

Direct Construction of ROBDDs

- In the past lectures, we learnt:
 - $ITTE(f, g, h) = f \cdot g + f' \cdot h$
 - $f \odot g = x'(f^{x'} \odot g^{x'}) + x(f^x \odot g^x)$
 - ITE at top-nodes of $f, g, h \rightarrow$ ITE at their cofactors
 - Operate on graphs of f, g, h and derive the graph for $Z = ite(f, g, h)$
 - Shannon's expansion always w.r.t. top-node of f, g, h
- Use of a symbol table as a unique table
 - Every time you create a new node, Reduce it
 - Then, compute its *Key*: $\{low(v), v, high(v)\}$
 - No duplicate keys to be stored in the hash-table

The ITE Algorithm

```
ITE( $f, g, h$ ) {  
  if (terminal case)  
    return trivial result;  
  else  
     $x$  = top variable of  $f, g, h$ ;  
     $e = ITE(f_{x'}, g_{x'}, h_{x'})$ ;  
     $t = ITE(f_x, g_x, h_x)$ ;  
    if ( $t == e$ ) /* redundant node */  
      return  $t$ ; /* or return  $e$  */  
    /* Look-up the unique table for isomorphic subtrees */  
     $r = find-or-add-in-unique-table(e, x, t)$ ;  
    Update unique table, if required;  
    return  $r$ ;  
  end if  
}
```

ROBDD Construction Example

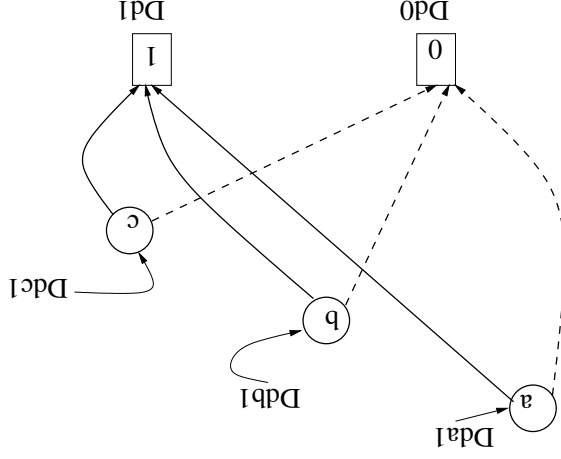
- $f = a + c; g = b + c; f \cdot g = ab + c$

- First construct trivial ROBDDs for a, b, c

- Assume var order a, b, c

- Fill-up the Unique Table (symbol/hash table)

Value	Key
Dd0	{NULL, 0, NULL}
Dd1	{NULL, 1, NULL}
Dda1	{Dd0, a, Dd1}
Ddb1	{Dd0, b, Dd1}
Ddc1	{Dd0, c, Dd1}

