEZ.** Opcode: multiply  
**CASEX.** Opcode: multiply  
L18 "testfixture.new": \$fin

## Get off my Case

```
// Case tests
// Allen Tanner
reg sel;
initial begin
 $display("We've only just begun");
 #1 $display ("Driving 0");
 sel = 0;
 #1 $display ("Driving 1");
 sel = 1;
 #1 $display ("Driving x");
 sel = 1'bxx;
 #1 $display ("Driving z");
 sel = 1'bzz;
 #1 $finish;
end

always @ (sel)
case (sel)
 1'b0 : $display("CASE : Logic 0 on sel");
 1'b1 : $display("CASE : Logic 1 on sel");
 1'b? : $display("CASE : Logic x on sel");
 1'bxx : $display("CASE : Logic z on sel");
endcase

always @ (sel)
casex (sel)
 1'b0 : $display("CASEX : Logic 0 on sel");
 1'b1 : $display("CASEX : Logic 1 on sel");
 1'b? : $display("CASEX : Logic x on sel");
 1'bxx : $display("CASEX : Logic z on sel");
endcase

always @ (sel)
casesel (sel)
 1'b0 : $display("CASEZ : Logic 0 on sel");
 1'b1 : $display("CASEZ : Logic 1 on sel");
 1'b? : $display("CASEZ : Logic x on sel");
 1'bxx : $display("CASEZ : Logic z on sel");
endcase

always @ (sel)
casex (sel)
 1'b0 : $display("CASEZA : Logic 0 on sel");
 1'b1 : $display("CASEZA : Logic 1 on sel");
 1'b? : $display("CASEZA : Logic x on sel");
 1'bxx : $display("CASEZA : Logic z on sel");
endcase
```

We've only just begun  
 Driving 0  
**CASEZa** : Logic 0 on sel  
**CASEX** : Logic 0 on sel  
**CASEZ** : Logic 0 on sel  
**CASE** : Logic 0 on sel

Driving 1  
**CASE** : Logic 1 on sel  
**CASEX** : Logic 1 on sel  
**CASEZ** : Logic 1 on sel  
**CASEza** : Logic 1 on sel

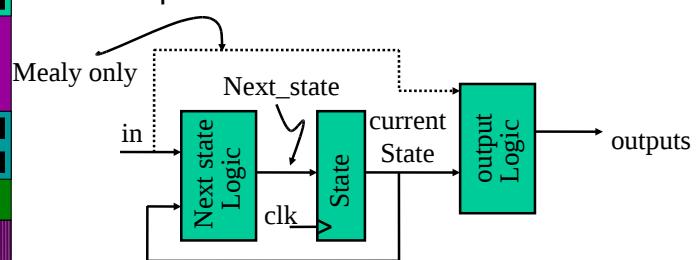
Driving x  
**CASEZa** : Logic x on sel  
**CASEX** : Logic 0 on sel  
**CASEZ** : Logic 0 on sel  
**CASE** : Logic x on sel

Driving z  
**CASE** : Logic z on sel  
**CASEX** : Logic 0 on sel  
**CASEZ** : Logic 0 on sel  
**CASEza** : Logic 1 on sel  
L18 "testfixture.new": \$fin

**Order Matters**

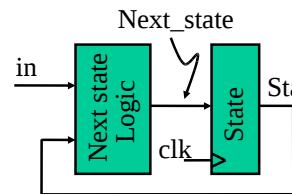
## FSM Description

- One simple way: break it up like a schematic
- A combinational block for next\_state generation
- A combinational block for output generation
- A sequential block to store the current state

