

Last Time

- ◆ Looked at ColdFire and ARM in depth

Today

- ◆ **Tools and toolchains for embedded systems**
 - **Linkers**
 - **Programmiers**
 - **Booting an embedded CPU**
 - **Debuggers**
 - **JTAG**
- ◆ **All of this stuff is “below” the C compiler in the stack of tools**
 - **Material on embedded C will follow**
- ◆ **Any weak link in the toolchain will hinder development**

Economic Context

- ◆ **Dev. tools for general-purpose systems:**
 - **Mass-market users: Lots of them, so compiler gets tested thoroughly**
 - **ISVs: Sell popular programs, so executables are widely tested**
- ◆ **Dev. tools for embedded systems:**
 - **One of these categories does not exist**
- ◆ **Hard to make money selling embedded toolchains**
 - **A few, large sales**
 - **In many cases, tools are thrown in with the architecture license**

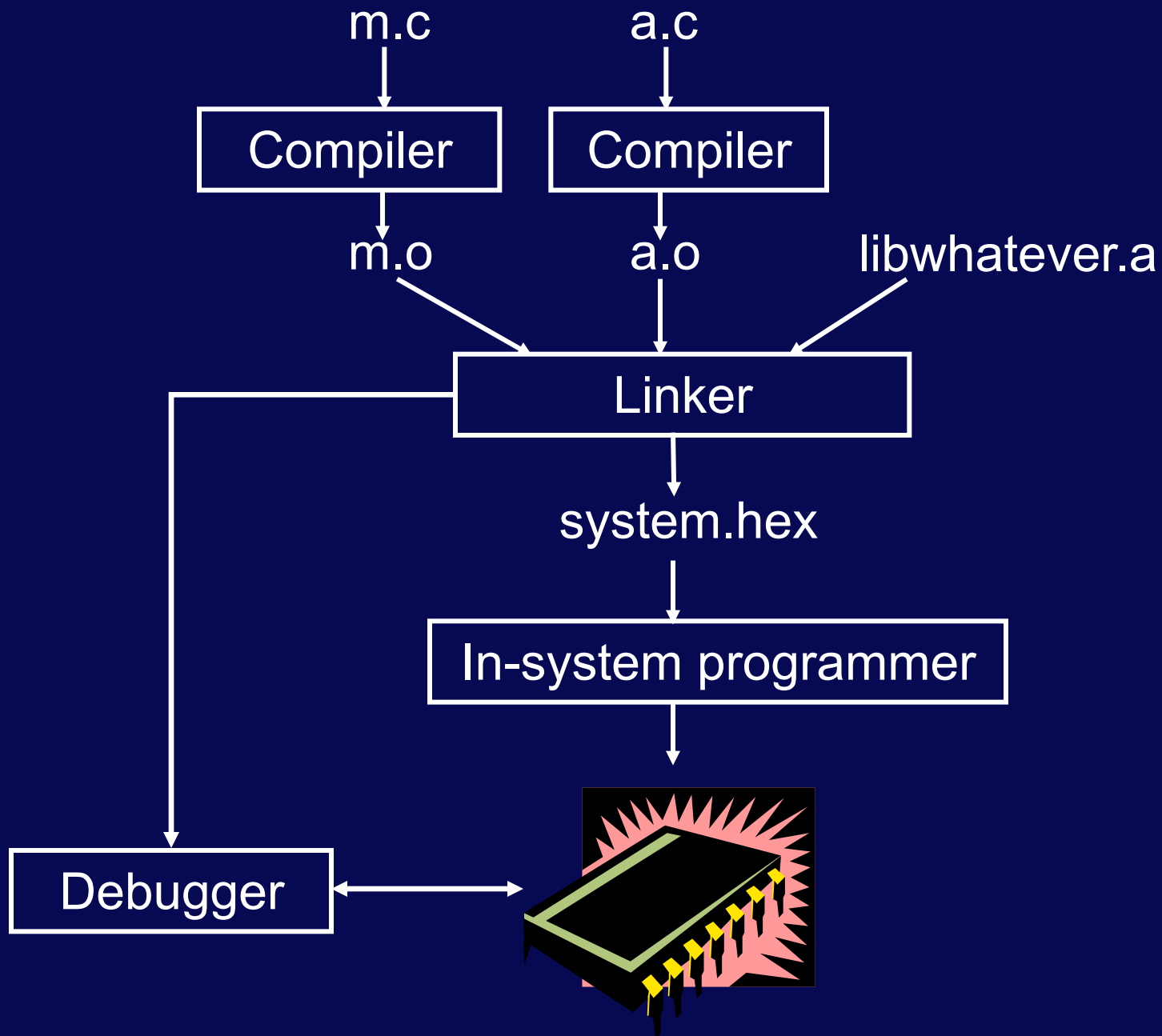
Economic Context

- ◆ **Open source tools have changed things quite a bit**
 - **GCC, mainly, but other tools too**
 - **GCC targets ~40 architectures**
- ◆ **Often, a company pays to have GCC ported to some embedded architecture**
 - **Eventually results are open sourced**
 - **ARM has ~8 people working full-time on GCC**
 - **Why would they do this? Keep in mind they also sell compiler tools**