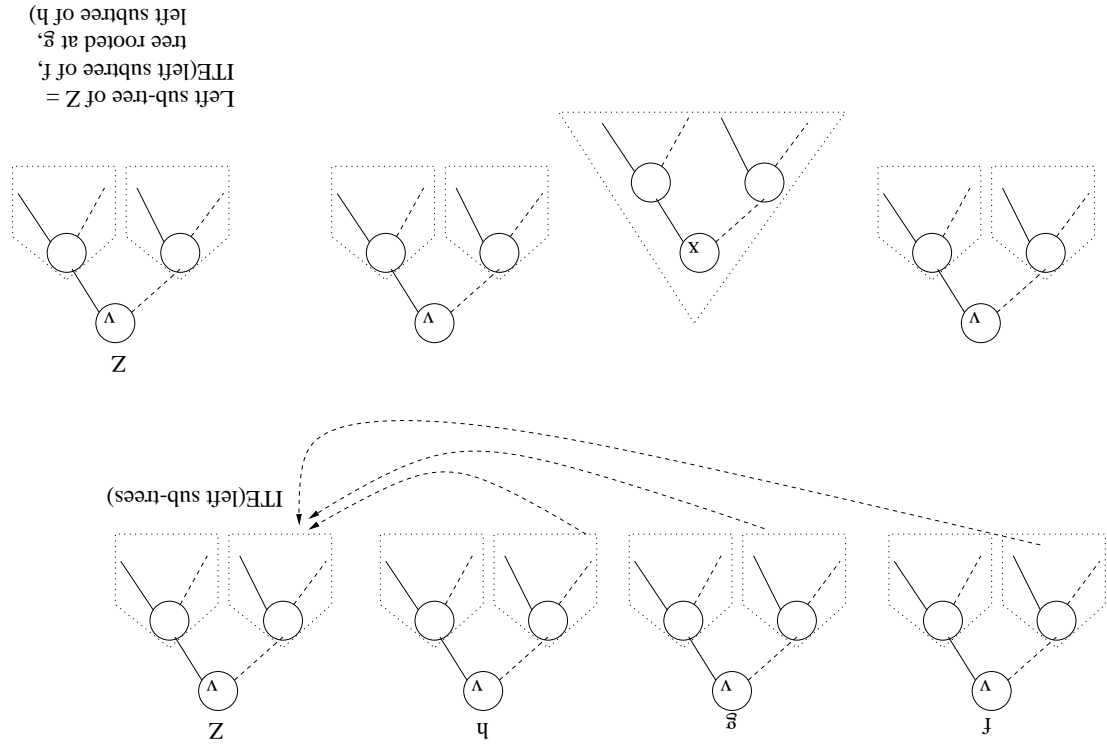


Constructing ROBDDs using ITE

- Few things to keep in mind...

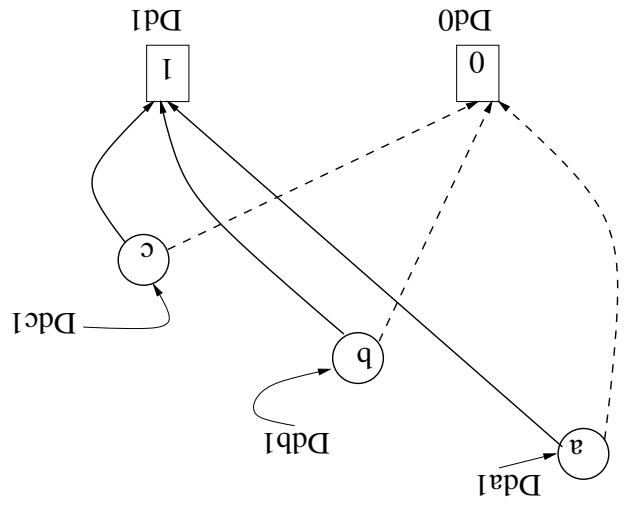
- When f, g, h have same top variable...

- When the f, g, h may not have same top-vars



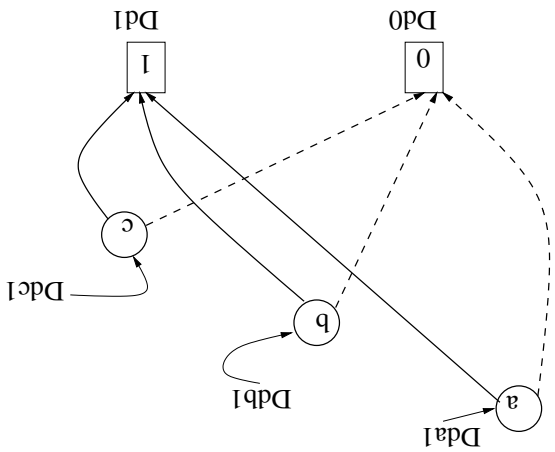
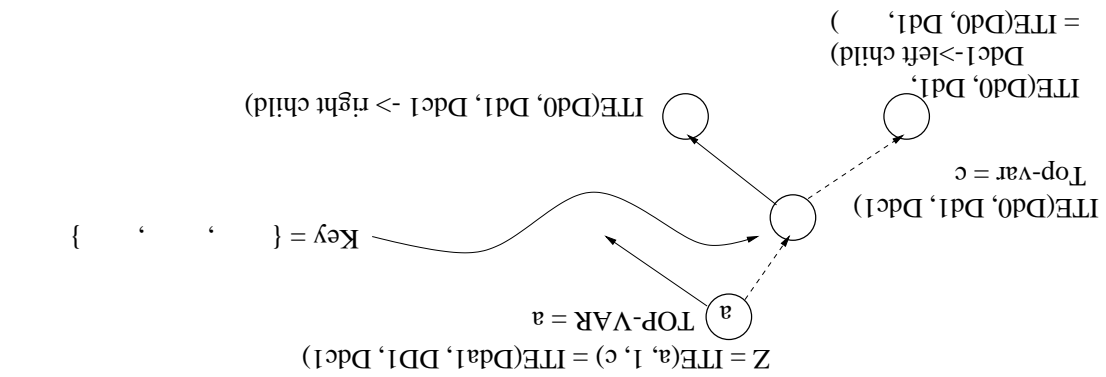
Construct $f = a + c$

