

# Part I

# Allocation

Constructor calls are allocation:

```
(define (interp)
  (type-case ExpD expr-reg
    ....
    [(lamD body-expr)
     (begin
       (set! v-reg (closV body-expr env-reg))
       (continue))]
    ....))

(define (continue)
  ....
  [(plusSecondK r env k)
   (begin
     (set! expr-reg r)
     (set! env-reg sc)
     (set! k-reg (doPlusK v-reg k))
     (interp))]
  ....)
```

# Deallocation

Where does **free** go?

```
(define (continue)
  ....
  [(doPlusK v1 k)
   (begin
    (set! v-reg (num+ v1 v-reg))
    (free k-reg) ; ???
    (set! k-reg k)
    (continue))])
  ....
  [(doAppK fun-val k)
   (begin
    (set! expr-reg (closV-body fun-val))
    (set! env-reg (cons v-reg
                        (closV-env fun-val)))

    (set! k-reg k)
    (free fun-val) ; ???
    (interp))]
  ....)
```

## Deallocation

```
[ (doPlusK v1 k)
  (begin
    (set! v-reg (num+ v1 v-reg))
    (free k-reg) ; ???
    (set! k-reg k)
    (continue))] ]
```

- Without `let/cc`, this `free` is fine, because the continuation can't be referenced anywhere else
- A continuation object is always freed as `(free k-reg)`, which is why many language implementations use a stack

## Deallocation

```
[(doAppK fun-val k)
 (begin
  (set! expr-reg (closV-body fun-val))
  (set! env-reg (cons v-reg
                     (closV-env fun-val)))

  (set! k-reg k)
  (free fun-val) ; ???
  (interp))]
```

- This free is *not* ok, because the closure might be kept in a environment somewhere
- Need to free only if no one else is using it...

## Code and Data

An **object** is any record allocated during `interp` and `continue`

Assume that expressions are allocated “statically”

- `compile` uses `code-malloc1`, etc.
- Only try to free objects allocated during `interp` and `continue`

## Part 2

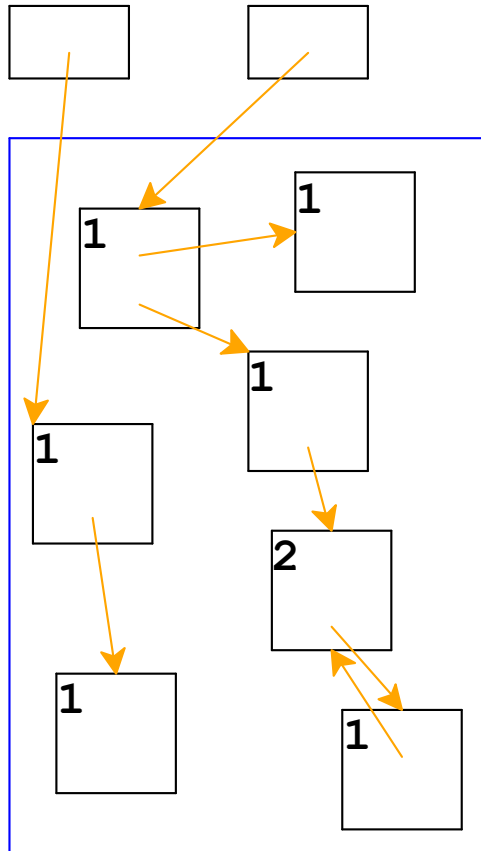
# Reference Counting

**Reference counting:** a way to know whether an object has other users

- Attach a count to every object, starting at 0
- When installing a pointer to an object (into a register or another object), increment its count
- When replacing a pointer to an object, decrement its count
- When a count is decremented to 0, decrement counts for other objects referenced by the object, then free



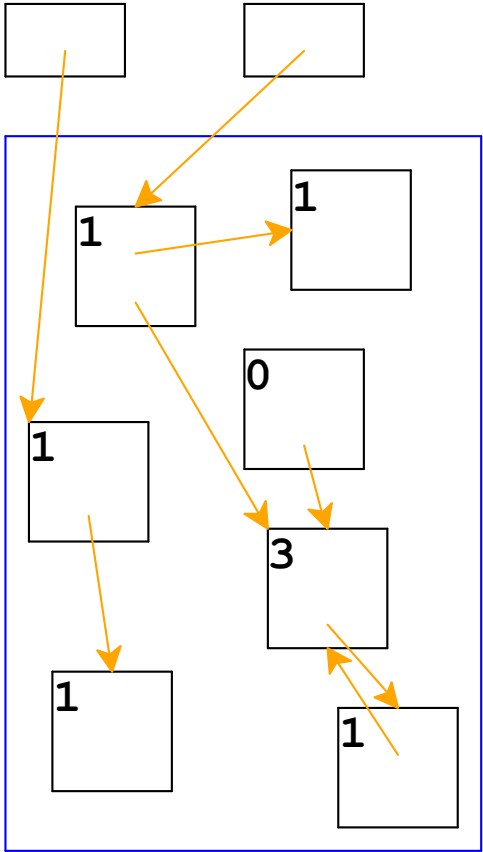
# Reference Counting



Top boxes are the registers  
**k-reg**, **v-reg**, etc.

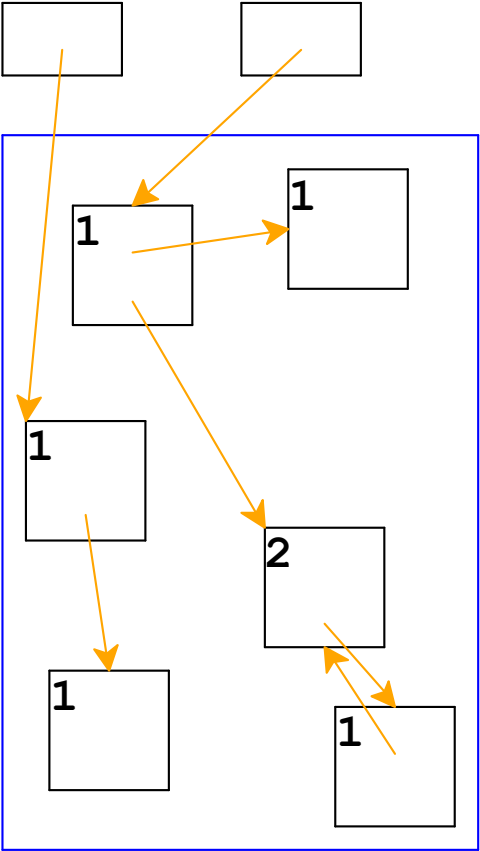
Boxes in the blue area are  
allocated with **malloc**

# Reference Counting



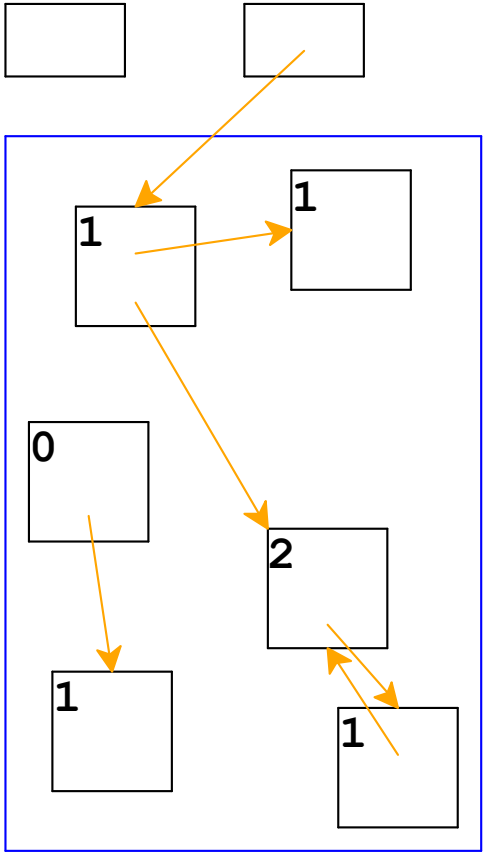
Adjust counts when a pointer is changed...

