

ECE 697B (667)

Spring 2003

Synthesis and Verification of Digital Systems

BDD-based Bi-decomposition *BDS system*

Outline

- Review of current decomposition methods
 - Algebraic
 - Boolean
- Theory of BDD decomposition [C. Yang 1999]
 - *Bi-decomposition* $F = D \oplus Q$
 - Boolean *AND/OR* Decomposition
 - Boolean *XOR* Decomposition
 - *MUX* Decomposition
- Logic optimization based on BDD decomposition

Functional Decomposition – previous work

- Ashenurst [1959], Curtis [1962]
 - Tabular method based on cut: bound/free variables
 - BDD implementation:
 - Lai *et al.* [1993, 1996], Chang *et al.* [1996]
 - Stanion *et al.* [1995]
- Roth, Karp [1962]
 - Similar to Ashenurst, but using cubes, covers
 - Also used by SIS
- Factorization based
 - SIS, algebraic factorization using cube notation
 - Bertacco *et al.* [1997], BDD-based recursive bidecomp.

Drawbacks of Traditional Synthesis Methods

- Weak *Boolean* factorization capability.
- Difficult to identify *XOR* and *MUX* decomposition.
- Separate platforms for Boolean operations and factorization.
- Our goal: use a common platform to carry out both Boolean operations and factorization: BDD's

What is wrong with Algebraic Division?

- Divisor and quotient are orthogonal!!
- Better factored form might be:

$$(q_1 + q_2 + \dots + q_n)(d_1 + d_2 + \dots + d_m)$$
 – g_i and d_j may share same or opposite literals

• Example:

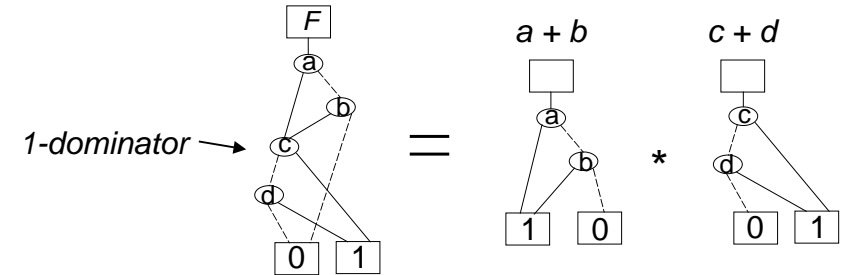
SOP form: $F = abg + acg + adf + aef + afg + bd + ce + be + cd.$ (23 lits)

Algebraic: $F = (b + c)(d + e + ag) + (d + e + g)af.$ (11 lits)

Boolean: $F = (af + b + c)(ag + d + e).$ (8 lits)

First work, Karplus [1988]: 1-dominator

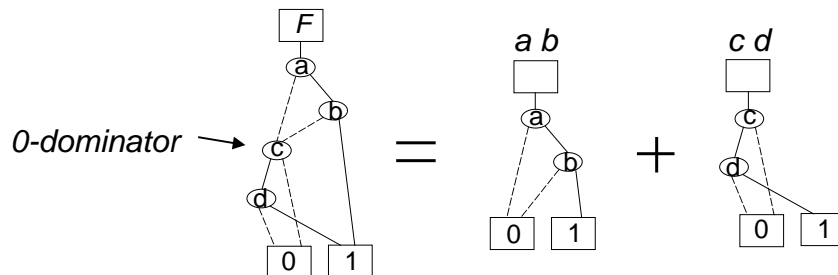
- Definition: *1-dominator* is a node that belongs to every path from root to terminal 1.



- *1-dominator* defines algebraic conjunctive (AND) decomposition: $F = (a+b)(c+d).$

Karplus: 0-dominator

- Definition: *0-dominator* is a node that belongs to every path from root to terminal 0.



- *0-dominator* defines algebraic disjunctive (OR) decomposition: $F = ab + cd.$

Bi-decomposition based on Dominators

- We can generalize the concept of algebraic decomposition and dominators to:
 - Generalized dominators
 - Boolean bi-decompositions (AND, OR, XOR)
- Bi-decomposition: $F = D \theta Q$
- First, let's review fundamental theorems for Boolean division and factoring.