Value	Key
Dd0	{NULL, 0, NULL}
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Dda1	{Dd0, a, Dd1}
Ddb1	{Dd0, b, Dd1}
Ddc1	{Dd0, c, Dd1}



$Z = \text{ITE}(a, 1, c) = \text{ITE}(Dda1, Dd1, Ddc1)$
 $\text{ITE}(Dda1 > \text{left child}, Dd1, Ddc1)$
 $= \text{ITE}(Dd0, Dd1, Ddc1)$
 TOP-VAR = a

- Go-up a recursion level and compute the right sub-tree

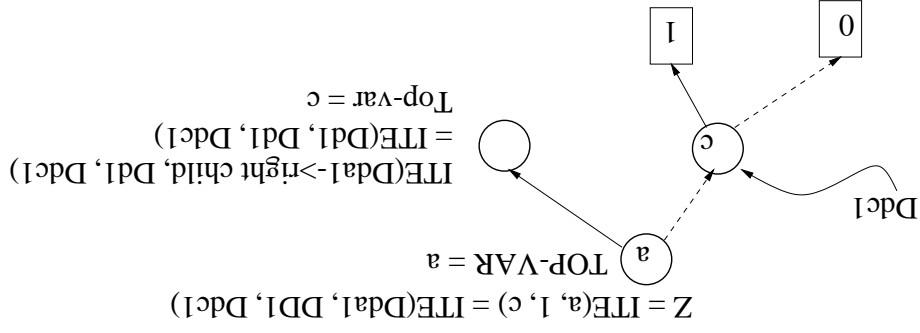
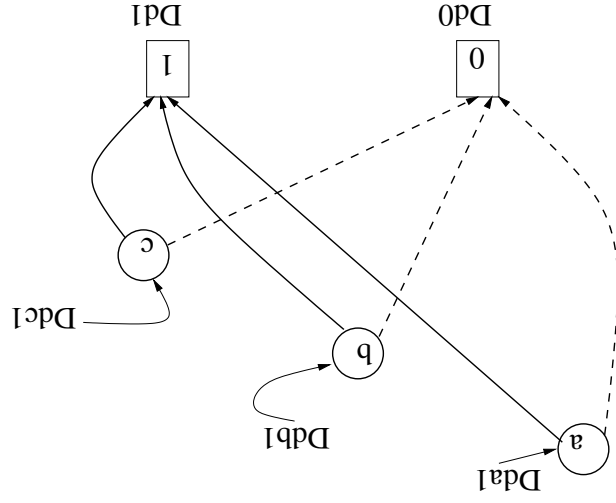


Value	Key
Dd0	{NULL, 0, NULL}
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$$f = a + c \text{ Contd...}$$

$f = a + c$ Contd. further...