# Reference Counting



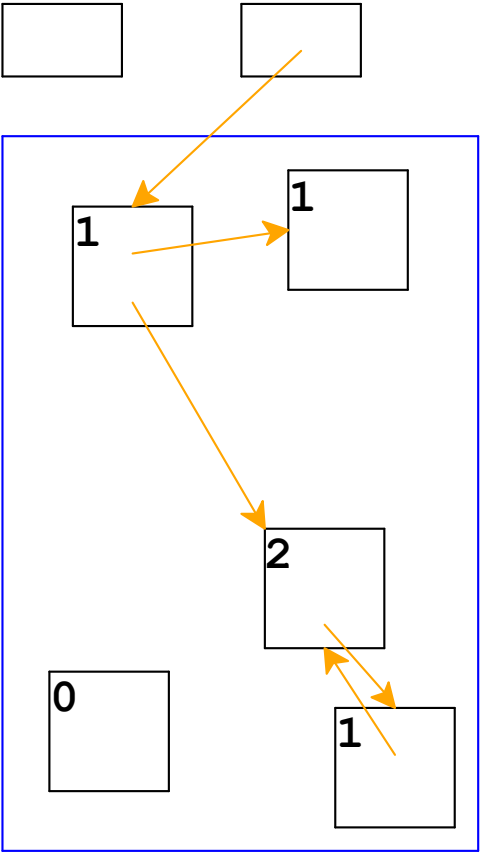
... freeing an object if its count goes to 0

# Reference Counting



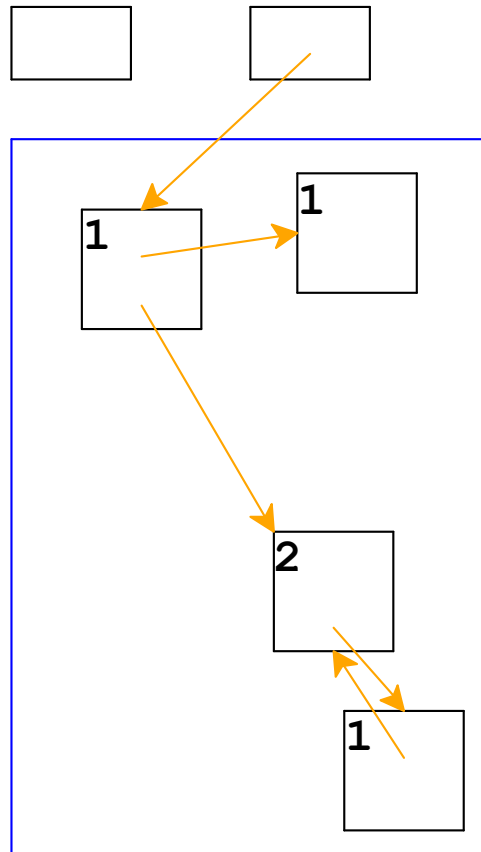
Same if the pointer is in a register

# Reference Counting



Adjust counts after frees, too...

# Reference Counting

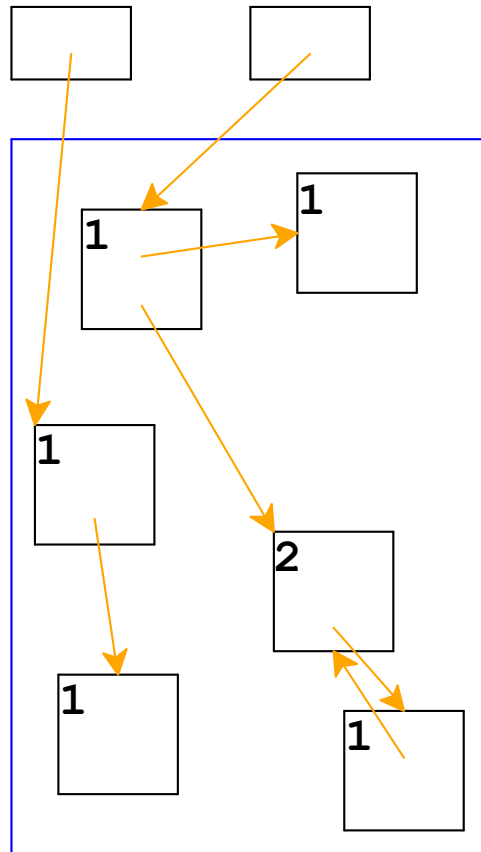


... which can trigger more frees

## Reference Counting in an Interpreter

```
...
[(lamE body-expr)
 (begin
  (ref- v-reg)
  (set! v-reg
   ; must ref+ env:
   (closV body-expr env-reg))
  (ref+ v-reg)
  (continue))]
...
[(doAppK fun-val k)
 (begin
  (set! expr-reg (closV-body fun-val)) ; code is static
  (ref- env-reg)
  (set! env-reg
   ; must ref+ each arg:
   (cons v-reg (closV-env fun-val)))
  (ref+ env-reg) ; => ref+ on v-reg
  (ref+ k)
  (ref- k-reg) ; => ref- on fun-val and k
  (set! k-reg k)
  (interp))]
```

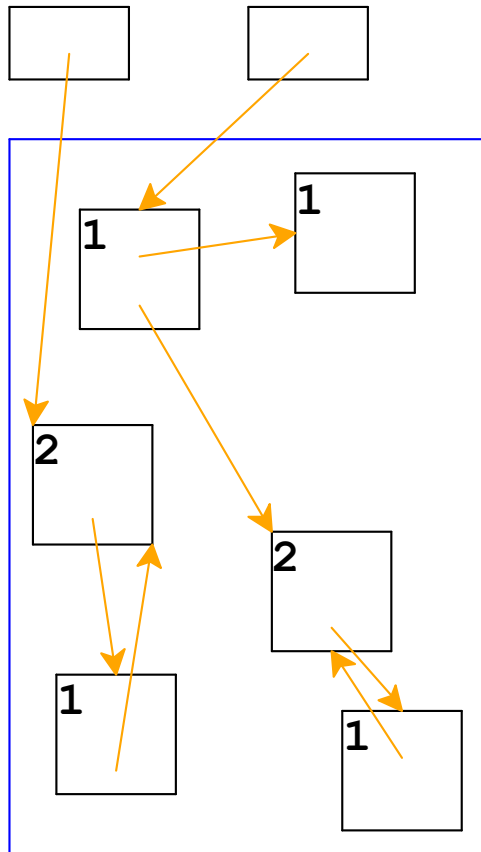
# Reference Counting And Cycles



An assignment can create a cycle...

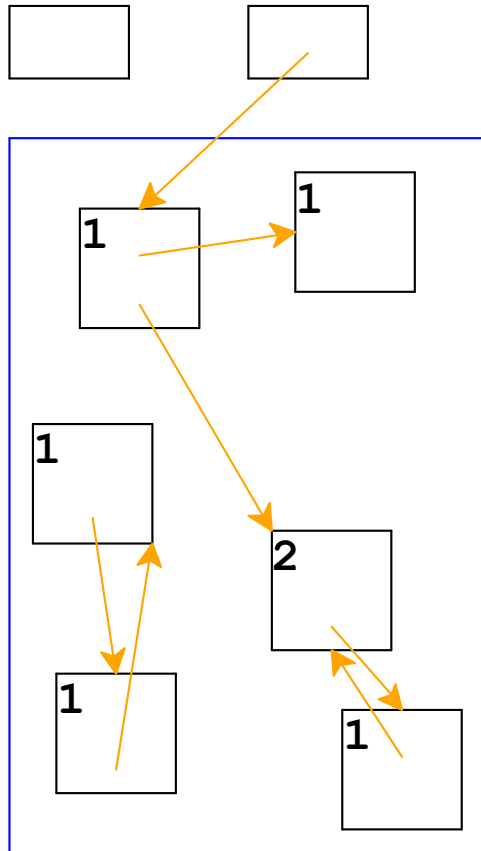


# Reference Counting And Cycles



Adding a reference increments a count

# Reference Counting And Cycles



Lower-left objects are inaccessible, but not deallocated

In general, cycles break reference counting

## Part 3

# Garbage Collection

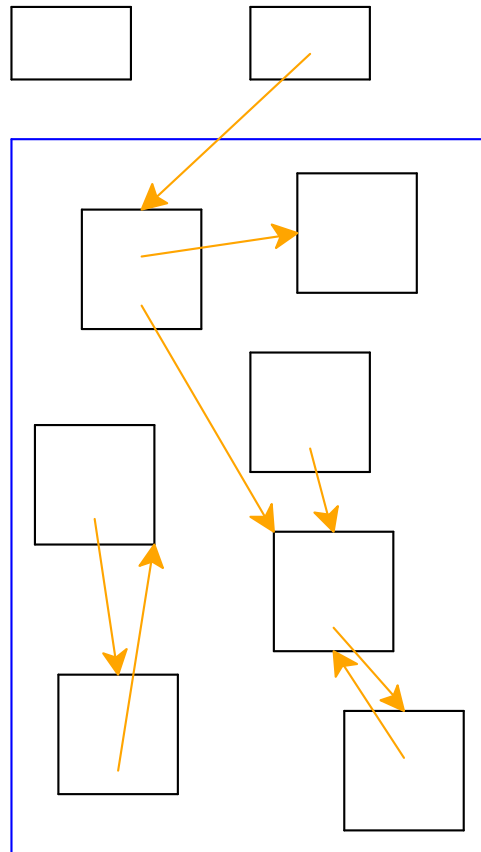
**Garbage collection:** a way to know whether an object is *accessible*

- An object referenced by a register is **live**
- An object referenced by a live object is also live
- A program can only possibly use live objects, because there is no way to get to other objects
- A garbage collector frees all objects that are not live
- Allocate until we run out of memory, then run a garbage collector to get more space