Boolean Division

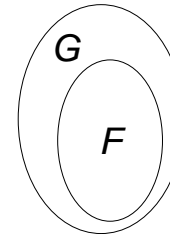
Definitions

- G is a *Boolean divisor* of F if there exist H and R such that $F = GH + R$, and $GH \neq 0$.
- G is said to be a *factor* of F if, in addition, $R=0$, that is: $F = GH$, where H is the quotient, R is the remainder.
Note: H and R may not be unique.

Boolean Factor - Theorem

Theorem:

Boolean function G is a *Boolean factor* of Boolean function F iff $F \subseteq G$, (i.e. $FG' = 0$, or $G' \subseteq F$).



Proof:

\Rightarrow : G is a Boolean factor of F . Then $\exists H$ s.t. $F = GH$; Hence, $F \subseteq G$ (as well as $F \subseteq H$).

\Leftarrow : $F \subseteq G \Rightarrow F = GF = G(F + R) = GH$. (Here R is any function $R \subseteq G'$.)

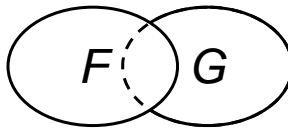
Notes:

- Given F and G , H is not unique.
- To get a small H is the same as getting a small $F + R$. Since $RG = 0$, this is the same as minimizing (simplifying) f with $DC = G'$.

Boolean Division - Theorem

Theorem:

G is a *Boolean divisor* of F if and only if $FG \neq 0$.



Proof:

\Rightarrow : $F = GH + R$, $GH \neq 0 \Rightarrow FG = GH + GR$. Since $GH \neq 0$, $FG \neq 0$.

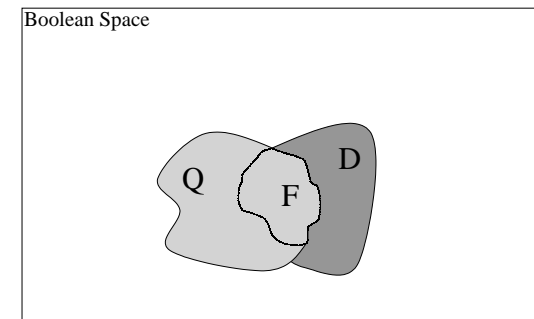
\Leftarrow : Assume that $FG \neq 0$. $F = FG + FG' = G(F + K) + FG'$. (Here $K \subseteq G'$.) Then $F = GH + R$, with $H = F + K$, $R = FG'$. Since $GH = FG \neq 0$, then $GH \neq 0$.

Note:

f has many divisors. We are looking for a g such that $f = gh + r$, where g, h, r are simple functions. (simplify f with $DC = g'$)

Boolean Division

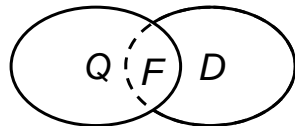
Goal : for a given F , find D and Q such that $F = Q \cdot D$.



$$F = e + bd, \quad D = e + d, \quad Q = e + b$$

Conjunctive (AND) Decomposition

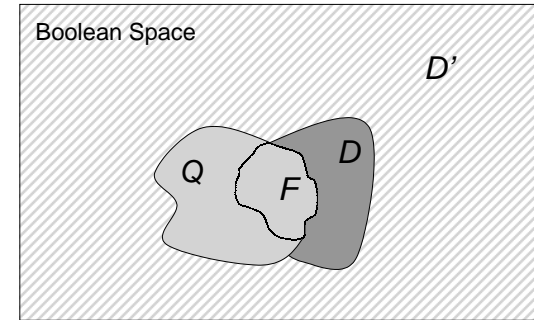
- Conjunctive (AND) decomposition: $F = D Q$.
- Theorem:
Boolean function F has conjunctive decomposition iff $F \subseteq D$. For a given choice of D , the quotient Q must satisfy: $F \subseteq Q \subseteq F + D'$.



- For a given pair (F, D) , this provides a recipe for Q .