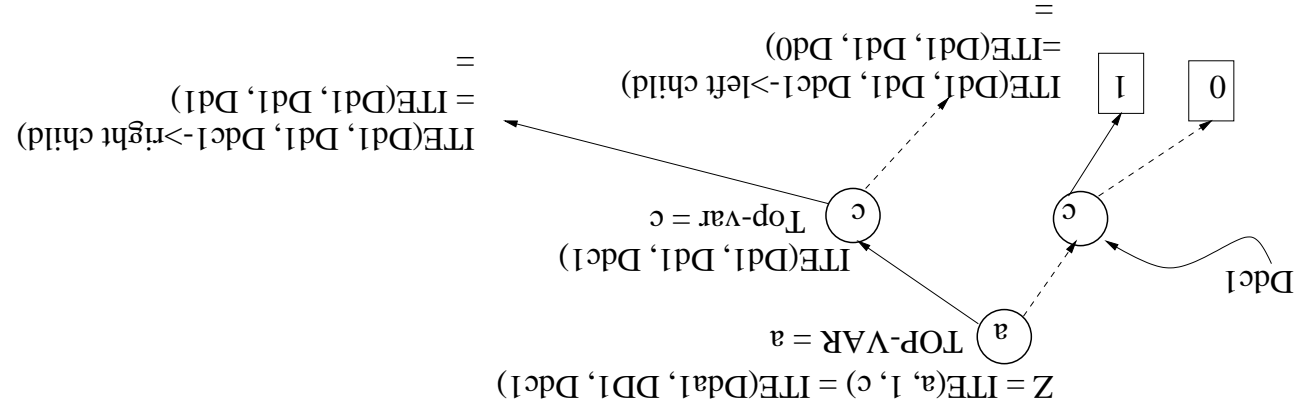
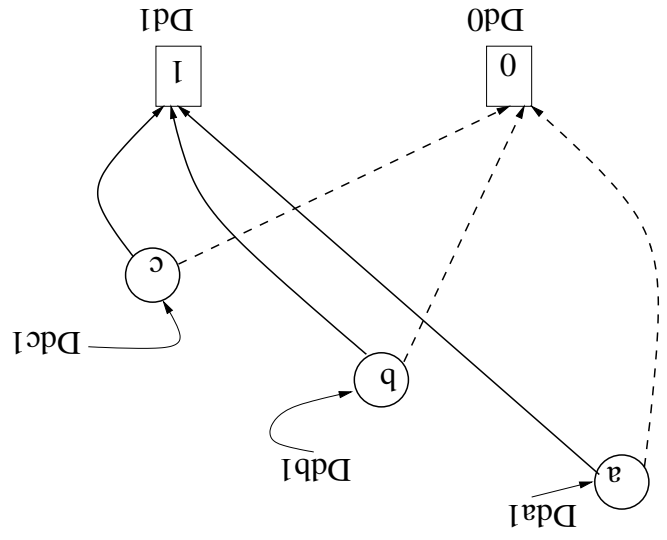
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Dd0	{NULL, 0, NULL}
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Dda1	{Dd0, a, Dd1}
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Ddc1	{Dd0, c, Dd1}



$$Z = ITE(a, 1, c) = ITE(Dda1, Dd1, Ddc1)$$

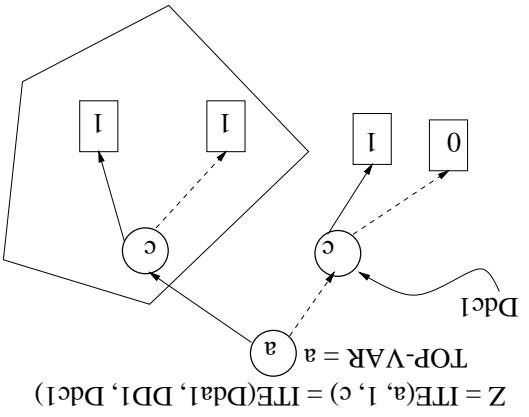
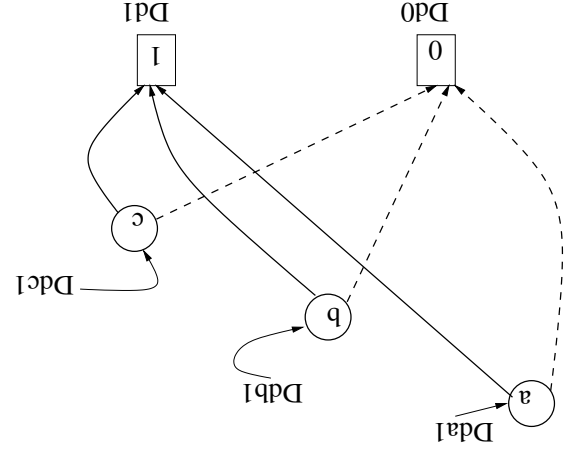
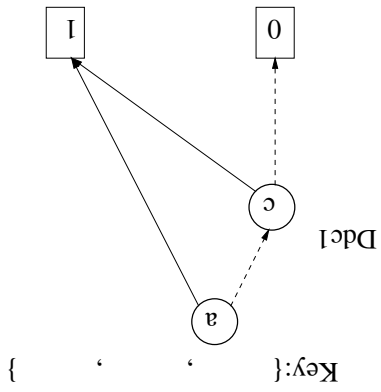
$f = a + c$ Contd. even further...