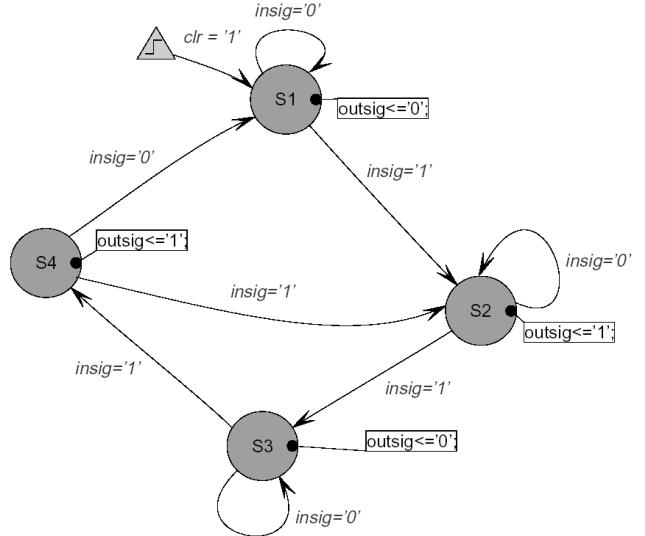


# Modeling State Machines

```
// General view
module FSM (clk, in, out);
 input clk, in;
 output out;
 reg out;
 // state variables
 reg [1:0] state;
 // next state variable
 reg [1:0] next_state;
 always @posedge(clk) // state register
 state = next_state;
 always @(state or in); // next-state logic
 // compute next state and output logic
 // make sure every local variable has an
 // assignment in this block
endmodule
```



# FSM Description



## Verilog Version

```
module moore (clk, clr, insig, outsig); // define combinational logic for
 input clk, clr, insig;
 output outsig;
 // define state encodings as
 parameters
 parameter [1:0] s0 = 2'b00,
 s1 = 2'b01, s2 = 2'b10, s3 = 2'b11;
 // define reg vars for state register
 // and next_state logic
 reg [1:0] state, next_state;
 //define state register (with
 //synchronous active-high clear)
 always @(posedge clk)
 begin
 if (clr) state = s0;
 else state = next_state;
 end
 // assign outsig as continuous assign
 assign outsig =
 ((state == s1) || (state == s3));
endmodule
```

## Verilog Version

```
module moore (clk, clr, insig, outsig);
 input clk, clr, insig;
 output outsig;
 // define state encodings as parameters
 parameter [1:0] s0 = 2'b00, s1 = 2'b01,
 s2 = 2'b10, s3 = 2'b11;
 // define reg vars for state register and next_state logic
 reg [1:0] state, next_state;
 //define state register (with synchronous active-high clear)
 always @(posedge clk)
 begin
 if (clr) state = s0;
 else state = next_state;
 end
```

## Verilog Version Continued...

```
// define combinational logic for next_state
always @(insig or state)
begin
 case (state)
 s0: if (insig) next_state = s1;
 else next_state = s0;
 s1: if (insig) next_state = s2;
 else next_state = s1;
 s2: if (insig) next_state = s3;
 else next_state = s2;
 s3: if (insig) next_state = s1;
 else next_state = s0;
 endcase
end
```

## Verilog Version Continued...

```
// now set the outsig. This could also be done in an always
// block... but in that case, outsig would have to be
// defined as a reg.
assign outsig = ((state == s1) || (state == s3));
endmodule
```

## Unsupported for Synthesis

- ▶ Delay (Synopsys will ignore #'s)
- ▶ initial blocks (use explicit resets)
- ▶ repeat
- ▶ wait
- ▶ fork
- ▶ event
- ▶ deassign
- ▶ force
- ▶ release

## More Unsupported Stuff

- ▶ You cannot assign the same reg variable in more than one procedural block

// don't do this...

```
always @(posedge a)
 out = in1;
always @(posedge b)
 out = in2;
```

## Combinational Always Blocks

- ▶ Be careful...

```
always @(sel)
 if (sel == 1)
 out = in1;
 else
 out = in2;
always @(sel or in1 or in2)
 if (sel == 1)
 out = in1;
 else
 out = in2;
```

- ▶ Which one is a good mux?

## Combinational Always Blocks

- ▶ Be careful...

```
always @(sel)
 if (sel == 1)
 out = in1;
 else
 out = in2;
always @(sel or in1 or in2)
 if (sel == 1)
 out = in1;
 else
 out = in2;
```

- ▶ Which one is a good mux?

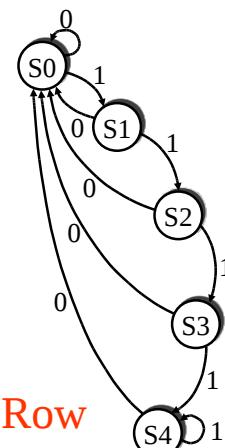
▶ Always @\*
 if (sel == 1) out = in1; else out = in2;

## Sync vs. Async Register Reset

```
// synchronous reset (active-high reset)
always @(posedge clk)
 if (reset) state = s0;
 else state = s1;
```

```
// async reset (active-low reset)
always @(posedge clk or negedge reset)
 if (reset == 0) state = s0;
 else state = s1;
```