## Garbage Collection Algorithm

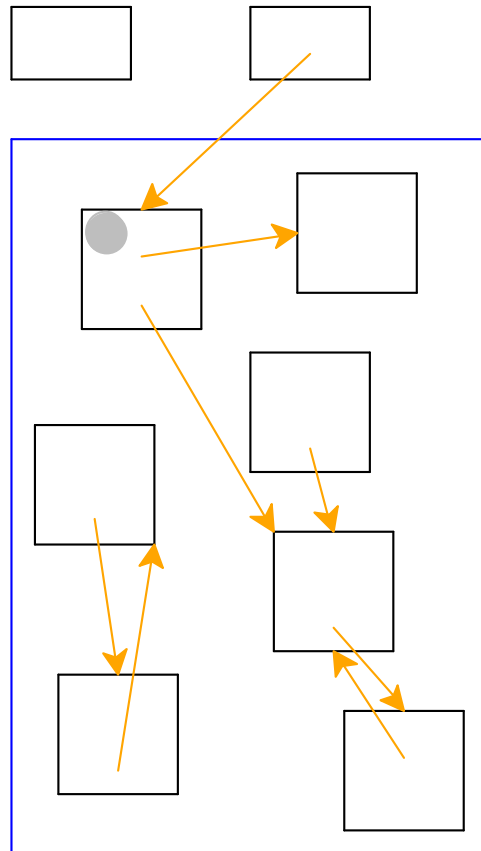
- Color all objects **white**
- Color objects referenced by registers **gray**
- Repeat until there are no gray objects:
  - Pick a gray object,  $r$
  - For each white object that  $r$  points to, make it gray
  - Color  $r$  **black**
- Deallocate all white objects