More Economic Context

- ◆ **Problem: Embedded tools often not very high quality**
 - Small number of expert users
 - Lots and lots of chips to support
 - This fact is independent of whether the tools are open-source or not

- ◆ **Read the Wolfe article linked to the course web page**
 - Great article



Linking Background

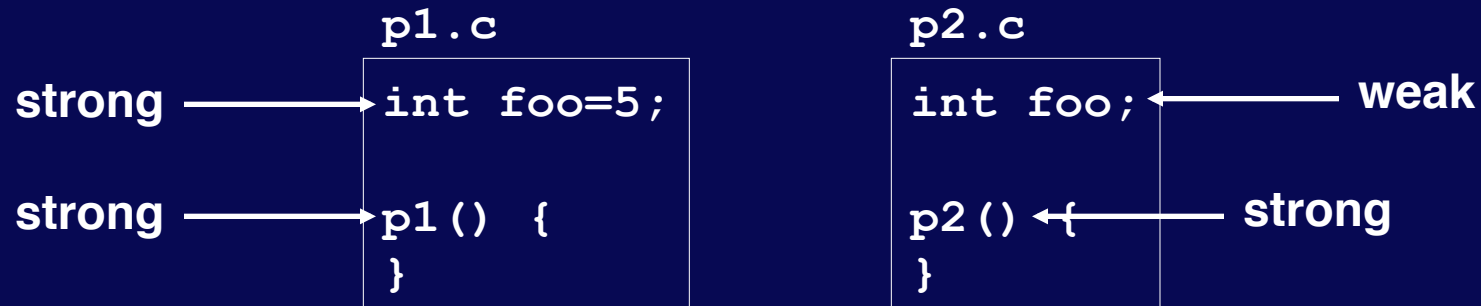
- ◆ Each .c file, plus any headers it includes, is called a “compilation unit”
 - Compiler turns compilation unit into an object file
- ◆ Each object (.o) file contains:
 - text segment – executable code
 - data segment – initialized data
 - BSS segment – uninitialized data
 - Other stuff – debugging symbols, etc.
- ◆ Object files:
 - Relocatable
 - Code and data addresses are symbolic – not yet bound to physical addresses
 - Contain unresolved references

Linking

- ◆ **Linker functions**
 1. **Merge text, data, BSS segments of individual object files**
 - Including libraries
 - Including processor boot code
 2. **Resolve references to code and data**
 - Report any errors
 3. **Locate relocatable code**
 - Follow instructions in linker script
 - Report any errors
- ◆ **Result: Binary image ready to be loaded onto the target system**

