

# Part I

# Values

A **value** is the result of an **expression**

- Expression: `1 + 2`
- Value: `3`

A value can be be  
the argument to a function,  
the right-hand side of a **let**,  
...

## Functions as Values?

Is a function a value in Moe?

**No**

You can define a function

```
fun double(x) : x + x
```

You can call a function

```
double(10)
```

You *cannot* use a function name without calling it

You *cannot* pass a function to another function

## Functions as Values?

Is a function a value in Shplait?

**Yes**

An expression can produce a function result

```
fun double(x) : x + x  
double
```

```
[min, max]
```

```
fun (x) : x + x
```

You can pass a function to a function:

```
map(fun (x) : x + x,  
     [1, 2, 3])
```

## Why Functions as Values

Abstraction is easier with functions as values

- `filter`, `map`, `foldl`, etc.

Separate `fun` definition form becomes unnecessary

```
fun f(x) : 1 + x
f(10)
```

⇒

```
let f = (fun (x) : 1 + x) :
  f(10)
```

## Why Functions as Values

Abstraction is easier with functions as values

- `filter`, `map`, `foldl`, etc.

Separate `fun` definition form becomes unnecessary

```
fun f(x) : 1 + x
f(10)
```

⇒

Historical name: `lambda` or  $\lambda$

```
let f = (fun (x) : 1 + x) :
  f(10)
```

## Part 2

## New Moe Grammar, Almost

```
<Exp> ::= <Int>
        | <Symbol>
        | <Exp> + <Exp>
        | <Exp> * <Exp>
        | let <Symbol> = <Exp>: <Exp>
        | <Symbol> (<Exp>)
        | fun (<Symbol>) : <Exp>
```

\*

NEW



## Evaluation

`10`  $\Rightarrow$  `10`

`y`  $\Rightarrow$  *free variable*

`1 + 2`  $\Rightarrow$  `3`

`2 * 3`  $\Rightarrow$  `6`

`let x = 7: x + 2`  $\Rightarrow$  `7 + 2`  $\Rightarrow$  `9`

`fun (x): 1 + x`  $\Rightarrow$  `fun (x): 1 + x`

Result is not always a number!

~~`interp :: (Exp, ....) -> Int`~~

`interp :: (Exp, ....) -> Value`

## Evaluation

`10`  $\Rightarrow$  `10`

`y`  $\Rightarrow$  *free variable*

`1 + 2`  $\Rightarrow$  `3`

`2 * 3`  $\Rightarrow$  `6`

`let x = 7: x + 2`  $\Rightarrow$  `7 + 2`  $\Rightarrow$  `9`

`fun (x): 1 + x`  $\Rightarrow$  `fun (x): 1 + x`

`let y = 10: fun (x): y + x`

$\Rightarrow$  `fun (x): 10 + x`

`let f = (fun (x): 1 + x): f(3)`

$\Rightarrow$  `(fun (x): 1 + x)(3)`

Doesn't match the grammar for `<Exp>`

# New Moe Grammar

```
<Exp> ::= <Int>  
        | <Symbol>  
        | <Exp> + <Exp>  
        | <Exp> * <Exp>  
        | let <Symbol> = <Exp>: <Exp>  
        | <Symbol> (<Exp>)  
        | fun (<Symbol>): <Exp>  
        | <Exp> (<Exp>)
```



## Evaluation

```
let f = (fun (x) : 1 + x) : f(3)
```

```
⇒ (fun (x) : 1 + x) (3)
```

```
⇒ 1 + 3 ⇒ 4
```

```
(fun (x) : 1 + x) (3) ⇒ 1 + 3 ⇒ 4
```

```
1 (2) ⇒ not a function
```

```
1 + (fun (x) : 10) ⇒ not a number
```

## Part 3

## Expression Datatype

```
type Exp
| intE(n :: Int)
| idE(s :: Symbol)
| plusE(l :: Exp,
        r :: Exp)
| multE(l :: Exp,
        r :: Exp)
| letE(n :: Symbol,
      rhs :: Exp,
      body :: Exp)
| funE(n :: Symbol,
      body :: Exp)
| appE(fn :: Exp,
      arg :: Exp)
```

```
check: parse('fun (x) : x + 1')
       ~is funE('#x, plusE(idE('#x), intE(1)))
```

## Expression Datatype

```
type Exp
| intE(n :: Int)
| idE(s :: Symbol)
| plusE(l :: Exp,
        r :: Exp)
| multE(l :: Exp,
        r :: Exp)
| letE(n :: Symbol,
      rhs :: Exp,
      body :: Exp)
| funE(n :: Symbol,
      body :: Exp)
| appE(fn :: Exp,
      arg :: Exp)
```

```
check: parse('f(10)')
       ~is appE(idE('#f'), intE(10))
```