## Finite State Machine



Four in a Row

## Textbook FSM

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner
module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 // define state encodings as parameters
 parameter [2:0] s0 = 3'b000, s1 = 3'b001, s2 = 3'b010, s3 = 3'b011, s4 = 3'b100;
 // define reg vars for state register and next_state logic
 reg [2:0] state, next_state;
 //define state register (with asynchronous active-low clear)
 always @ (posedge clk or negedge clr)
 begin
 if (clr==0) state = s0;
 else state = next_state;
 end

 // define combinational logic for next_state
 always @ (insig or state)
 begin
 case (state)
 s0: if (insig) next_state = s1;
 else next_state = s0;
 s1: if (insig) next_state = s2;
 else next_state = s0;
 s2: if (insig) next_state = s3;
 else next_state = s0;
 s3: if (insig) next_state = s4;
 else next_state = s0;
 s4: if (insig) next_state = s4;
 else next_state = s0;
 default: next_state = s0;
 endcase
 end

 // now set the saw4. This could also be done in an always
 // block... but in that case, saw4 would have to be
 // defined as a reg.
 assign saw4 = state == s4;
endmodule
```

## Textbook FSM

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner
module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 // define state encodings as parameters
 parameter [2:0] s0 = 3'b000, s1 = 3'b001, s2 = 3'b010, s3 = 3'b011, s4 = 3'b100;
 // define reg vars for state register and next_state logic
 reg [2:0] state, next_state;
 //define state register (with asynchronous active-low clear)
 always @ (posedge clk or negedge clr)
 begin
 if (clr==0) state = s0;
 else state = next_state;
 end

 // define combinational logic for next_state
 always @ (insig or state)
 begin
 case (state)
 s0: if (insig) next_state = s1;
 else next_state = s0;
 s1: if (insig) next_state = s2;
 else next_state = s0;
 s2: if (insig) next_state = s3;
 else next_state = s0;
 s3: if (insig) next_state = s4;
 else next_state = s0;
 s4: if (insig) next_state = s4;
 else next_state = s0;
 default: next_state = s0;
 endcase
 end

 // now set the saw4. This could also be done in an always
 // block... but in that case, saw4 would have to be
 // defined as a reg.
 assign saw4 = state == s4;
endmodule
```

**Comments**

**Polarity?**

**Always use <= for FF?**

## Documented FSM

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner
module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 parameter [2:0] s0 = 3'b000; // initial state, saw at least 1 zero
 parameter [2:0] s1 = 3'b001; // saw 1 one
 parameter [2:0] s2 = 3'b010; // saw 2 ones
 parameter [2:0] s3 = 3'b011; // saw 3 ones
 parameter [2:0] s4 = 3'b100; // saw at least, 4 ones

 reg [2:0] state, next_state;
 // state register
 always @ (posedge clk or posedge clr)
 begin
 if (clr) state <= s0;
 else state <= next_state;
 end

 // next state logic
 always @ (insig or state)
 begin
 case (state)
 s0: if (insig) next_state = s1;
 else next_state = s0;
 s1: if (insig) next_state = s2;
 else next_state = s0;
 s2: if (insig) next_state = s3;
 else next_state = s0;
 s3: if (insig) next_state = s4;
 else next_state = s0;
 s4: if (insig) next_state = s4;
 else next_state = s0;
 default: next_state = s0;
 endcase
 end

 assign saw4 = state == s4;
endmodule
```

## Waveform Test Bench

```
// Four ones in a row detector.
// Test bench
// Allen Tanner
initial
begin
 clk = 1'b0;
 clr = 1'b0;
 insig = 1'b0;

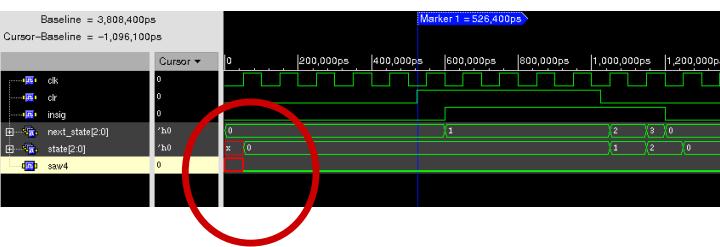
 send_message(32'b0001_1000_1010_1111_0000_0111_1110_0000);
 send_message(32'b0001_1000_1010_1111_0000_0111_1110_0000);
 $finish;
end

always #50 clk = ~clk;