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Dd0	{NULL, 0, NULL}
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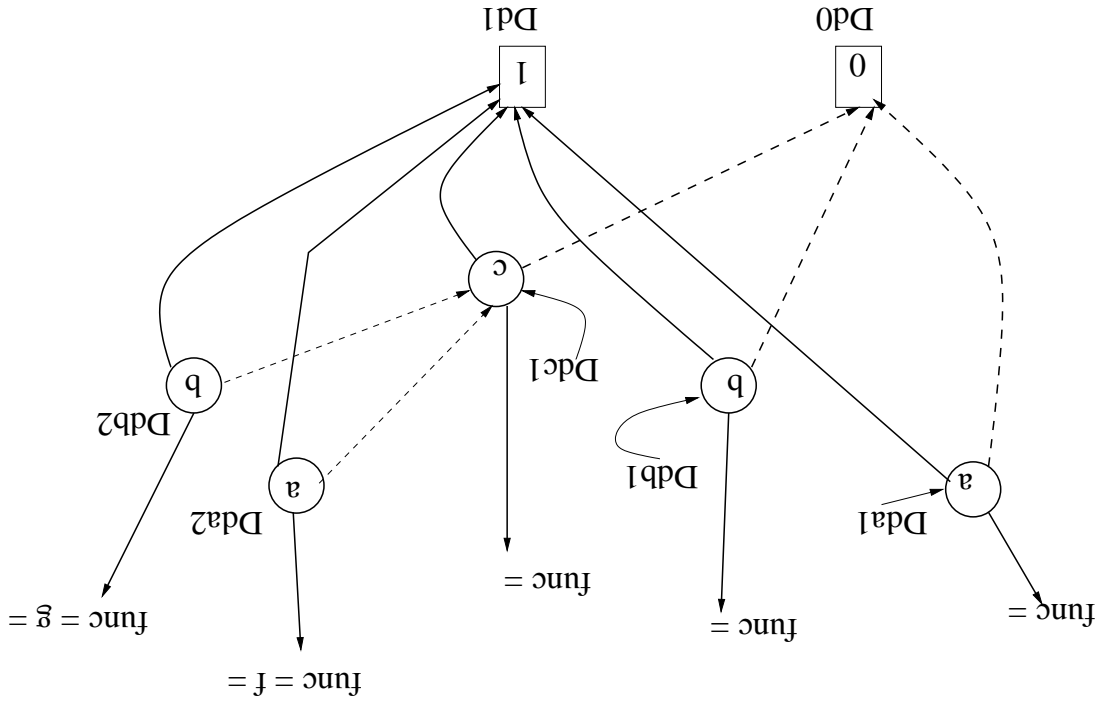


Final ROBDD for $f = a + c$

Value	Key
Dd0	{NULL, 0, NULL}
Dd1	{NULL, 1, NULL}
Dda1	{Dd0, a, Dd1}
Ddb1	{Dd0, b, Dd1}
Ddc1	{Dd0, c, Dd1}