Boolean Division \Rightarrow AND decomposition

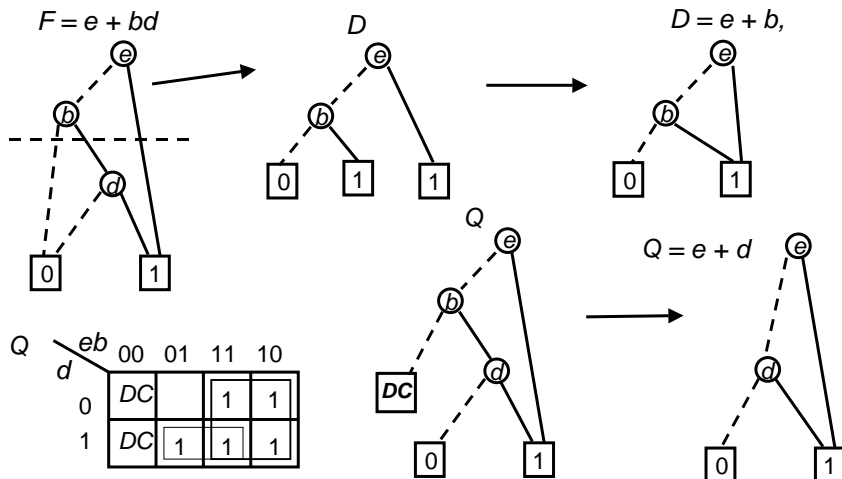
- Given function F and divisor $D \supseteq F$, find Q such that:
 $F \subseteq Q \subseteq F + D'$.



$$F = e + bd, \quad D = e + d, \quad Q = e + b \subseteq F + D'$$

AND Decomposition ($F = D Q$): Example

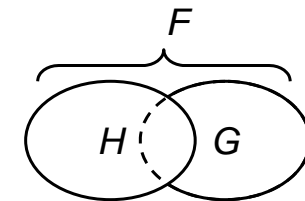
- Recall: given (F, D) , quotient Q must satisfy: $F \subseteq Q \subseteq F + D'$.



Disjunctive (OR) Decomposition

- Disjunctive (OR) decomposition: $F = G + H$.
- Theorem:
Boolean function F has disjunctive decomposition iff $F \supseteq G$. For a given choice of G , the term H must satisfy: $F' \subseteq H' \subseteq F' + G$.

Dual to conjunctive decomposition.

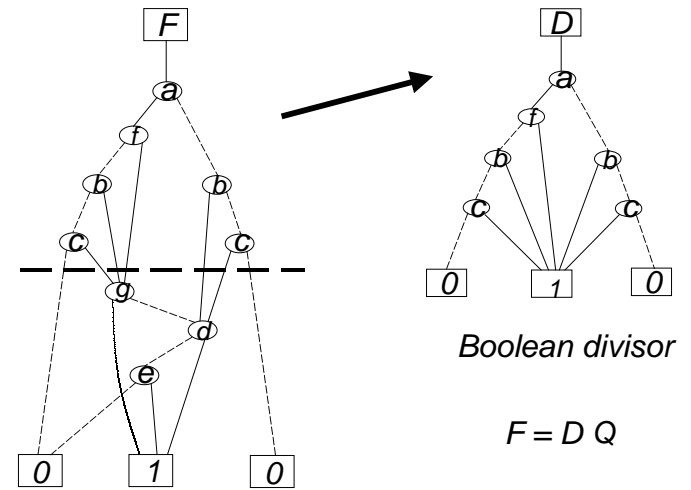


- For a given (F, G) , this provides a recipe for H .