# Garbage Collection



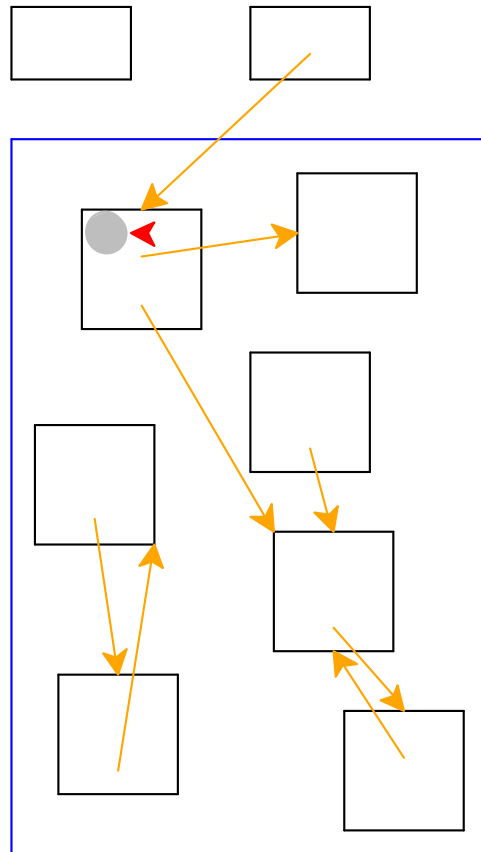
All objects are marked white

# Garbage Collection



Mark objects referenced by registers as gray

# Garbage Collection

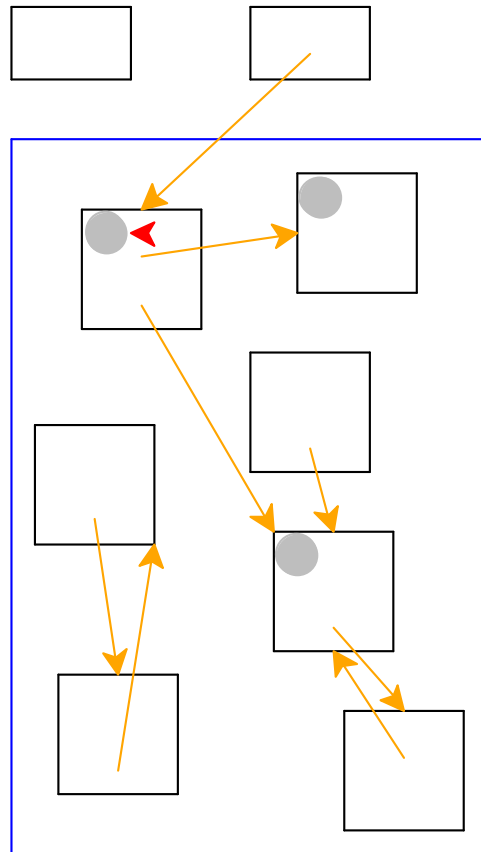


Need to pick a gray object

Red arrow indicates the chosen object



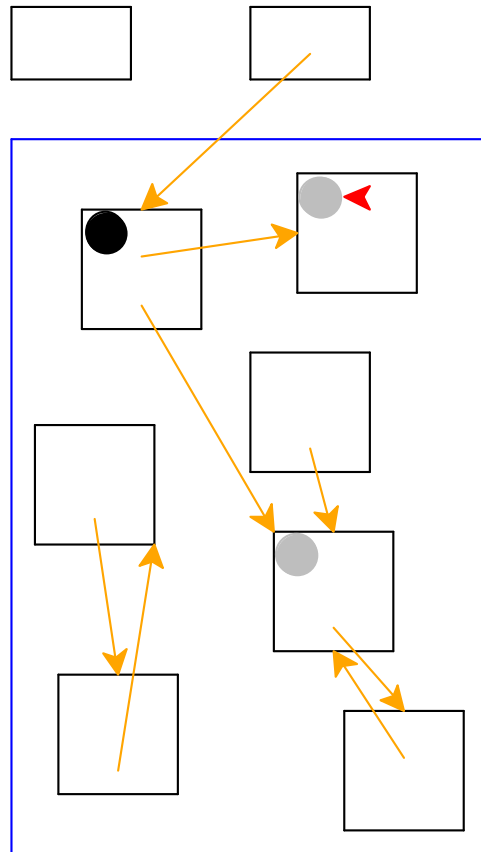
# Garbage Collection



Mark white objects referenced by  
chosen object as gray

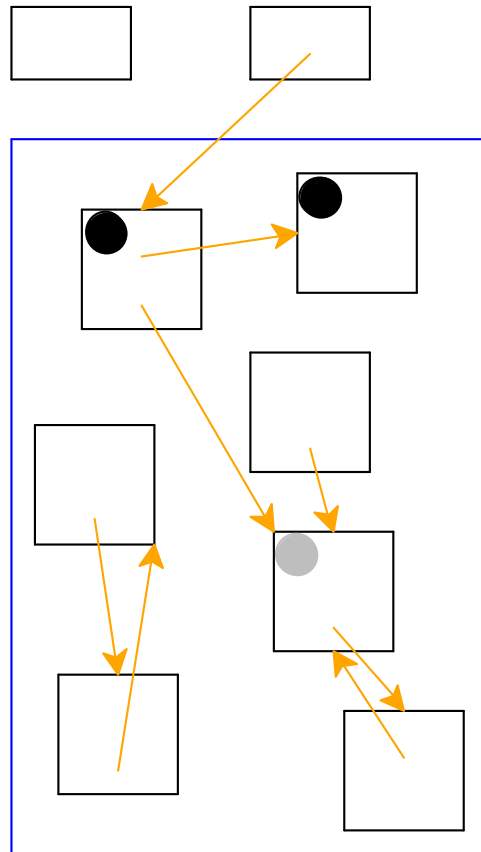


# Garbage Collection



Start again: pick a gray object

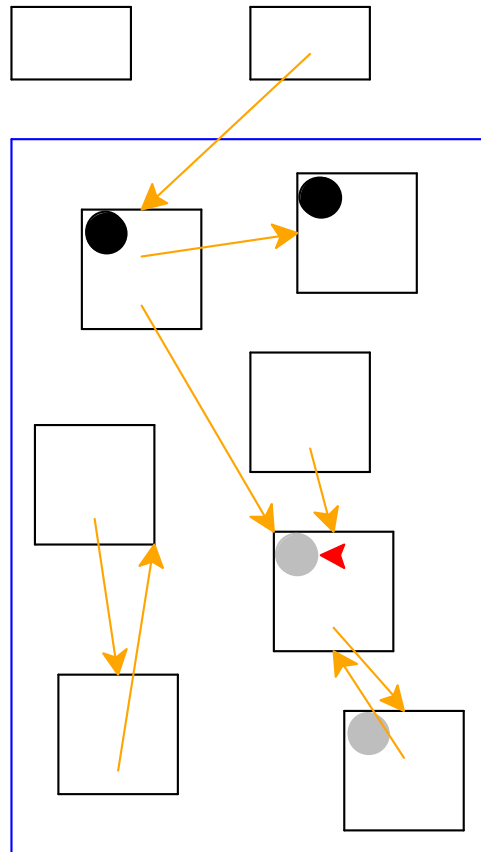
# Garbage Collection



No referenced objects; mark black

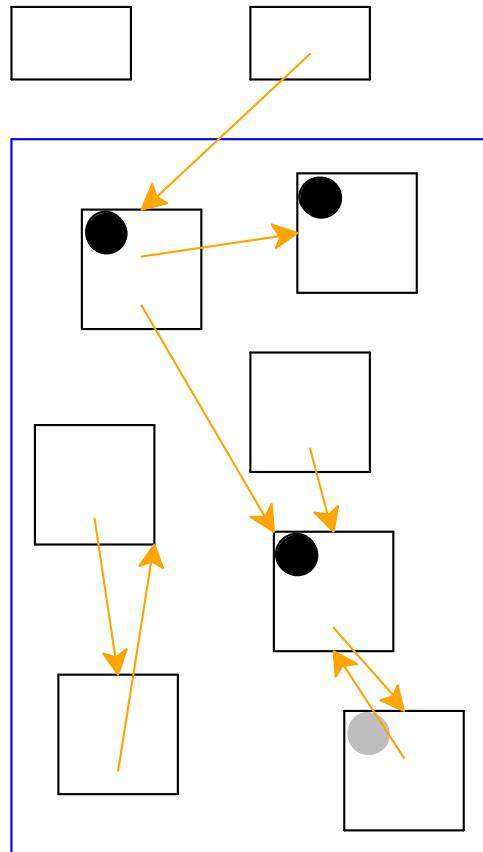


# Garbage Collection



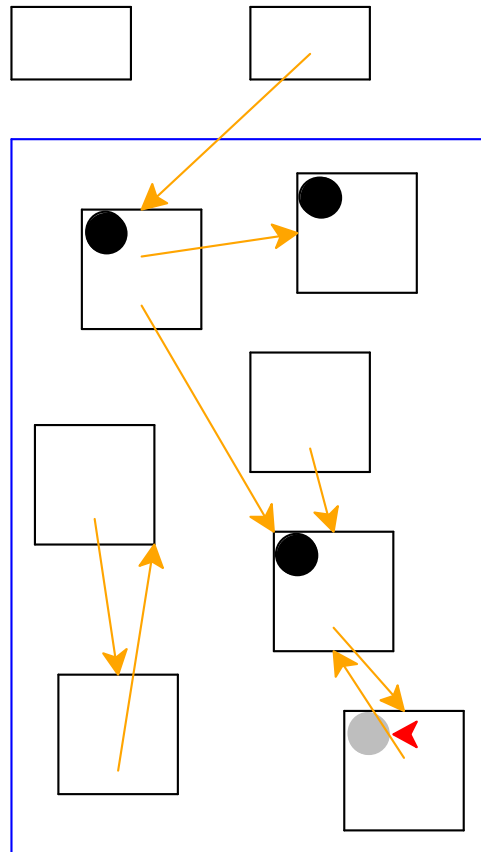
Mark white objects referenced by chosen object as gray

# Garbage Collection



Mark chosen object black

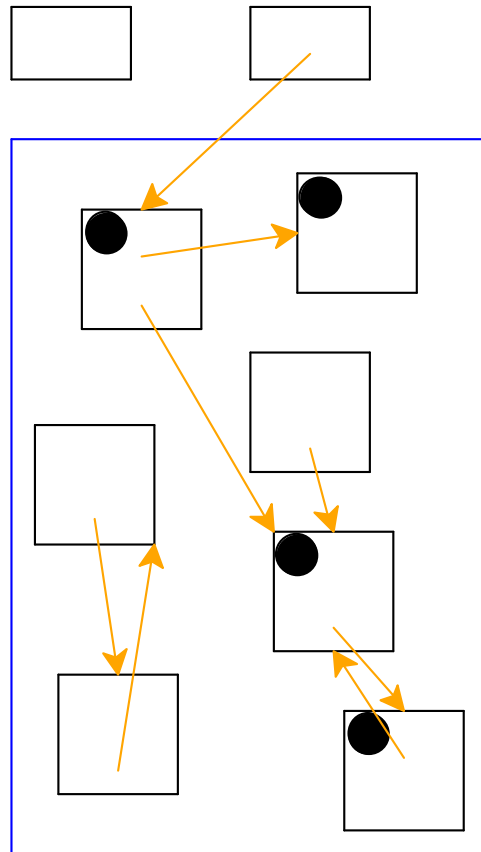
# Garbage Collection



Start again: pick a gray object

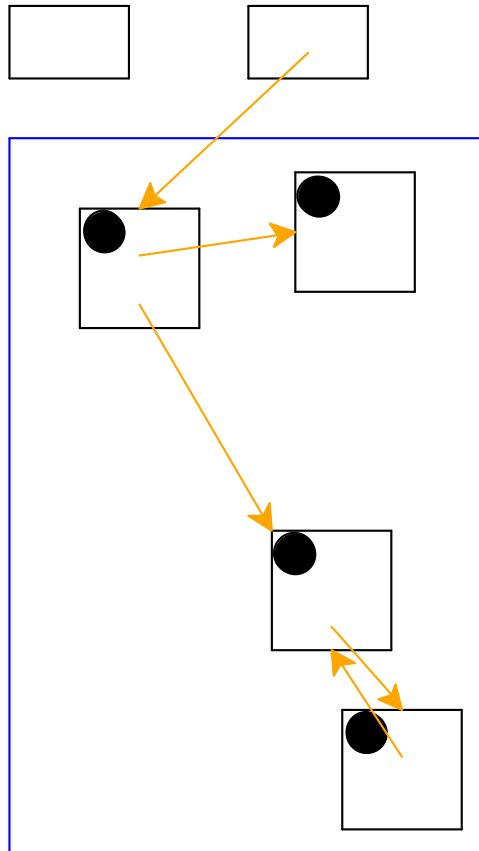


# Garbage Collection



No referenced white objects;  
mark black

# Garbage Collection



No more gray objects; deallocate white objects

Cycles **do not** break garbage collection

## Part 4

## Two-Space Copying Collectors

A **two-space** copying collector compacts memory as it collects, making allocation easier.

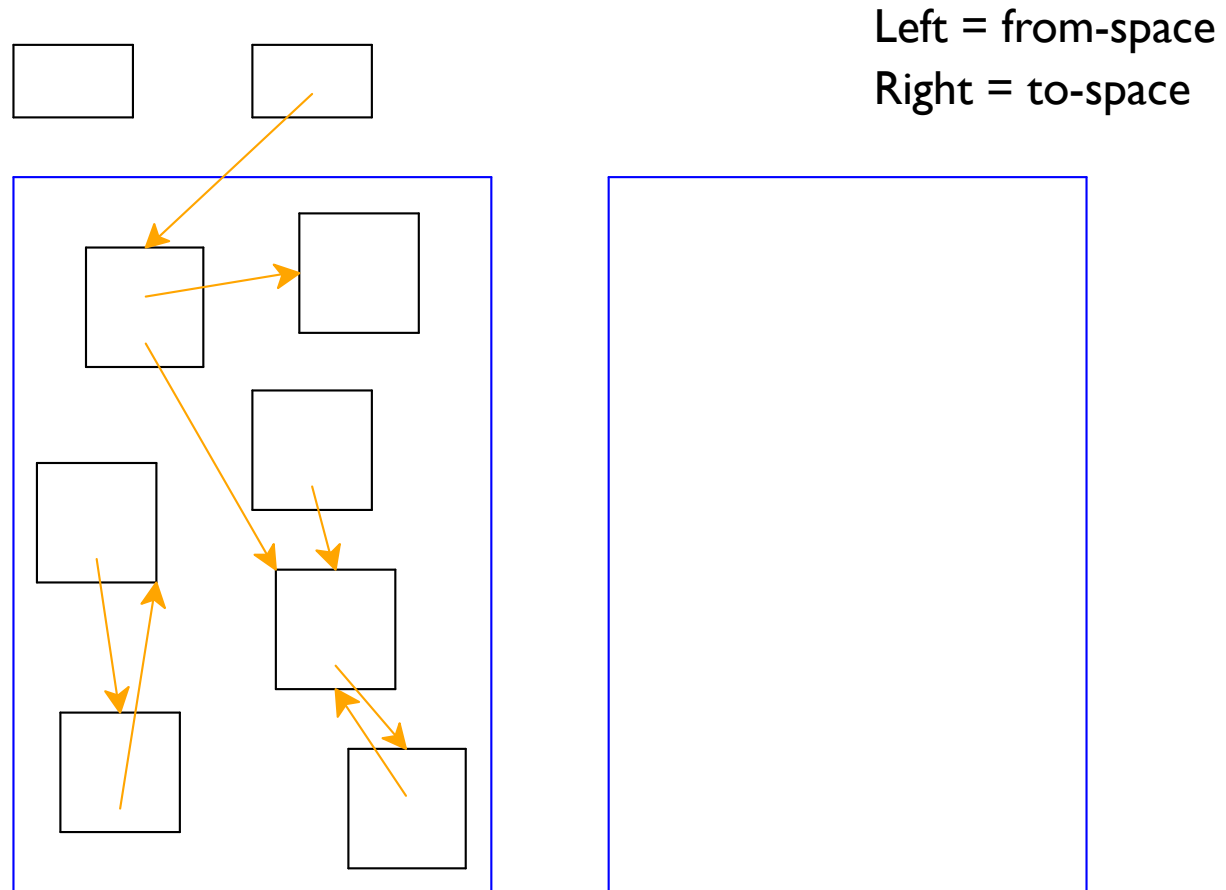
### Allocator:

- Partitions memory into **to-space** and **from-space**
- Allocates only in **to-space**