## Expression Datatype

```
type Exp
| intE(n :: Int)
| idE(s :: Symbol)
| plusE(l :: Exp,
        r :: Exp)
| multE(l :: Exp,
        r :: Exp)
| letE(n :: Symbol,
      rhs :: Exp,
      body :: Exp)
| funE(n :: Symbol,
      body :: Exp)
| appE(fn :: Exp,
      arg :: Exp)
```

```
check: parse('(fun (x): x + 1)(10)')
       ~is appE(funE('#x', plusE(idE('#x'), intE(1))),
               intE(10))
```



## Part 4

## Functions with Substitutions

```
interp( let y = 10:  
        fun (x) : y + x )
```

## Functions with Substitutions

```
interp ( let y = 10:  
         fun (x) : y + x )
```

## Functions with Substitutions

```
interp ( let y = 10:  
         fun (x) : y + x )
```

⇒

```
fun (x) : 10 + x
```

## Functions with Substitutions

```
interp ( let y = 10 : fun (x) : y + x )
```

⇒

```
fun (x) : 10 + x
```

## Functions with Deferred Substitution

```
interp (let y = 10: fun (x): y + x)
```

⇒

```
interp (fun (x): y + x)
```

y = 10

## Functions with Deferred Substitution

```
interp( (let y = 10: fun (x): y + x) (let y = 7: y) )
```

Argument expression:

```
interp( let y = 7: y )
```

⇒

```
interp( y = 7 y ) ⇒ 7
```

Function expression:

```
interp( let y = 10: fun (x): y + x )
```

⇒

```
interp( y = 10 fun (x): y + x ) ⇒ ?
```

## Functions with Deferred Substitution

```
interp( (let y = 10: fun (x): y + x) (let y = 7: y) )
```

Argument expression:

```
interp( let y = 7: y )
```

⇒

```
interp( y = 7 y ) ⇒ 7
```

Function expression:

```
interp( let y = 10: fun (x): y + x )
```

⇒

```
interp( y = 10 fun (x): y + x ) ⇒ ?
```

A **closure** combines an expression with an environment



## Representing Values

```
type Value
| intV(n :: Int)
| closV(arg :: Symbol,
        body :: Exp,
        env :: Env)
```

```
type Binding
| bind(name :: Symbol,
       val :: Value)
```

```
check: interp(let y = 10: fun (x): y + x,
             mt_env)
~is closV(#'x, y + x,
         extend_env(bind(#'y, intV(10)),
                   mt_env))
```

## Continuing Evaluation

Argument: `interp(y)` y = intV(7)  
⇒ `intV(7)`

Function: `interp(fun (x) : y + x)` y = intV(10)  
⇒ `closV('#x, y + x, extend_env(bind('#y, intV(10)), mt_env))`

To apply, interpret the function body with the given argument:

`interp(y + x)` x = intV(7) y = intV(10)

## Part 5

# Interpreter

```
fun interp(a :: Exp, env :: Env) :: Value:
  match a
  | intE(n): intV(n)
  | idE(s): lookup(s, env)
  | plusE(l, r): num_plus(interp(l, env), interp(r, env))
  | multE(l, r): ....
  | letE(n, rhs, body):
    ....
  | funE(n, body): ....
  | appE(fn, arg):
    ....
```

## Add and Multiply

```
fun num_plus(l :: Value, r :: Value) :: Value:
  cond
  | l is_a intV && r is_a intV:
    intV(intV.n(l) + intV.n(r))
  | ~else:
    error(#'interp, "not a number")

fun num_mult(l :: Value, r :: Value) :: Value:
  cond
  | l is_a intV && r is_a intV:
    intV(intV.n(l) * intV.n(r))
  | ~else:
    error(#'interp, "not a number")
```

## Add and Multiply

```
fun num_op(l, r, op):  
  cond  
  | l is_a intV && r is_a intV:  
    intV(op(intV.n(l), intV.n(r)))  
  | ~else:  
    error('#interp, "not a number")  
  
fun num_plus(l :: Value, r :: Value) :: Value:  
  num_op(l, r, fun (ln, rn): ln + rn)  
  
fun num_mult(l :: Value, r :: Value) :: Value:  
  num_op(l, r, fun (ln, rn): ln * rn)
```

# Interpreter

```
fun interp(a :: Exp, env :: Env) :: Value:
  match a
  | intE(n): intV(n)
  | idE(s): lookup(s, env)
  | plusE(l, r): num_plus(interp(l, env), interp(r, env))
  | multE(l, r): num_mult(interp(l, env), interp(r, env))
  | letE(n, rhs, body):
      interp(body, extend_env(bind(n, interp(rhs, env)),
                              env))
  | funE(n, body): closV(n, body, env)
  | appE(fn, arg):
      match interp(fn, env)
      | closV(n, body, c_env):
          interp(body,
                 extend_env(bind(n, interp(arg, env)),
                             c_env))
      | ~else: error(#'interp, "not a function")
```