Linker Operation

- ◆ **Classify all program symbols as either:**
 - **Weak** – uninitialized globals
 - **Strong** – functions and initialized globals



- ◆ **Scan object files in order supplied to the linker, applying linker rules**
 - **Bizarre consequence: Same object file might have to appear on command line multiple times**

Linker Operation

1. **A strong symbol can only appear once**
 - otherwise error
 2. **A weak symbol is overridden by a strong symbol of the same name**
 - I.e. all references to that name resolve to the strong symbol
 3. **If there are multiple weak symbols, the linker can pick an arbitrary one**
 - uh oh
- ◆ **Lots more details in CS 4400**

Linker Scripts

- ◆ **GNU linker is flexible and powerful**
 - Needs a “program” to tell it how to link for a given embedded platform
- ◆ **Linker script functionality:**
 - Put parts of executable into the right parts of memory
 - Insert padding to meet alignment requirements
 - Define extra symbols
 - Do arithmetic
 - Keep track of current position in memory as “.”

RPi bare metal linker script

```
SECTIONS {
    .init 0x0000 : {
        *(.init)
    }
    .text 0x8000 : {
        *(.text)
    }
    .data : {
        *(.data)
    }
    /DISCARD/ : {
        *(*)
    }
}
```

MCF52233 Linker Script

```
MEMORY {
    code (RX) : ORIGIN = 0x00000500,
               LENGTH = 0x0003FB00
    userram (RWX) : ORIGIN = 0x20000400,
                  LENGTH = 0x00007C00
}

SECTIONS {
    ___heap_size = 0x1000;
    ___stack_size = 0x1000;
}
```

More Linker Script

```
RAMBAR          = 0x20000000;
RAMBAR_SIZE     = 0x00008000;

FLASHBAR       = 0x00000000;
FLASHBAR_SIZE  = 0x00040000;

.vectors :
{
    mcf5xxx_vectors.s (.text)
    . = ALIGN (0x4);
} >> vectorrom

.text :
{
    *(.text)
    . = ALIGN (0x4);
    *(.rodata)
    . = ALIGN (0x4);
    ___ROM_AT = .;
    ___DATA_ROM = .;
} >> code
```

More Linker Script

```
data :
{
    __DATA_RAM = .;
    . = ALIGN(0x4);

    __sinit__ = .;
    STATICINIT
    __START_DATA = .;

    *(.data)
    . = ALIGN (0x4);
    __END_DATA = .;
} >> userram
```

More Linker Script

```
.bss :  
{  
    __START_BSS = . ;  
    *(.bss)  
    . = ALIGN (0x4) ;  
    *(COMMON)  
    __END_BSS = . ;  
    __BSS_END = . ;  
  
    . = ALIGN(0x4) ;  
} >> userram
```


Loading Programs

- ◆ **Goal: Set things up so CPU runs the desired program when powered up**
- ◆ **How is this done?**
 - **Make a ROM, plug it in**
 - **Burn a PROM / EPROM / EEPROM, plug it in**
 - **Download into RAM or flash ROM using an ISP**
 - **“In system programmer”**
 - **Load new code over a network**
- ◆ **Pros and cons of each?**

Booting a CPU

- ◆ **Execute a sequence of steps**
 - **May run in different orders in different systems**
 - **Some steps optional**
- ◆ **Usually want to cope with both hard and soft boot**

Bootup Steps

1. **Disable all interrupts**
 - ◆ Most processors power up with interrupts off
 - ◆ However – may be a soft reboot
2. **Perform RAM and ROM checks**
 - ◆ RAM – “walking 1s” test or similar
 - ◆ ROM – checksum
 - ◆ No point proceeding if one of these fails
3. **Initialize devices to known states**
4. **Copy initialized data segment from ROM to RAM**
5. **Clear BSS – uninitialized data segment**

