### Software Verification

CS 6110, taught by Pavel Panchekha



### Pavel Panchekha

T/R 14:00-15:30 WEB L126

### Software Verification

# NO BUGS Bugs are bad-

## Bugs are bad

Software bugs are **dangerous and costly** 

Security bugs, down-time, crashes, lost productivity **\$59,000,000,000 annual cost** (NIST estimate)

## Bugs are bad



### Therac 25



#### **Radiation therapy software**

Atomic Energy of Canada Race condition **kills 5** 

## **ION Therapy Planning**



#### **Dosage calculation for radio-therapy**

Panamanian health organization **8 killed**, 20 injured, FDA investigation

### Ariane 5



#### **Ariane 4 rocket controller reused**

Faster rocket causes integer overflow Flight computer crashes 40 seconds in

## 2003 BMW Crash



### Thai Finance Minister **trapped in car** Computer crashes, doors stuck Firemen **smashed windows**

### 2005 Prius Recall



### Cars stall, stop at highway speeds

Software bug causes computer crash **75,000 cars recalled** 

## What is a Bug?

- Feature request
- Confusing features
- Ugly, unintuitive
- Unusably slow, memory hog
- Confusing error messages
- Unhandled exceptions
- Segmentation fault

## What is a Bug?

Feature request

Confusing features

### Specification

A bug is a mismatch between code and its specification

Unhandled exceptions

Segmentation fault

## Specifications

Exact description of what a program should do

What it does, not how it does it

quicksort(l) // Spec: sorted(output)

Intermediate specifications

What needs to be true of **partition** for **quicksort** to work?

## **Program Reasoning**

Spec: what a program **should** do-what **does** it do?

Matching program execution to clauses in spec

// len(l) == 3 and l[0] == 5
quicksort(l)
// len(l) == 3 and ???

Program reasoning is running the program "in your head"

You can run it **backwards or forwards** 

## Scaling up

### How do you choose which specs to prove?



#### Choose the specificity of your specification

Fine-grained or coarse-grained specifications can be best

### About this Course

Lectures, Assignments, Projects, Grading

## Your Instructor



### Pavel Panchekha PhD in 2019 < I'm new here! U. of Washington

## I work on **the browser as a platform** for executing programs with PL techniques

## Your TA



### Manasij Mukherjee PhD student Advisor: John Regehr

Works on compiler optimization using **program synthesis** and **static analysis**.

### **Three Parts**

## Write a specification

### Propagate into program

Scale to systems

### **Three Parts**

## Logical reasoning

Program logics

Static analysis

## **Six Topics**



Deciding truth

Axiomatic reasoning

Dependent types

## Six Assignments

Logical reasoning

Program logics

## Static analysis

First-order logic
N-Queens with miniSAT

Deciding truth

Kenken with Z3

Symbolic execution
Bug-finding with KLEE

Abstract domains Analysis with Checkers

Axiomatic reasoning

Verification with Dafny

Dependent types

## Six Assignments

Assignments every two weeks ish

### Due Thursdays at midnight

Static analysis

Late policy: 20 point deduction without 48hr notice

Each assignment **10% of your final grade** 

No assignment for dependent types because projects

Kenken with Z3

**/erification with Dafny** 

## **Course Project**

### Verify something, build a verifier, contribute, etc...

Proposal Draft	Early February
<b>Proposal Presentation</b>	Mid-February
Milestone 1	Early March
Milestone 2	Late March
Final presentation	Mid April

### Project is 40% of your final grade

**Question:** Team projects? Alternative assignment?

## Participation

10% of your final grade, half attendance half feedback

**24hr notice** for absence from class (email is fine)

Feedback form mandatory every class

Answers aren't graded, I'll pretend they're anonymous

A few standard questions plus **class polls** 

## Web Pages



### my.eng.utah.edu/~cs6110



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### How to Verify

From specification to verification

def **quicksort**(1):

**pivot** = l[len(l)//2]

left, right = partition(l, pivot)

left2 = quicksort(left)

right2 = quicksort(right)

**return** left2 + right2

Post: sorted(output)
def quicksort(l):



**pivot** = 1[len(1)//2]

left, right = partition(l, pivot)

left2 = quicksort(left)

right2 = quicksort(right)

**return** left2 + right2

Post: sorted(output)
def quicksort(l):

**pivot** = 1[len(1)//2] **Prove:** left[i] < right[j]

left, right = partition(l, pivot)

?

left2 = **quicksort**(left) ← sorted(left2)

**return** left2 + right2

Post: sorted(output)
def quicksort(l):

Nope! Is <u>left2[i] < right2[j]</u>, or just <u>left</u> and <u>right</u>?

**pivot** = l[len(l)//2]

right2 = **quicksort**(right) ← sorted(right2)

return left2 + right2 ← sorted(left2 + right2)

Post: sorted(output) and output[i] in I def quicksort(1):

**pivot** = 1[len(1)/2] **Prove:** left[i] in I and right[j] in I

right2 = **quicksort**(right) ← sorted(right2)

return left2 + right2 ← sorted(left2 + right2)

Post: sorted(output) and output[i] in I def quicksort(1): pivot = 1[len(1)//2] The right specification?

left[i] < right[j] and left, right = **partition**(l, pivot) ← left[i] in I and right[j] in I

left2 = **quicksort**(left) ← sorted(left2) and left2[i] in I

right2 = **quicksort**(right) ← sorted(right2) and ...

**return** left2 + right2 ← **sorted**(left2 + right2) and ...

Post: sorted(output) and output[i] in I and I[i] in output def **quicksort**(1):

return []

Nope!

Post: sorted(output) and output[i] in I and I[i] in output def quicksort(1): and len(I) = len(output)

return [min(l)] + sort(l)

Nope!

Post: sorted(output) and output[i] in I and I[i] in output def **quicksort**(1): and **len(I) = len(output)** 

### if **1 == [1, 2, 3, 3, 4]**:

return [1, 2, 3, 4, 4]

else:

return sort(l)

Nope!

Post: sorted(output) and l.count(x) == output.count(x) def quicksort(1):

### if 1 == [1, 2, 3, 3, 4]: return [1, 2, 3, 4, 4]

else:

Is that finally enough?

return sort(l)

## Challenges

Writing a specification for quicksort

Reasoning about predicates like sorted

**Combining facts** about lists and predicates

Propagating facts through the program

Expanding the specification to make it provable

### **Verification is hard!**

### **Recent Successes**



## SAGE

Static analysis based fuzzing for Windows applications

500+ years of machine time

**Hundreds** of security issues identified **One in three** bugs in Win7 WEX identified by SAGE

### Millions of dollars saved



### Verify drivers don't crash Windows

... in some very key areas, for example, driver verification we're building tools that can do actual proof about the software ...

- Bill Gates, WinHec 2002 Keynote

### Standard part of driver SDK

Prevented hundreds of blue screen bugs

### IronFleet

### Verified implementation of **distributed system**



Provably resistant to hardware failure

Found bugs in standard algorithms

### Next class: Boolean Logic

To do:
Course feedback
Register for Piazza
Find course webpage