initial
begin
 #525
 clr = 1'b1;
 #500
 clr = 1'b0;
end

task send_message;
 input [31:0] pattern;
 integer i;
 begin
 for(i=0;i<32; i=i+1)
 @(negedge clk) insig = pattern[i];
 end
endtask // send_message
```

## Waveform



Pay attention to first few cycles...

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 input clk, clr, insig;
 output saw4;

 parameter [2:0] s0 = 3'b000; // initial state, saw at least 1 zero
 parameter [2:0] s1 = 3'b001; // saw 1 one
 parameter [2:0] s2 = 3'b010; // saw 2 ones
 parameter [2:0] s3 = 3'b011; // saw 3 ones
 parameter [2:0] s4 = 3'b100; // saw at least, 4 ones

 reg [2:0] state, next_state;
 // state register
 always @ (posedge clk or posedge clr)
 begin
 if (clr) state <= s0;
 else
 case (state)
 s0: if (insig) state <= s1;
 else state <= s0;
 s1: if (insig) state <= s2;
 else state <= s0;
 s2: if (insig) state <= s3;
 else state <= s0;
 s3: if (insig) state <= s4;
 else state <= s0;
 s4: if (insig) state <= s4;
 else state <= s0;
 default: state <= s0;
 endcase
 end

 assign saw4 = state == s4;
endmodule
```

## FSM

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 parameter [2:0] s3 = 3'b011; // saw 3 ones
 parameter [2:0] s4 = 3'b100; // saw at least, 4 ones

 reg [2:0] state;
 wire [2:0] next_state;

 assign next_state = {3{((state == s0) & !insig)} & s0
 | 3{((state == s0) & insig)} & s1
 | 3{((state == s1) & !insig)} & s0
 | 3{((state == s1) & insig)} & s2
 | 3{((state == s2) & !insig)} & s0
 | 3{((state == s2) & insig)} & s3
 | 3{((state == s3) & !insig)} & s0
 | 3{((state == s3) & insig)} & s4
 | 3{((state == s4) & !insig)} & s0
 | 3{((state == s4) & insig)} & s4};

 always @ (posedge clk or posedge clr) // state register
 begin
 if (clr) state <= s0;
 else state <= next_state;
 end

 assign saw4 = state == s4;
endmodule
```

## One-Hot FSM

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner

module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 reg s0; // initial state, saw at least 1 zero
 reg s1; // saw 1 one
 reg s2; // saw 2 ones
 reg s3; // saw 3 ones
 reg s4; // saw at least, 4 ones

 always @ (posedge clk or posedge clr) // state register
 begin
 if (clr)
 begin
 s0 <= 1'b1;
 s1 <= 1'b0;
 s2 <= 1'b0;
 s3 <= 1'b0;
 s4 <= 1'b0;
 end
 else
 begin
 s0 <= (s0 | s1 | s2 | s3 | s4) & !insig;
 s1 <= s0 & insig;
 s2 <= s1 & insig;
 s3 <= s2 & insig;
 s4 <= s3 & insig
 | s4 & insig;
 end
 end

 assign saw4 = s4;
endmodule
```

## Oops

```
module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;
 wire s0, s1, s2, N2, N3, N4, N5, n9, n10, n11, n14, n15, n16, n17, n18,
 n19, n20, n21, n22;

 /*^FFGEN^*/ s0_reg (.next_state(n14), .clocked_on(clk), .force_00(n11),
 .force_01(n11), .force_10(clr), .force_11(n11), .Q(s0));
 DFF s2_reg (.D(N3), .CLK(clk), .nCLR(n15), .Q(s2));
 DFF s1_reg (.D(N2), .CLK(clk), .nCLR(n15), .Q(s1));
 DFF s4_reg (.D(N5), .CLK(clk), .nCLR(n15), .Q(saw4), .QB(n9));
 DFF s3_reg (.D(N4), .CLK(clk), .nCLR(n15), .Q(B(n10)));
 TIEQ U20 (.Y(n11));
 INV U21 (.A(clr), .Y(n15));
 AOI U22 (.A(n16), .B(n17), .C(insig), .Y(n14));
 NOR2 U23 (.A(s2), .B(s1), .Y(n17));
 NOR2 U24 (.A(s0), .B(n18), .Y(n16));
 INV U25 (.A(n19), .Y(N5));
 NAND2 U26 (.A(n18), .B(insig), .Y(n19));
 NAND2 U27 (.A(n9), .B(n10), .Y(n18));
 INV U28 (.A(n20), .Y(N4));
 NAND2 U29 (.A(insig), .B(s2), .Y(n20));
 INV U30 (.A(n21), .Y(N3));
 NAND2 U31 (.A(insig), .B(s1), .Y(n21));
 INV U32 (.A(n22), .Y(N2));
 NAND2 U33 (.A(insig), .B(s0), .Y(n22));
endmodule
```