More Booting

6. Initialize the stack
 - ◆ Initialize the stack pointer
 - ◆ Create initial stack frame
7. Initialize the heap
8. Execute constructors and initializers for all global variables
9. Enable interrupts
10. Call main()
11. Deal with the fact that main exited

Debugging

- ◆ **Important capabilities:**
 - **Observability – See internal processor state**
 - **Real-time analysis – Follow execution without slowing it down or stopping it**
 - **Run control – Start and stop the processor, set breakpoints, watches, etc.**

- ◆ **For each debugging method:**
 - **Which capabilities does it provide?**
 - **What are its other pros and cons?**

Debugging Methods

- ◆ **LEDs under software control**
 - Minimal workable debugging environment
 - A most unpleasant way to debug complex software

- ◆ **printf() to serial console or LCD**
 - Severely perturbs timing, typically
 - Generally, a debug printf() is synchronous
 - Means: Hangs the system until the printf completes
 - Why?

More Debugging

- ◆ **Logic analyzer hooked to external pins**
 - **Timing mode – displays logic transitions on pins**
 - **State mode – decode executing instructions, bus transactions, etc.**
 - **Triggers – give the analyzer conditions on which to start a detailed trace**
 - **Triggers can be highly elaborate**

- ◆ **Remote debugger**
 - **Debugging stub runs on embedded processor**
 - **Main debugger (e.g., GDB) runs on a separate machine**
 - **The two communicate using Ethernet, serial line, or whatever**

More Debugging

- ◆ **JTAG, BDM, Nexus**
 - Basically just hardware implementations of debugging stubs
- ◆ **ICE – in-circuit emulator**
 - Acts like your embedded processor but provides lots of extra functionality
 - Runs at full speed
 - Typically expensive

More Debugging

◆ ROM emulator

- Looks like ROM, actually RAM + processor
- At minimum supports rapid loading of new SW
- Can implement breakpoints, execution tracing