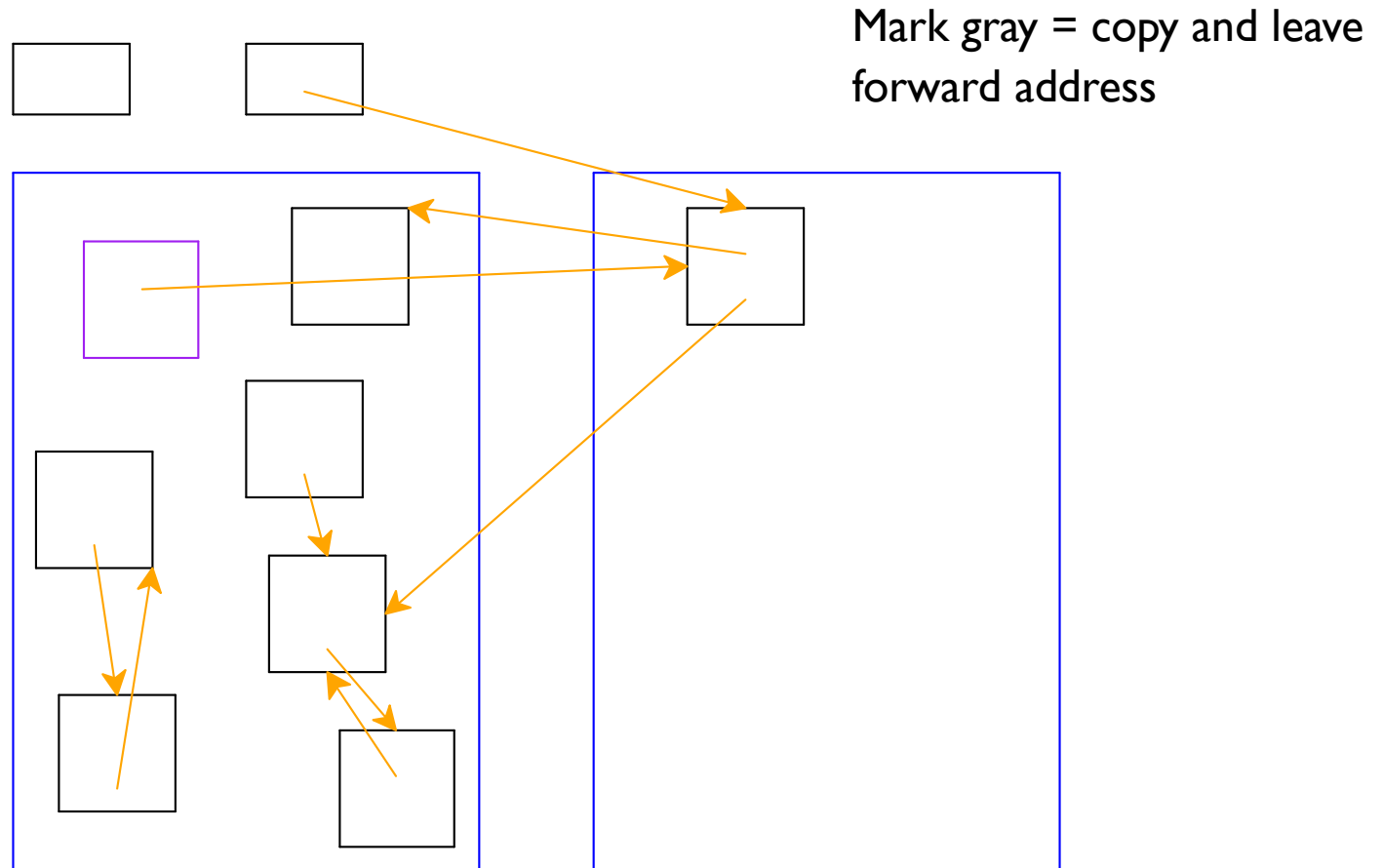
### Collector:

- Starts by swapping **to-space** and **from-space**
- Coloring gray  $\Rightarrow$  copy from **from-space** to **to-space**
- Choosing a gray object  $\Rightarrow$  walk once through the new **to-space**, update pointers

