## Mid-Term Exam 2

## CS 3520, Fall 2005

## (take-home, but Fall 2009 exam will be in-class)

Name:

Instructions: You have ninety minutes to complete this open-book, open-note, closed-computer, take-home exam. Please write you start and finish times above, and write all answers in the provided space, plus the back of the exam if necessary. The Fall 2009 Mid-Term 2 exam will be in-class.

- 1) Which of the following produce different results in a call-by-value language and a call-by-name language? Both produce the same result if they both produce the same number or they both produce a procedure (even if the procedure doesn't behave exactly the same when applied).
  - a) {{fun {y} 12} {1 2}}
  - b) {fun {x} {{fun {y} 12} {1 2}}
  - c) {+ 1 {fun {y} 12}}
  - d) {+ 1 {{fun {x} {+ 1 13}} {+ 1 {fun {z} 12}}}
  - e) {+ 1 {{fun {x} {+ x 13}} {+ 1 {fun {z} 12}}}

2) The following web servlet implementation (main handler plus helper function) uses web-read, which takes only a prompt and uses let/cc internally to obtain a continuation. Convert the servlet (both functions) to instead use web-read/k, which takes a prompt and an explicit continuation procedure (and does not use let/cc internally). You should assume that the correct-password? function requires no interaction with the user. The Fall 2009 version of this question will be more difficult.

```
(define (pw-handler base args)
 (get-pw (web-read "Name")))
(define (get-pw name)
  (local [(define pw (web-read "Password"))]
   (if (correct-password? name pw)
        (format "Hello, ~a" name)
        (get-pw name))))
```

**3**) Given the following expression:

{{fun {x} {x x}} {fun {y} 12}}

Describe a trace of the evaluation in terms of arguments to interp and continue functions for every call of each. (There will be 7 calls to interp and 5 calls to continue.) The interp function takes three arguments — an expression, a substitution cache, and a continuation — so show all three for each interp call. The continue function takes two arguments — a value and a continuation — so show both for each continue call. Represent continuations using records. The Fall 2009 version of this question will involve the store-passing interpreter of HW6 instead of the interp-continue interpreter. The trace will be shorter than in Mid-Term 1.

- 4) Suppose a garbage-collected interepreter uses the following three kinds of records:
  - Tag 1: a record containing two pointers
  - Tag 2: a record containing one pointer and one integer
  - Tag 3: a record containing one integer

The interpreter has one register, which always contains a pointer, and a memory pool of size 22. The allocator/collector is a two-space copying collector, so each space is of size 11. Records are allocated consecutively in to-space, starting from the first memory location, 0.

The following is a snapshot of memory just before a collection where all memory has been allocated:

- Register: 8
- To space: 1 3 8 3 0 2 3 7 2 0 8

What are the values in the register and the new to-space (which is also addressed starting from 0) after collection? Assume that unallocated memory in to-space contains 0.

- Register:
- To space:

## Answers

**2**)

 $\mathbf{3})$ 

```
1) a and d.
     (define (pw-handler base args)
        (web-read/k "Name" get-pw))
     (define (get-pw name)
         (web-read/k "Password"
                        (lambda (pw)
                          (if (correct-password? name pw)
                               (format "Hello, ~a" name)
                               (get-pw name)))))
                  interp
                           expr
                                  =
                                       \{\{fun \{x\} \{x x\}\} \{fun \{y\} 12\}\}
                                       (mtSub)
                           subs
                                  =
                           k
                                  =
                                       (mtK)
                                       \{fun \{x\} \{x x\}\}
                  interp
                           expr
                                  =
                                       (mtSub)
                           subs
                                  =
                           k
                                       (appArgK | {fun {y} 12} | (mtSub) (mtK))
                                  =
                                       (closureV 'x | \{x x\} \rangle = v_1
                  \operatorname{cont}
                           val
                                  =
                           k
                                       (appArgK | {fun {y} 12} | (mtSub) (mtK))
                                  =
                  interp
                           expr
                                       {fun {y} 12}
                                  =
                           subs
                                       (mtSub)
                                  =
                           k
                                       (doAppK v_1 (mtK))
                                  =
                                       (closureV 'y 12) = v_2
                  \operatorname{cont}
                           val
                                  =
                           k
                                       (doAppK v_1 (mtK))
                                  =
                  interp
                           expr
                                  =
                                       \{x x\}
                           ds
                                  =
                                       (aSub 'x v_2 (mtSub)) = ds_1
                           k
                                       (mtK)
                                  =
                  interp
                           expr
                                  =
                                       x
                           ds
                                       ds_1
                                  =
                           k
                                  =
                                       (appArgk x ds_1 (mtK))
                  \operatorname{cont}
                           val
                                  =
                                       v_2
                           k
                                       (appArgK x ds_1 (mtK))
                                  =
                  interp
                           expr
                                       x
                                  =
                           ds
                                       ds_1
                                  =
                           k
                                       (doAppK v_2 (mtK))
                                  =
                           val
                  \operatorname{cont}
                                  =
                                       v_2
```

= (doAppK  $v_2$  (mtK)) k interp  $\operatorname{expr}$ 12 = ds (aSub 'y  $v_2$  (mtSub)) = (mtK) k =  $\operatorname{cont}$  $\operatorname{val}$ (numV 12) = (mtK) k =

4) Register: 0, To space:  $2\ 3\ 8\ 1\ 6\ 0\ 3\ 0\ 0\ 0$