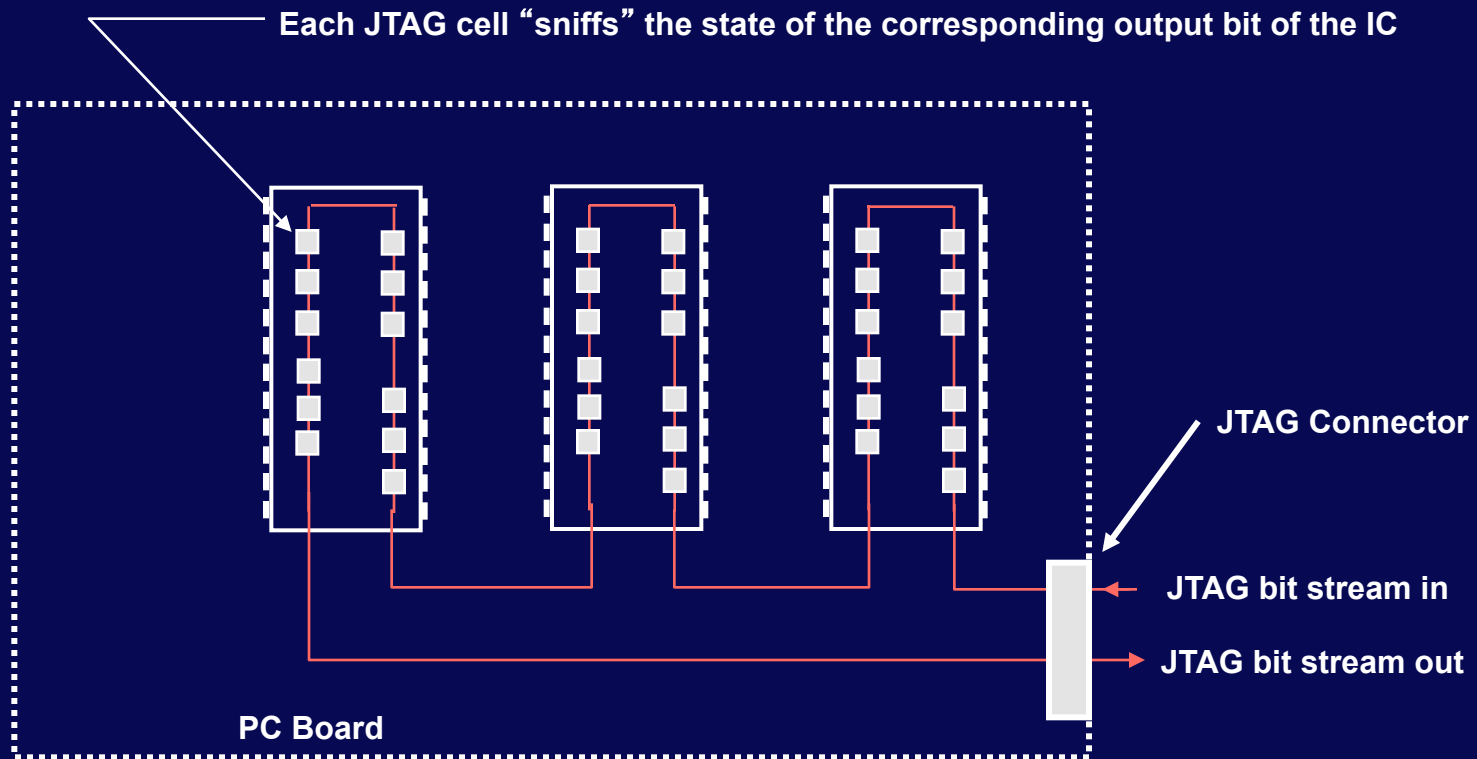
◆ Simulator

- Maximum controllability and observability
- Often slow
- Hard to interface to the real world
- Easy to simulate the CPU, hard to simulate everything else