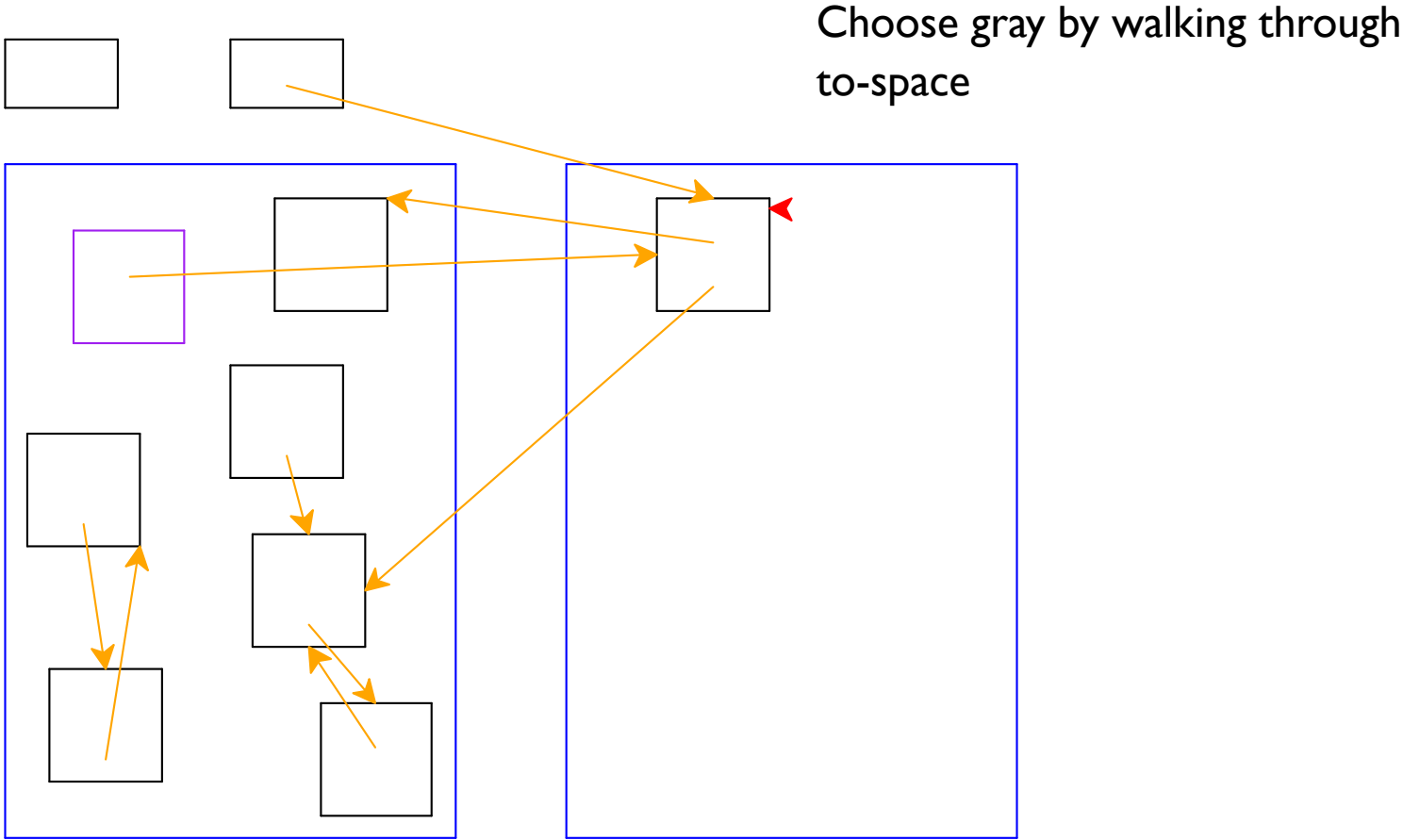
# Two-Space Collection



# Two-Space Collection

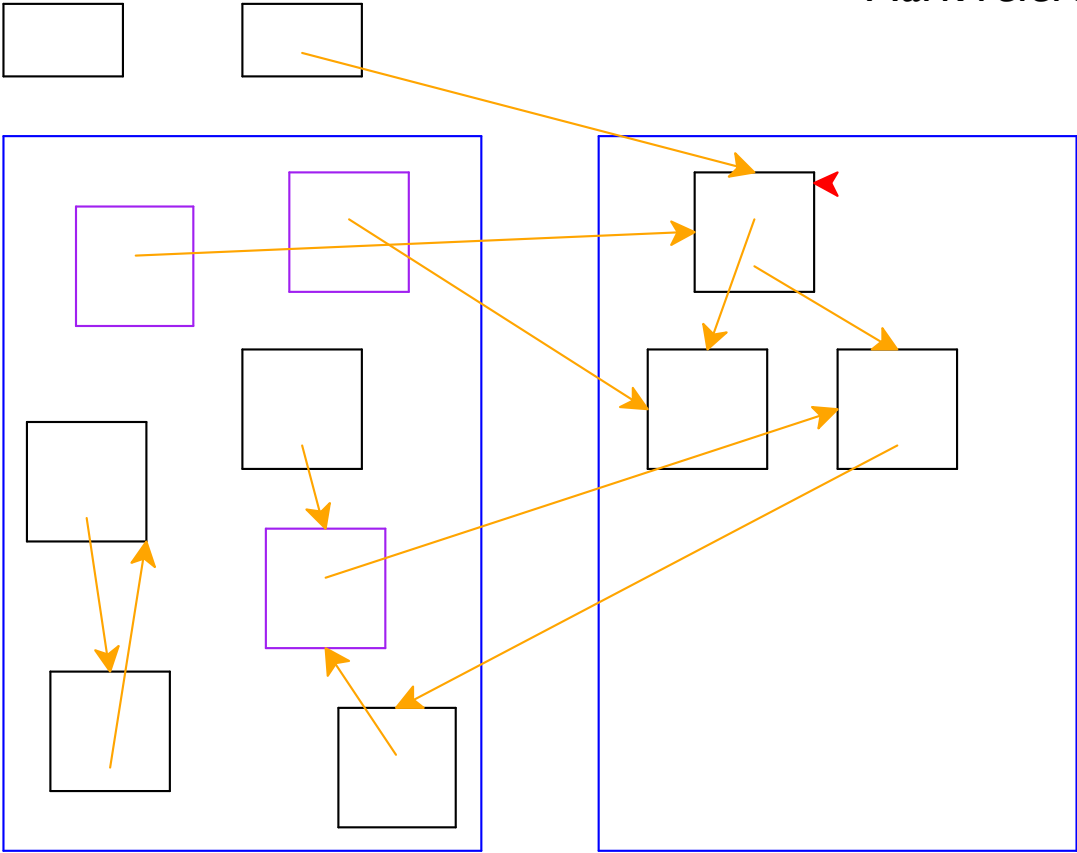


# Two-Space Collection



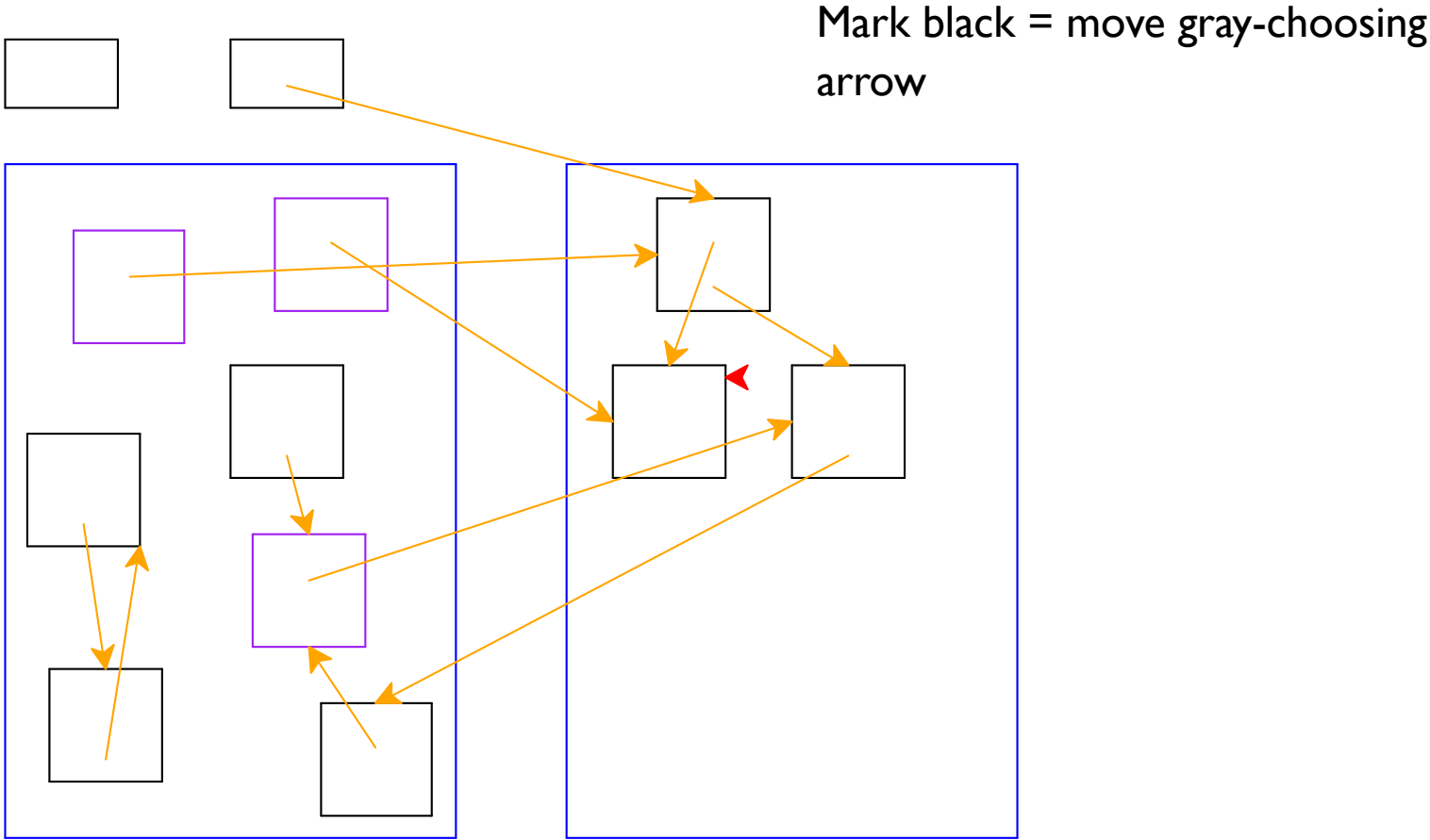
# Two-Space Collection

Mark referenced as gray