## FSM

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner

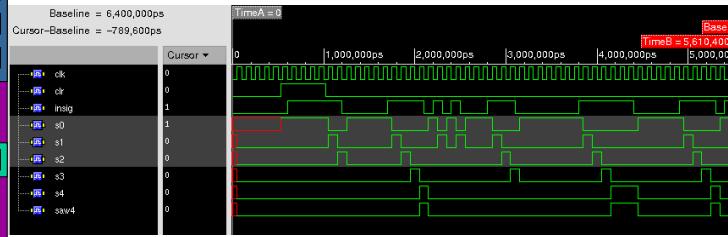
module seed4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 parameter [2:0] s0 = 3'b000; // initial state, saw at least 1 zero
 parameter [2:0] s1 = 3'b001; // saw 1 one
 parameter [2:0] s2 = 3'b010; // saw 2 ones
 parameter [2:0] s3 = 3'b011; // saw 3 ones
 parameter [2:0] s4 = 3'b100; // saw at least, 4 ones

 reg [2:0] state;
 always @ (posedge clk or posedge clr) // state register
 begin
 if (clr) state <= s0;
 else state <= {3{((state == s0) & !insig)} & s0
 | 3{((state == s0) & insig)} & s1
 | 3{((state == s1) & !insig)} & s0
 | 3{((state == s1) & insig)} & s2
 | 3{((state == s2) & !insig)} & s0
 | 3{((state == s2) & insig)} & s3
 | 3{((state == s3) & !insig)} & s0
 | 3{((state == s3) & insig)} & s4
 | 3{((state == s4) & !insig)} & s0
 | 3{((state == s4) & insig)} & s4};
 end

 assign saw4 = state == s4;
endmodule
```

## One-Hot FSM Counting



## No Asynchronous Sets

```
// Verilog HDL for "Ax", "see4" "behavioral"
// Four in a row detector - Allen Tanner

module seed4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 reg ns0; // initial state, saw at least 1 zero
 reg s1; // saw 1 one
 reg s2; // saw 2 ones
 reg s3; // saw 3 ones
 reg s4; // saw at least, 4 ones

 wire s0; // alias for !ns0 (ns0 used to avoid FFGEN in beh2str)
 assign s0 = !ns0;

 always @ (posedge clk or posedge clr) // state register
 begin
 if (clr)
 begin
 ns0 <= 1'b0;
 s1 <= 1'b0;
 s2 <= 1'b0;
 s3 <= 1'b0;
 s4 <= 1'b0;
 end
 else
 begin
 ns0 <= ~((s0 | s1 | s2 | s3 | s4) & !insig);
 s1 <= s0 & insig;
 s2 <= s1 & insig;
 s3 <= s2 & insig;
 s4 <= s3 & insig
 | s4 & insig;
 end
 end

 assign saw4 = s4;
```

# That's better

```

module seq4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;
 wire ns0, N0, N1, N2, N4, N5, n1, n9, n10, n11, n12, n13, n14, n15;

 DFF s4_reg (.D(N5), .CLK(clk), .nCLR(n1), .Q(saw4), .QB(n9));
 DFF s3_reg (.D(N2), .CLK(clk), .nCLR(n1), .QB(n12));
 DFF s2_reg (.D(N1), .CLK(clk), .nCLR(n1), .QB(n10));
 DFF s1_reg (.D(N0), .CLK(clk), .nCLR(n1), .QB(n11));
 DFF ns0_reg (.D(N4), .CLK(clk), .nCLR(n1), .Q(ns0));
 INV U12 (.A(clr), .Y(n1));
 AOI U13 (.A(n9), .B(n12), .C(n13), .Y(N5));
 OAI U14 (.A(n14), .B(n15), .C(n13), .Y(N4));
 NAND2 U15 (.A(ns0), .B(n12), .Y(n15));
 NAND3 U16 (.A(n10), .B(n9), .C(n11), .Y(n14));
 NOR2 U17 (.A(n13), .B(n10), .Y(N2));
 NOR2 U18 (.A(n13), .B(n11), .Y(N1));
 NOR2 U19 (.A(ns0), .B(n13), .Y(N0));
 INV U20 (.A(insig), .Y(n13));
endmodule

```

## Synchronous Clear