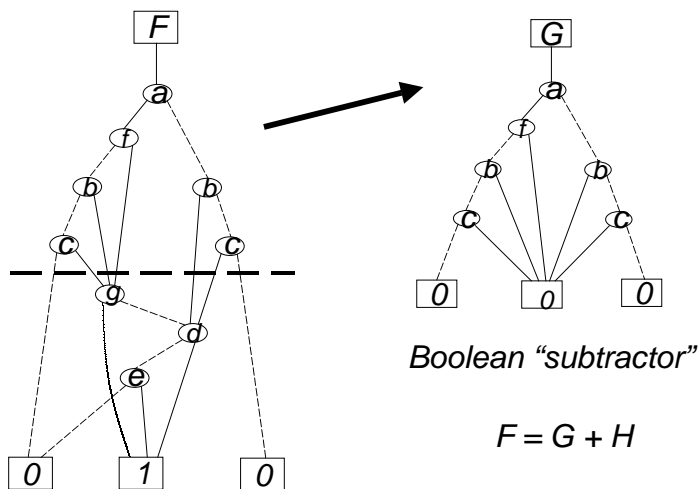
Boolean AND/OR Bi-decompositions

- Conjunctive (AND) decomposition
If $D \supseteq F$, $F = FD = QD$.
- Disjunctive (OR) decomposition
If $G \subseteq F$, $F = F + G = H + G$.
- D, G = generalized dominators

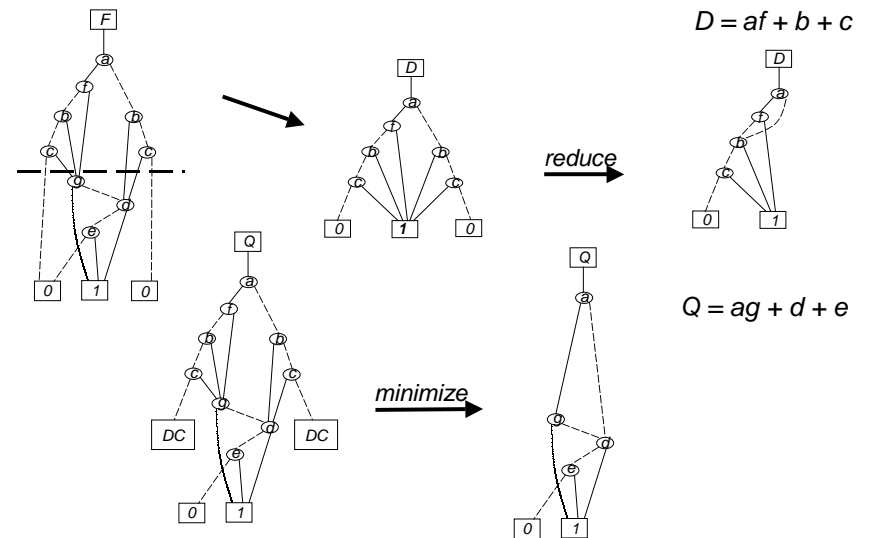
Generalized Dominator D



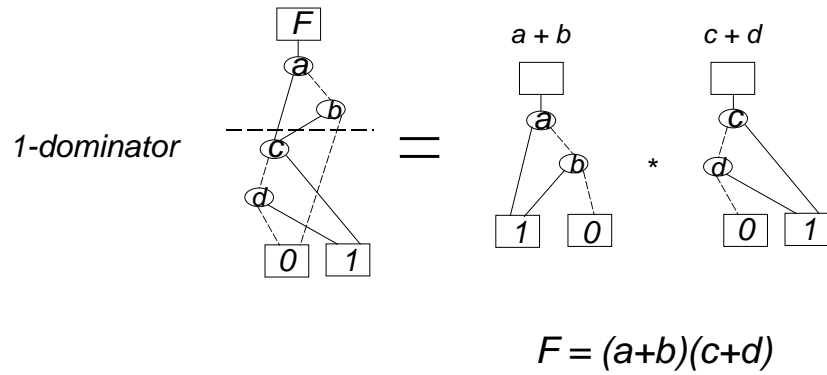
Generalized Dominator G



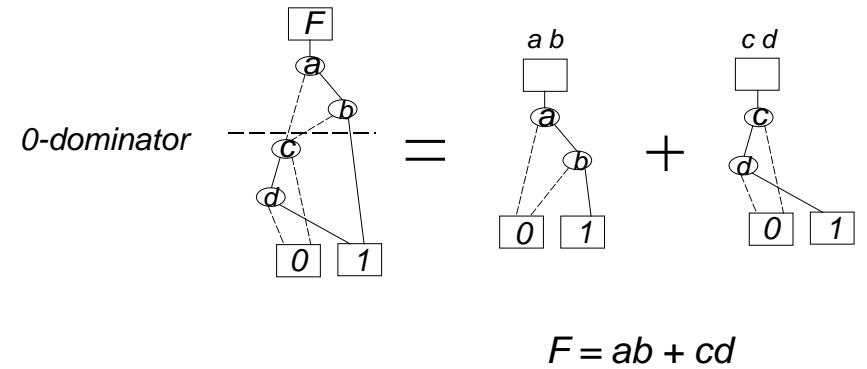
Boolean Division Based on Generalized Dominator