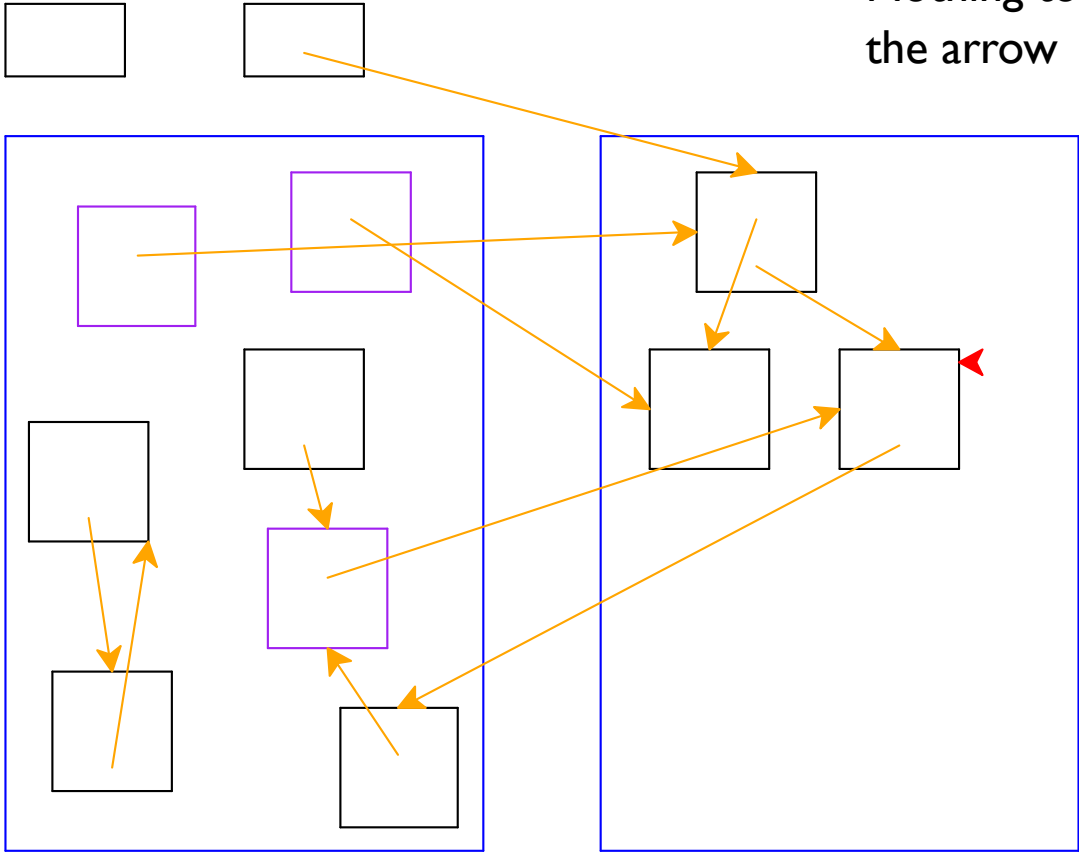




# Two-Space Collection



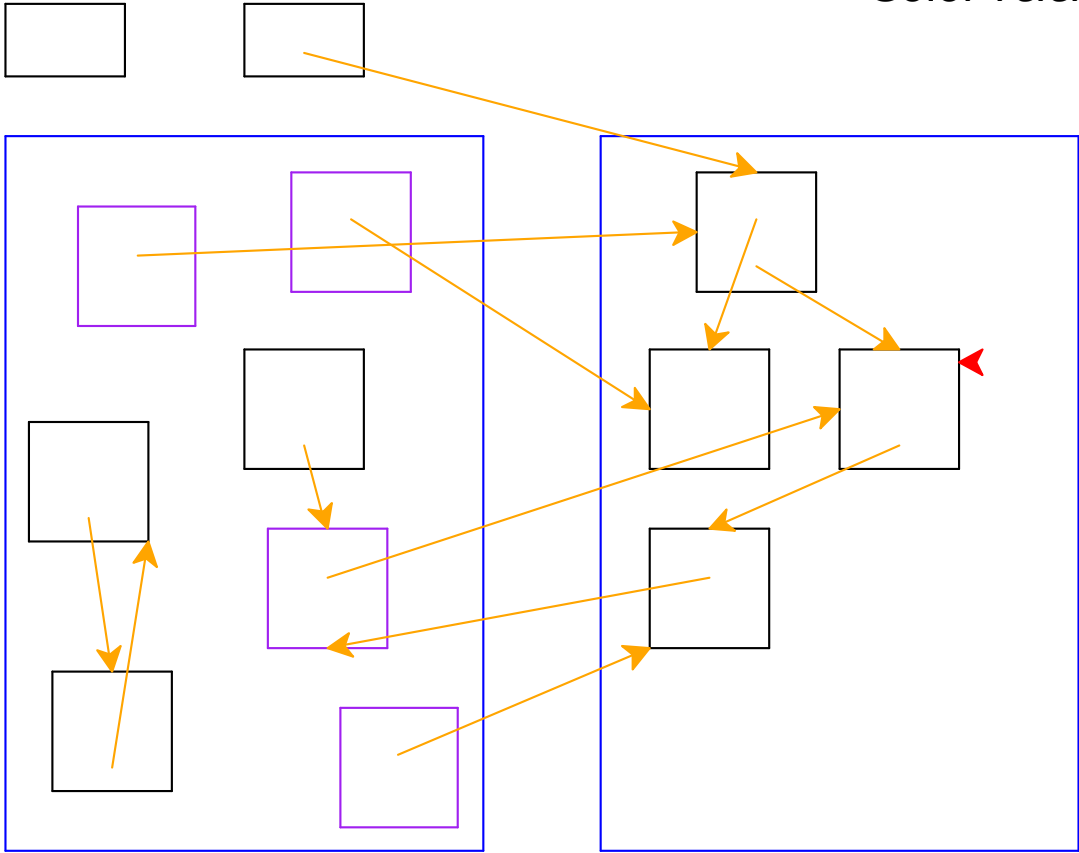
# Two-Space Collection



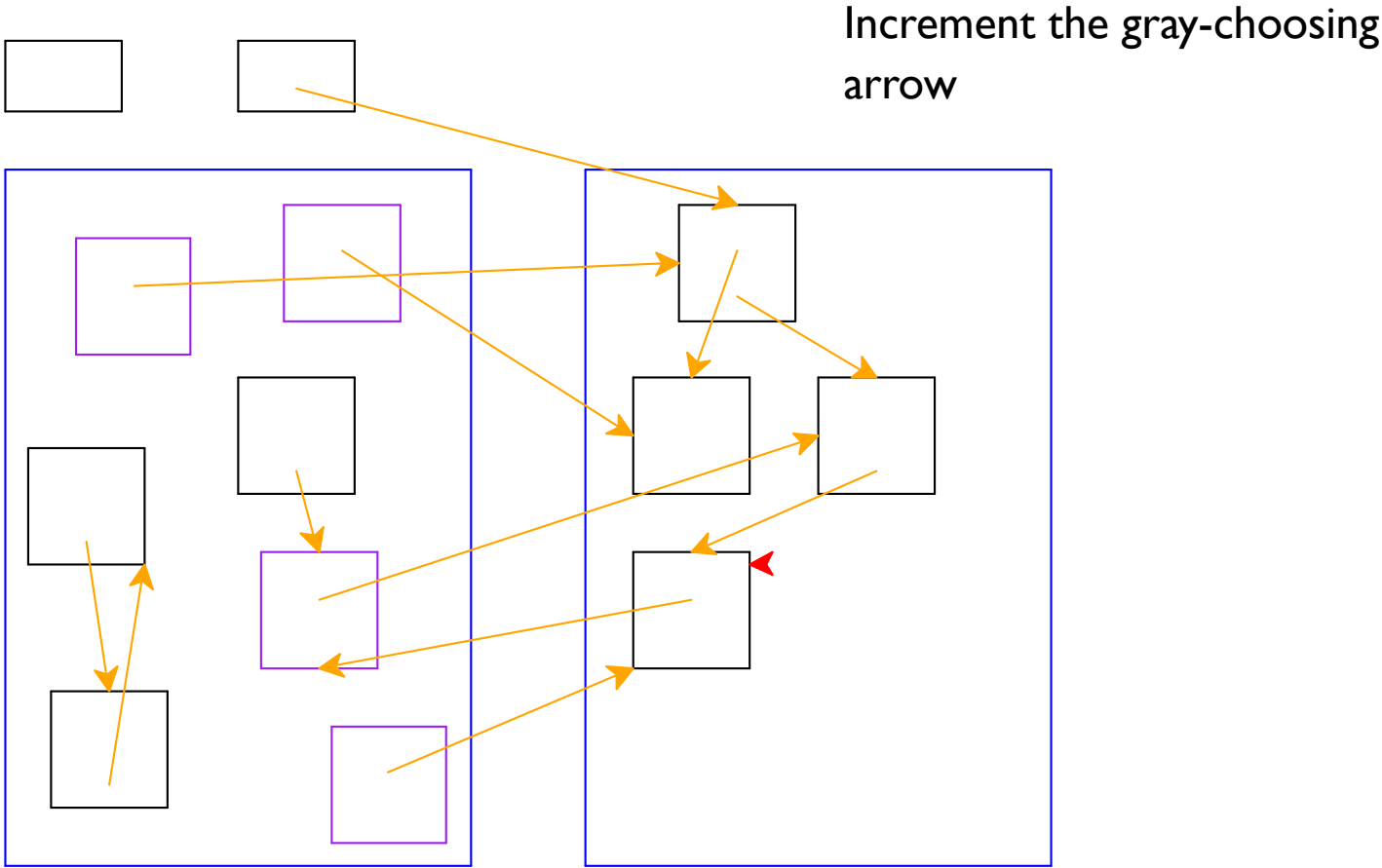
Nothing to color gray; increment the arrow