```

// Verilog HDL for "Ax", "see4" "behavioral"
// Synchronous clear - Allen Tanner

module seed4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;

 reg s0; // initial state, saw at least 1 zero
 reg s1; // saw 1 one
 reg s2; // saw 2 ones
 reg s3; // saw 3 ones
 reg s4; // saw at least, 4 ones

 always @ (posedge clk) // state register with synchronous clear
 begin
 if (clr)
 begin
 s0 <= 1'b1;
 s1 <= 1'b0;
 s2 <= 1'b0;
 s3 <= 1'b0;
 s4 <= 1'b0;
 end
 else
 begin
 s0 <= (s0 | s1 | s2 | s3 | s4) & !insig;
 s1 <= s0 & insig;
 s2 <= s1 & insig;
 s3 <= s2 & insig;
 s4 <= s3 & insig
 | s4 & insig;
 end
 end

 assign saw4 = s4;
endmodule

```

# Synchronous Clear

```

module see4 (clk, clr, insig, saw4);
 input clk, clr, insig;
 output saw4;
 wire N9, N10, N11, N12, N13, n2, n18, n19, n20, n21, n22, n23, n24, n25,
 n26, n27, n28, net12, net11, net9, net8;

 DFF s4_reg (.D(N13), .CLK(clk), .nCLR(n2), .Q(saw4), .QB(net8));
 DFF s3_reg (.D(N12), .CLK(clk), .nCLR(n2), .QB(net9));
 DFF s2_reg (.D(N11), .CLK(clk), .nCLR(n2), .Q(net11), .QB(n19));
 DFF s1_reg (.D(N10), .CLK(clk), .nCLR(n2), .QB(n18));
 DFF s0_reg (.D(N9), .CLK(clk), .nCLR(n2), .Q(net12), .QB(n20));
 TIEHI U21 (.Y(n2));
 NAND2 U22 (.A(n21), .B(n22), .Y(N9));
 NAND2 U23 (.A(n23), .B(n24), .Y(n22));
 INV U24 (.A(insig), .Y(n24));
 NAND3 U25 (.A(n25), .B(n18), .C(n26), .Y(n23));
 NOR2 U26 (.A(net12), .B(net11), .Y(n26));
 NOR2 U27 (.A(n25), .B(n27), .Y(N13));
 INV U28 (.A(n28), .Y(n25));
 NAND2 U29 (.A(net8), .B(net9), .Y(n28));
 NOR2 U30 (.A(n27), .B(n19), .Y(N12));
 NOR2 U31 (.A(n27), .B(n18), .Y(N11));
 NOR2 U32 (.A(n27), .B(n20), .Y(N10));
 NAND2 U33 (.A(insig), .B(n21), .Y(n27));
 INV U34 (.A(clr), .Y(n21));

endmodule

```

# Synchronous Clear



- ▶ Is asynchronous clear really asynchronous?
  - ▶ What about set-up & hold with respect to clock edge?

## ROM vs. Verilog

```

module mywords(addr, char);
 input [4:0] addr;
 output reg [5:0] char;

 always @ (addr)
 begin
 case(addr)
 'h00 : char = 'h00 ;
 'h01 : char = 'h24 ;
 'h02 : char = 'h28 ;
 'h03 : char = 'h25 ;
 'h04 : char = 'h2C ;
 'h05 : char = 'h2C ;
 'h06 : char = 'h2F ;
 'h07 : char = 'h00 ;
 'h08 : char = 'h37 ;
 'h09 : char = 'h2F ;
 'h0A : char = 'h32 ;
 'h0B : char = 'h2C ;
 'h0C : char = 'h24 ;
 'h0D : char = 'h00 ;
 'h0E : char = 'h00 ;
 'h0F : char = 'h00 ;

 'h10 : char = 'h00 ;
 'h11 : char = 'h00 ;
 'h12 : char = 'h33 ;
 'h13 : char = 'h25 ;
 'h14 : char = 'h34 ;
 'h15 : char = 'h00 ;
 'h16 : char = 'h2D ;
 'h17 : char = 'h25 ;
 'h18 : char = 'h00 ;
 'h19 : char = 'h36 ;
 'h1A : char = 'h25 ;
 'h1B : char = 'h25 ;
 'h1C : char = 'h25 ;
 'h1D : char = 'h00 ;
 'h1E : char = 'h00 ;
 'h1F : char = 'h00 ;
 endcase
 end
endmodule // mywords

```

## ROM vs. Verilog