JTAG – IEEE1149.1

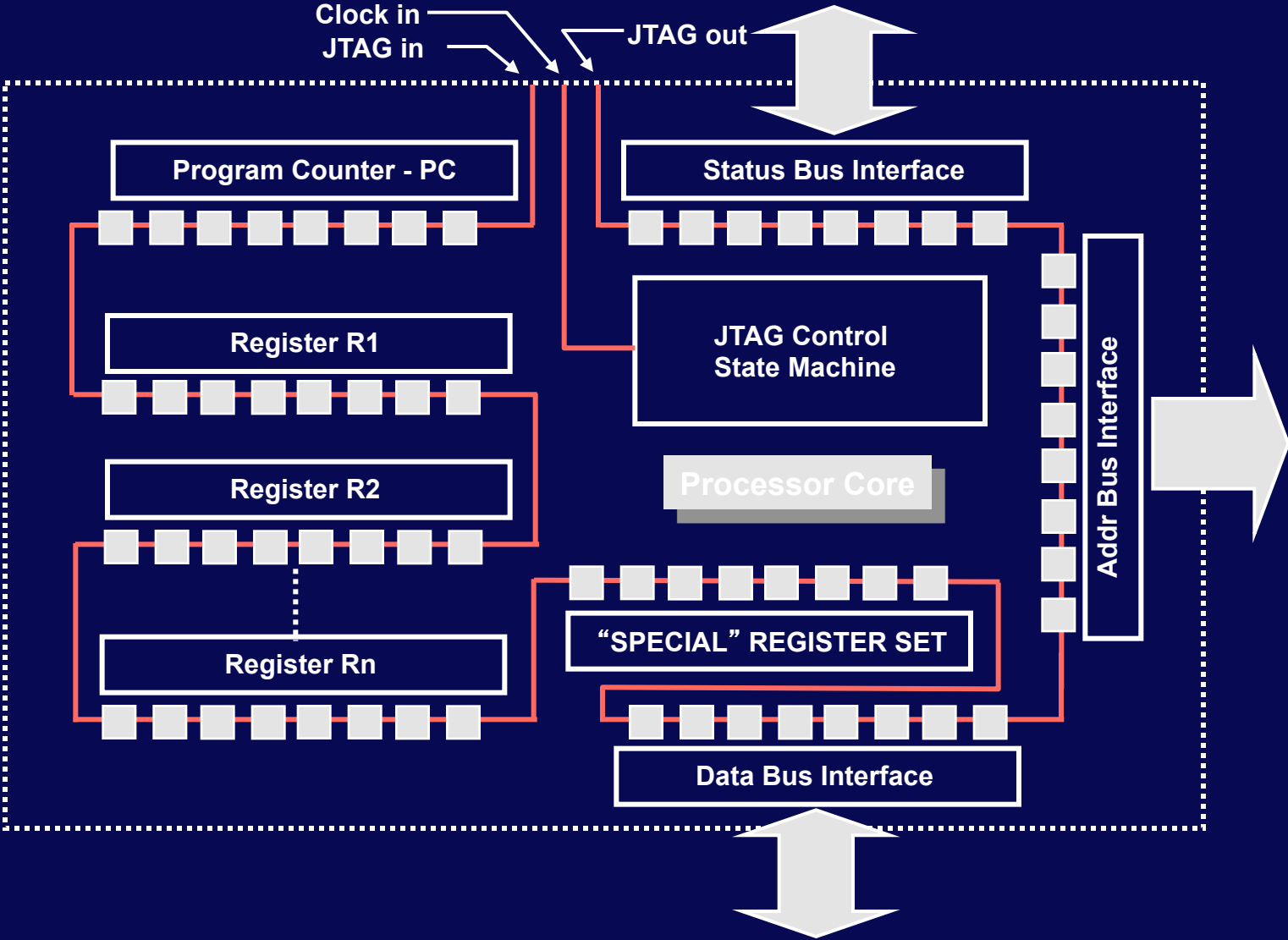
- ◆ Initially for hardware testing, evolved to support software testing
- ◆ Basic idea:
 - Each I/O pin, register, etc. can be “sniffed” by a JTAG cell
 - JTAG cells are connected in a “JTAG loop”
 - Contents of entire JTAG loop can be read using a shift register
 - Can also be written
 - External tool can reconstruct machine state from the JTAG bit stream

JTAG Hardware Debugging



Bit stream forms one long shift-register

JTAG Software Debugging



More JTAG

- ◆ **Advantages of shift-register approach:**
 - Simple
 - Requires few pins
- ◆ **Disadvantage of shift-register approach:**
 - End up reading and writing a lot of data just to change one register
- ◆ **JTAG optimizations:**
 - **Commands** – Directly change a single register or memory cell
 - **Addressable loops** – smaller JTAG loops each containing a subset of the machine state

More JTAG

- ◆ **Pins:**
 - TCK – clock
 - TDI – input data stream, sampled on rising edge of TCK
 - TDO – output data stream, updated on falling edge of TCK
 - TRST – Resets JTAG state machine (optional)
 - TMS – Test mode select: advances JTAG state machine
- ◆ **JTAG interface modules tend to be expensive**
 - “Low cost” solutions may be \$2000
 - However, all-software solutions (on the host side) exist
 - MCF52233 has JTAG

Summary

- ◆ **Embedded system development is strongly dependent on good tools**
 - **There is huge variation in tool quality**
 - **Lots of times free tools can be found**
 - **Sometimes they suck**
 - **Non-free tools can be really expensive**
 - **E.g., more than \$10K per developer seat**
 - **These can suck too**
- ◆ **You need to understand what the tools do, what the tradeoffs are, etc.**
- ◆ **Generally it's far better to buy the right tools up front**
 - **Saving \$\$ not worth if it makes the product ship late**