# Two-Space Collection

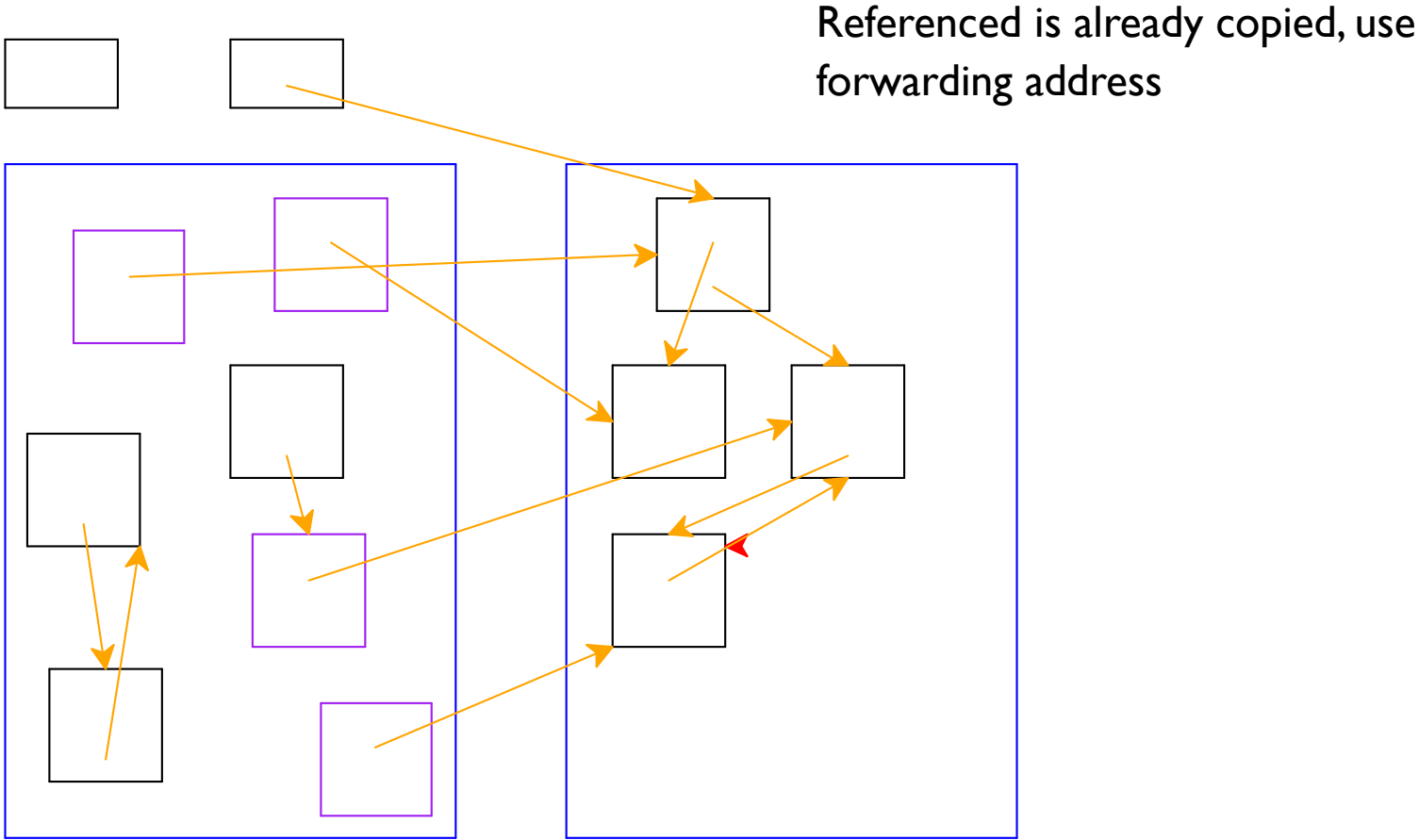
Color referenced object gray



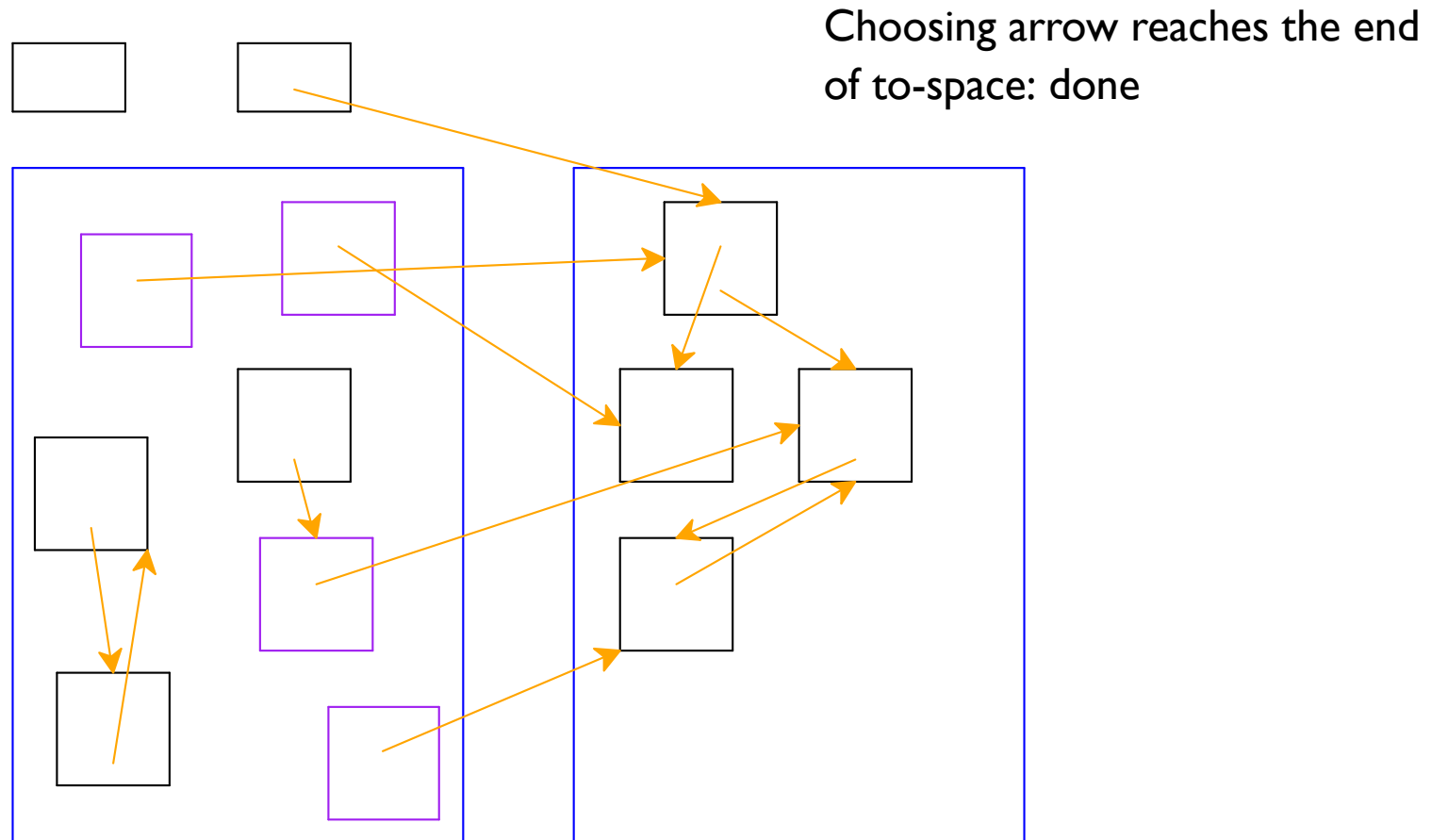
# Two-Space Collection



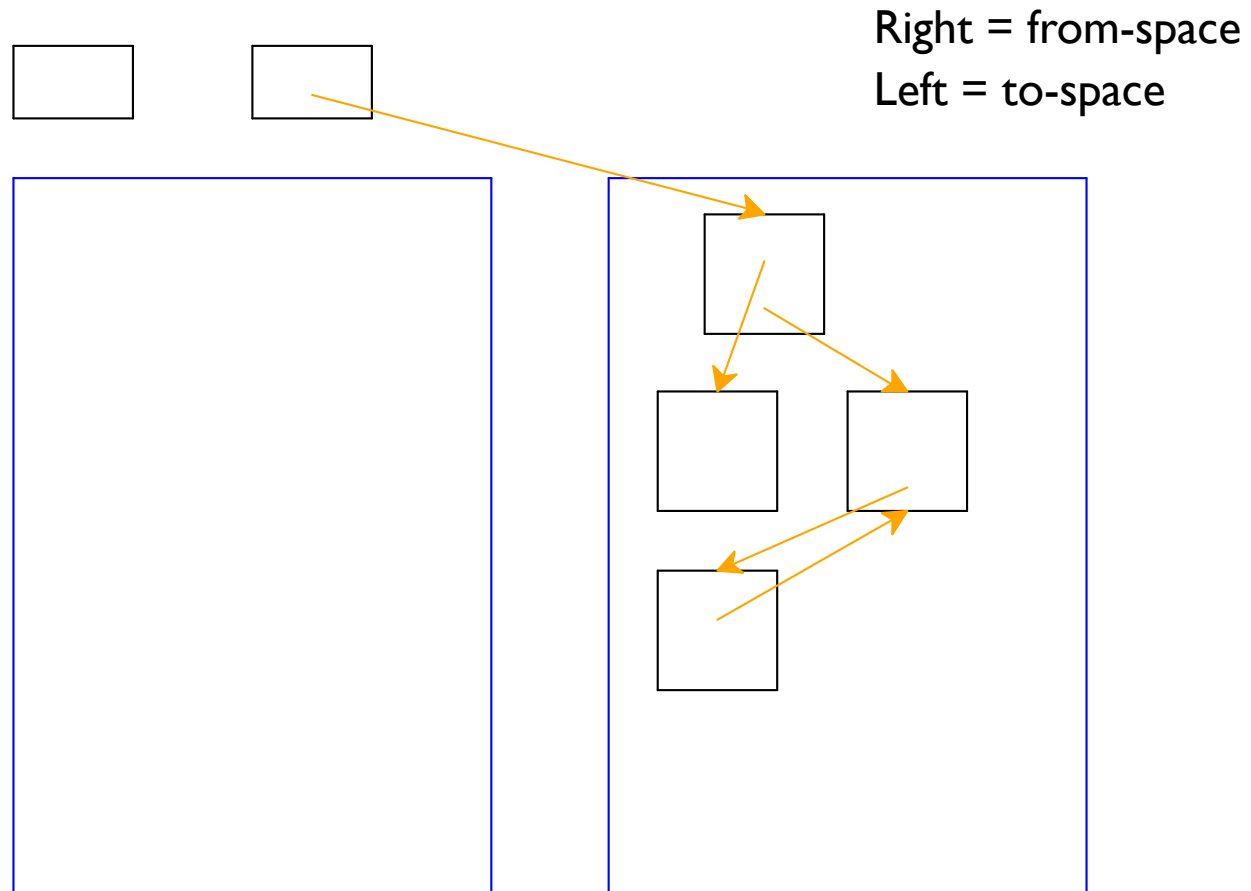
# Two-Space Collection



# Two-Space Collection



# Two-Space Collection



## Part 5



## Two-Space Collection on Vectors

- Everything is a number:
  - Some numbers are immediate integers
  - Some numbers are pointers
- An allocated object in memory starts with a tag, followed by a sequence of pointers and immediate integers
  - The tag describes the shape

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

Register 1: 7

Register 2: 0

From: 1 75 2 0 3 2 10 3 2 2 3 1 4

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 7						Register 2: 0						
From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 7						Register 2: 0						
From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 7						Register 2: 0						
From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	0	0	0	0	0	0	0	0	0	0	0	0	0
	^												
	^												

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 0						
From:	1	75	2	0	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	2	0	0	0	0	0	0	0	0	0	0
	^												
				^									

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 3						
From:	99	3	2	0	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	2	1	75	0	0	0	0	0	0	0	0
	^												
						^							

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0							Register 2: 3					
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	0	0	0	0	0	0	0
				^									
								^					



## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0							Register 2: 3					
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	0	0	0	0	0	0	0
						^							
								^					

## Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0							Register 2: 3					
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	3	0	0	0	0	0	0
								^					
								^					