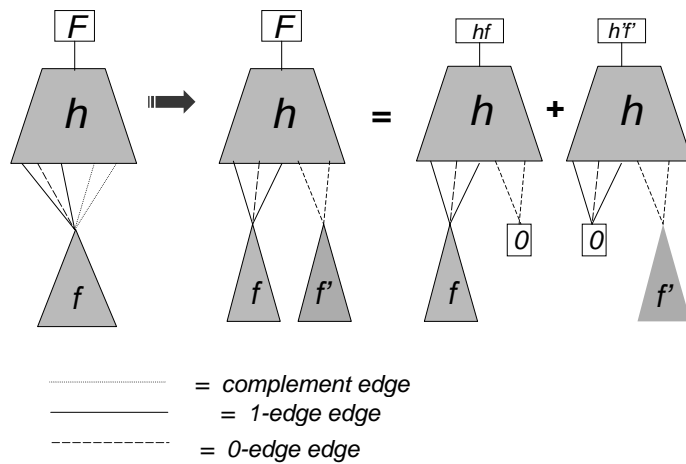
Special Case: 1-dominator



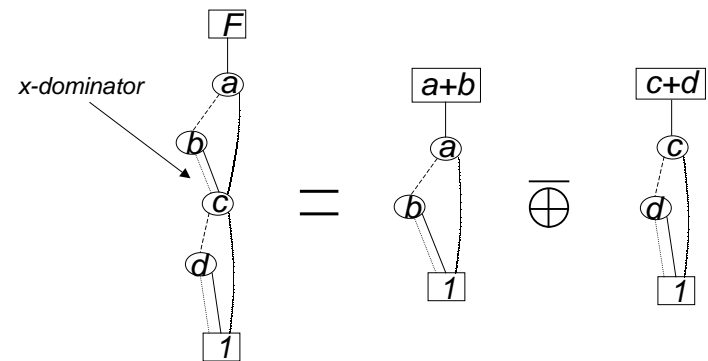
Special Case: 0-dominators



Algebraic XOR Decomposition

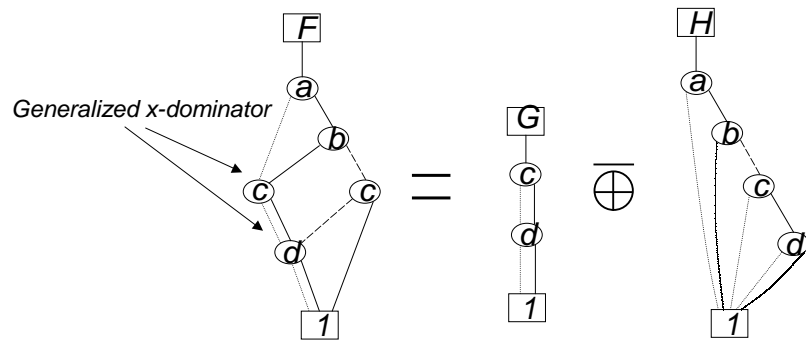


Algebraic XOR Decomposition: x -dominators



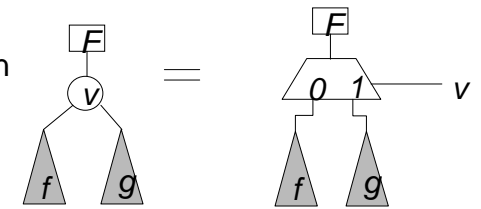
Boolean XOR Decomposition: Generalized x -dominators

Given F and G , there exists $H: F = G \otimes H; H = F \otimes G$.