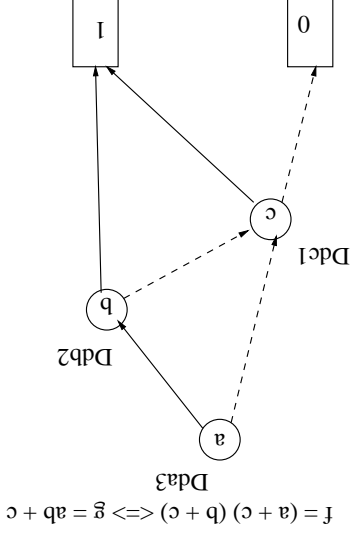
- Homework: Compute $f \cdot g$ yourself!
- Homework: Compute $(f = ab) \cdot (g = ab')$ yourself and notice how the graph reduces to $Dd0$.



$$f \cdot g = (a + c) \cdot (b + c)$$

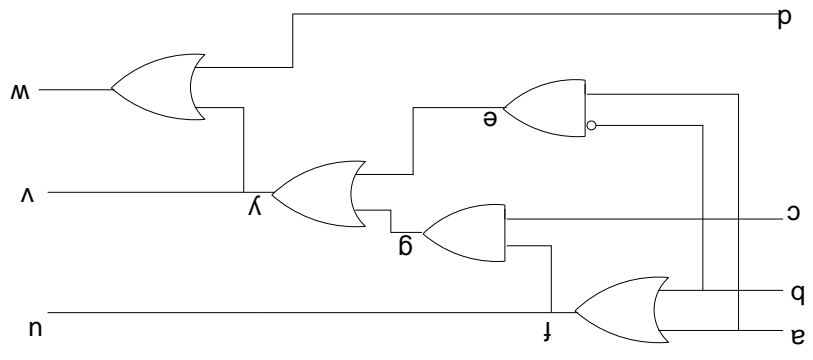
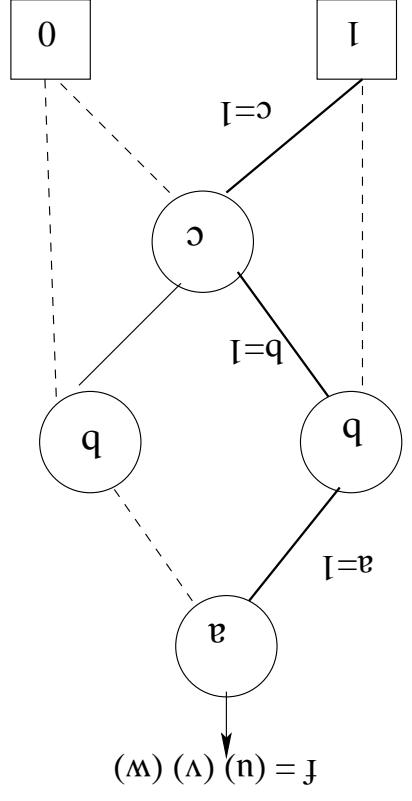
Equivalence Verification

- $f = (a + c)(b + c)$; $g = ab + c$
- Compute f , compute keys of every node, update symbol table
- Do the same with g and prove to yourself that $f = g$



Application to SAT

- For the circuit shown, $\text{SAT}(u=v=w=1) = ac + bc + ab$
- BDD $(u \cdot v \cdot w)$ shown. How to pick a solution?



BDD Size Problems and Variable Order

- Change Var order \rightarrow ROBDD structure and size changes!
- Worst-case ROBDD size \rightarrow no reduction \rightarrow full-blown tree.
- Given a Boolean function, how do we forecast a good var order?
- Intractable Problem! Though some heuristics exist....
- Var $x \in$ many cubes \rightarrow keep it up in the tree
- **Variable Ordering:** Interchange order of variables and see the change in size of ROBDD
- **Dynamic Var Ordering:** Do this while constructing OBDDs
- **Careful:** Var ordering of ALL the BDDs in the manager has to go through re-ordering.
- **Multipliers:** There exists NO good ordering. Take any order, at least one output would have worst-case scenario!