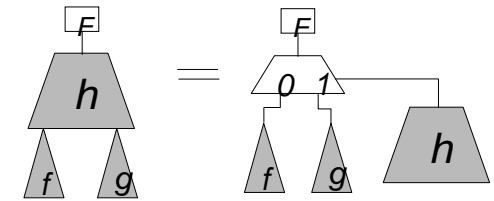


MUX Decomposition

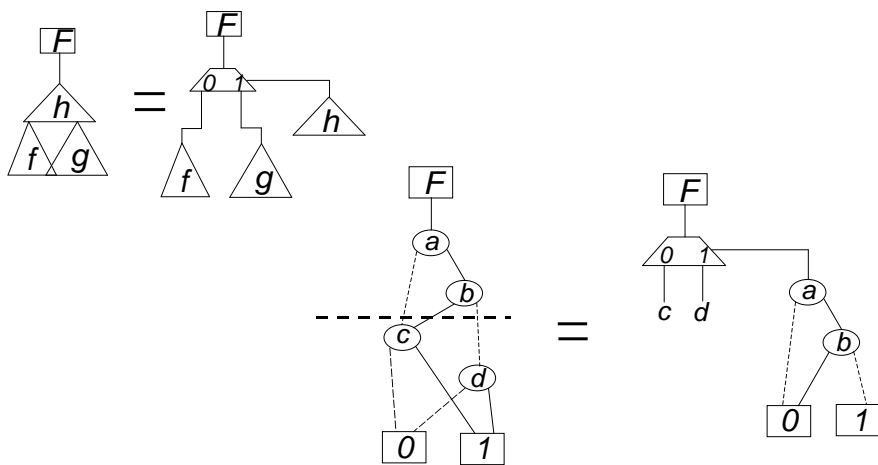
- Simple MUX decomposition



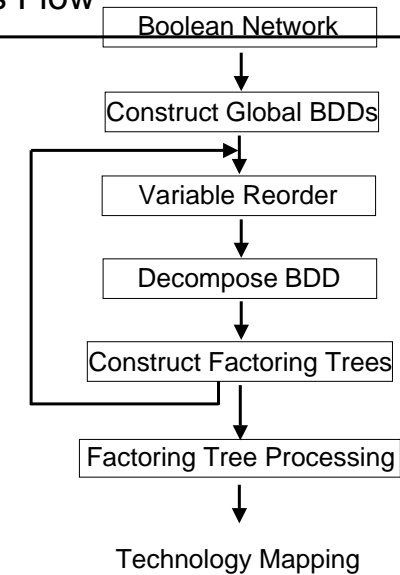
- Complex MUX decomposition



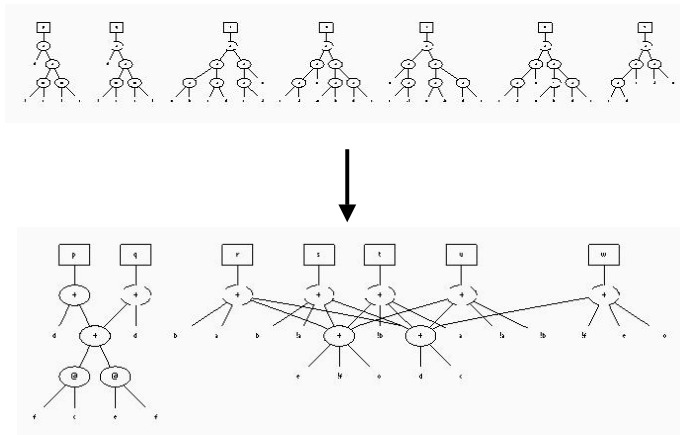
Functional MUX Decomposition - example



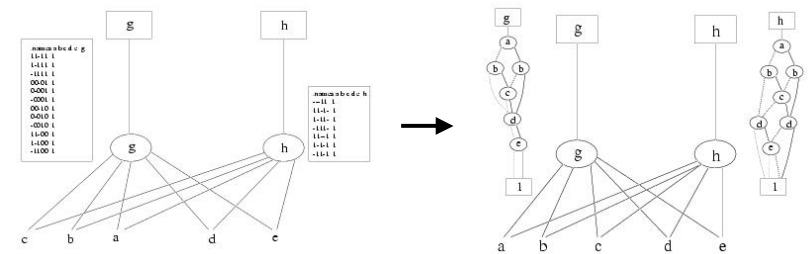
Synthesis Flow



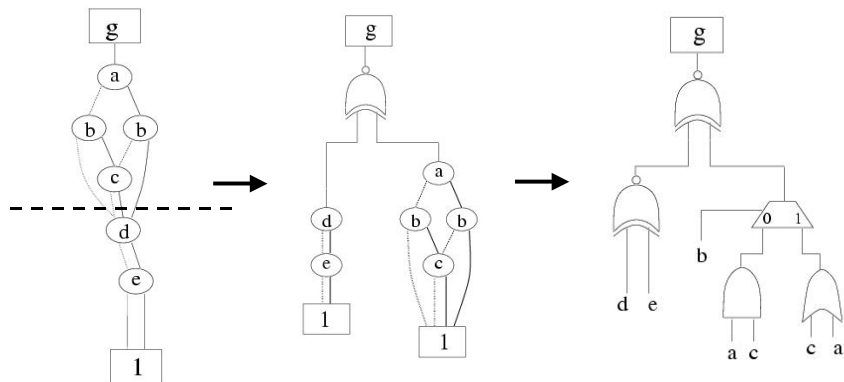
Factoring Tree Processing:



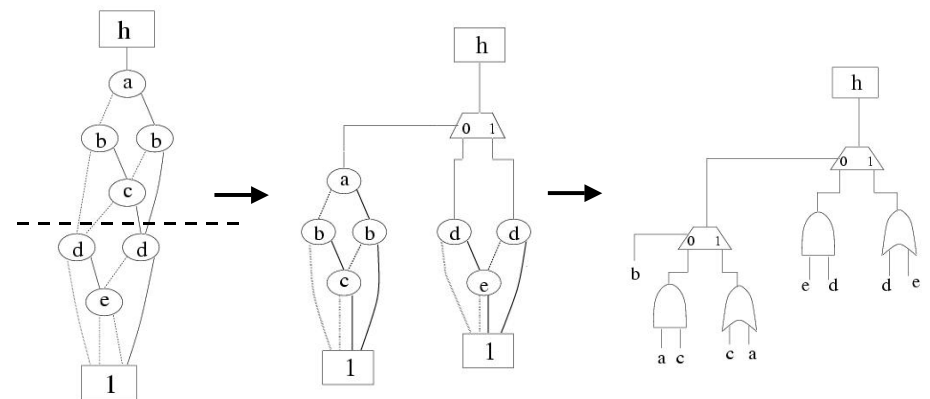
A Complete Synthesis Example



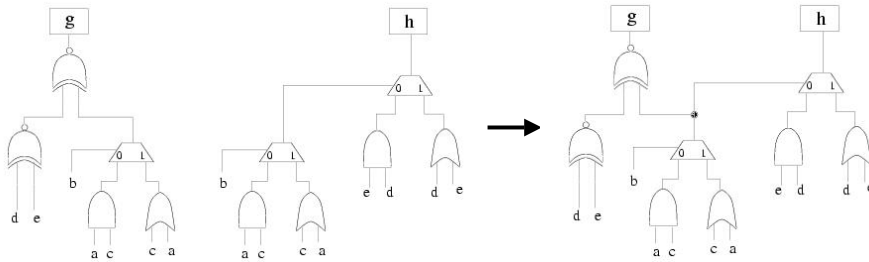
A Complete Synthesis Example (Decompose function g)



A Complete Synthesis Example (Decompose function h)



A Complete Synthesis Example (Sharing Extraction)



Conclusions

- BDD-based *bi-decomposition* is a good alternative to traditional, algebraic logic optimization
 - Produces *Boolean* decomposition
 - Several types: *AND, OR, XOR, MUX*
- BDD decomposition-based logic optimization is *fast*.
- Stand-alone BDD decomposition scheme is not amenable to large circuits
 - *Global BDD* too large
 - Must partition into network of BDDs (local BDDs)